## Temporal Implications of Database Information Accountability

Kyriacos E. Pavlou and Richard T. Snodgrass

Computer Science Department The University of Arizona TIME 2012



#### Motivation

## Motivation

Corporate collusion has given rise to regulations for trustworthy long-term data management.

 Code of Federal Regulations of FDA: Clinical trials

## Motivation

Corporate collusion has given rise to regulations for trustworthy long-term data management.

- Code of Federal Regulations of FDA: Clinical trials
- Sarbanes-Oxley Act: Financial transactions
- HIPAA Health Insurance Portability and Accountability Act; Canada's PIPEDA: Disclosure of medical information

• File systems & DB communities

- File systems & DB communities
  - tamper detection / prevention mechanisms

- File systems & DB communities
  - tamper detection / prevention mechanisms
- Audit log security & compliant records
  - Creation
  - Storage
  - Access
  - Maintenance / Retention

- File systems & DB communities
  - tamper detection / prevention mechanisms
- Audit log security & compliant records
  - Creation
  - Storage
  - Access
  - Maintenance / Retention

Governed by laws & regulations



#### Information Restriction

immutable retained records access control

#### Information Restriction

immutable retained records access control

#### Information **Accountability**

transparent information set of rules easily determine appropriate use



immutable retained records access control

#### Information Accountability

transparent information set of rules easily determine appropriate use

"[Information] accountability must become a primary means through which society addresses appropriate use." (Weitzner et al., CACM 2008)

### **Restriction vs Accountability**

## **Restriction vs Accountability**

- Home Security
  - Locked doors and windows (restriction)
  - Sweeping front yard, cameras (accountability)

# **Restriction vs Accountability**

- Home Security
  - Locked doors and windows (restriction)
  - Sweeping front yard, cameras (accountability)
- Bank Security
  - The vault is unlocked during business hours.
  - Easy access
  - CCTV cameras everywhere

• Tried and tested idea

- Tried and tested idea
- Example: Bullae, sigils, seals, etc

- Tried and tested idea
- Example: Bullae, sigils, seals, etc



- Tried and tested idea
- Example: Bullae, sigils, seals, etc



- Tried and tested idea
- Example: Bullae, sigils, seals, etc



- Tried and tested idea
- Example: Bullae, sigils, seals, etc



- Tried and tested idea
- Example: Bullae, sigils, seals, etc



- Tried and tested idea
- Example: Bullae, sigils, seals, etc



• Tried and tested idea

- Tried and tested idea
- Example: Fair Credit Reporting Act (1970)



No rules on the collection of data and analysis but on their use (credit report).

- Tried and tested idea
- Example: Fair Credit Reporting Act (1970)



No rules on the collection of data and analysis but on their use (credit report).

The consumers are allowed access to the data.

- Tried and tested idea
- Example: Fair Credit Reporting Act (1970)



No rules on the collection of data and analysis but on their use (credit report).

The consumers are allowed access to the data.

Agencies using credit reports are accountable for their decisions.

• Tried and tested idea

- Tried and tested idea
- Example: Creative Commons Licensing

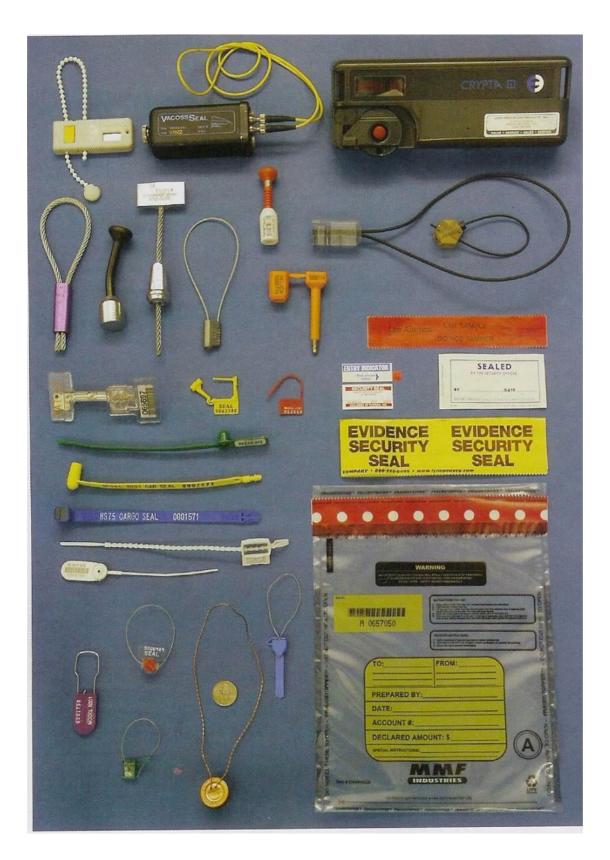


- Tried and tested idea
- Example: Creative Commons Licensing

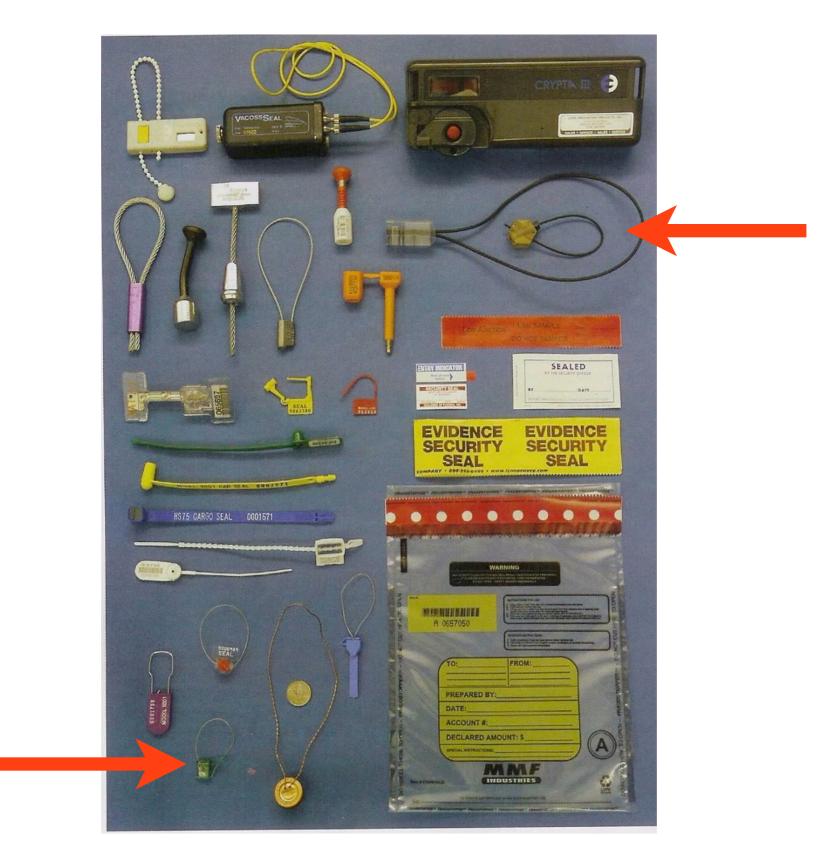


Do not attempt to prevent the lawful use of works they protect by using technology, but rather set forth rules regulating the use of the works.

#### Tamper-Indicating Seals for Nuclear Safeguarding



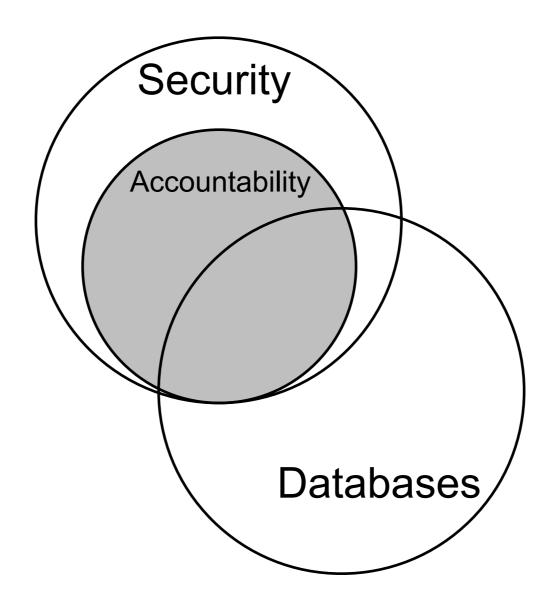
#### Tamper-Indicating Seals for Nuclear Safeguarding



# Accountability \begin{bmatrix} Databases \end{bmatrix}



# Accountability \begin{bmatrix} Databases \end{bmatrix}



Methodology accountability

restriction



accountability

restriction

coarse

**EMC**<sup>2</sup>

where information lives'

bc

NetApp<sup>\*</sup>



restriction

#### Fossilization (Hsu & Ong)



Investigative Data Mining (Mena)

Fossilization (Hsu & Ong)



accountability

Methodology

restriction

Forensic Server Project (Carvey & Kleiman)

Investigative Data Mining (Mena)

Fossilization (Hsu & Ong)



restriction

Forensic Server Project (Carvey & Kleiman)

Investigative Data Mining (Mena) Indexing Structures (Goodrich et al.)

Fossilization (Hsu & Ong)



restriction

Forensic Server Project (Carvey & Kleiman)

Investigative Data Mining (Mena) Indexing Structures (Goodrich et al.)

Fossilization (Hsu & Ong)

ECCC where information lives betAppp\* SarbOx Workflows (Agrawal et al.)

restriction

Forensic Server Project (Carvey & Kleiman)

Investigative Data Mining (Mena) Indexing Structures (Goodrich et al.)

Fossilization (Hsu & Ong)

<image>

SarbOx Workflows (Agrawal et al.)

ORACLE Total Recall

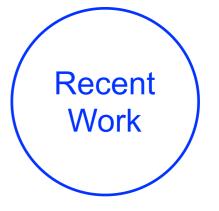


**Data Granularity** 

restriction

Forensic Server Project (Carvey & Kleiman)

Investigative Data Mining (Mena) Indexing Structures (Goodrich et al.)



restriction

accountability

#### Fossilization (Hsu & Ong)

<image>

SarbOx Workflows (Agrawal et al.)

ORACLE Total Recall



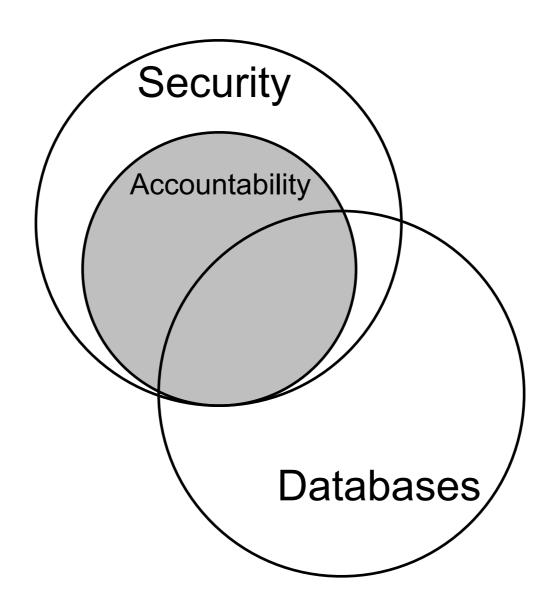
fine

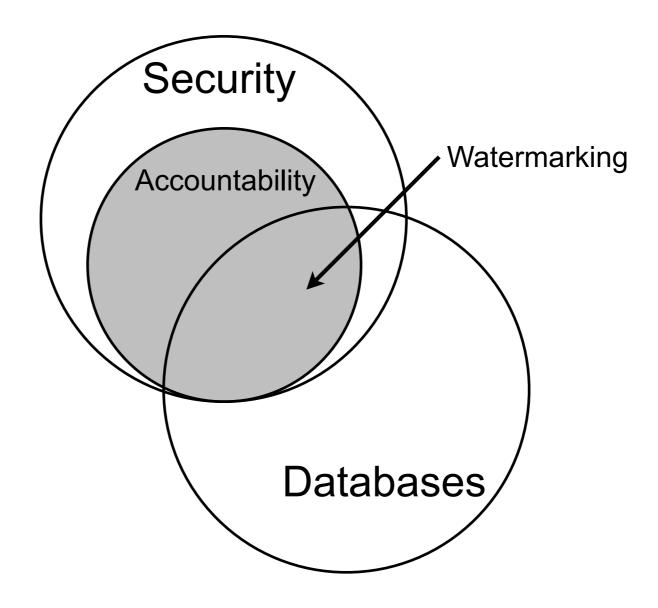
#### Info Accountability of Fine-Grained Data

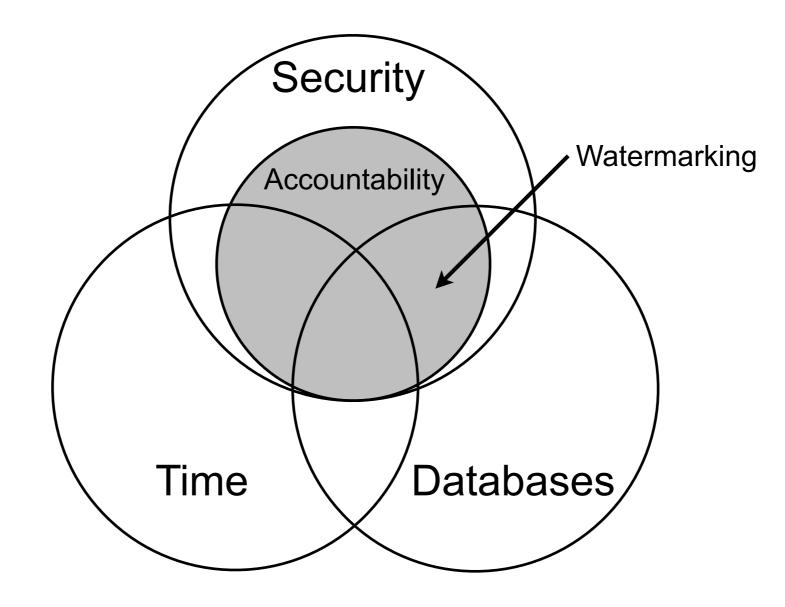
 Fragile watermarking scheme for detecting malicious modifications of database relations [Guo, Li, Liu, and Jajodia 2006].

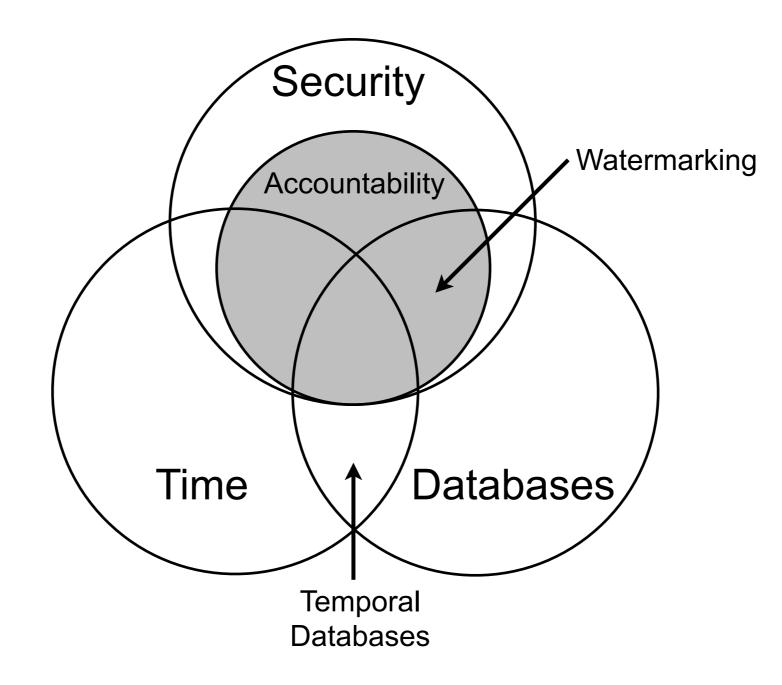
- Efficient audit-based compliance for relational data retention [Hasan, Winslett, and Mitra 2009].
- Tamper detection in audit logs [Snodgrass, Yao, and Collberg 2004].

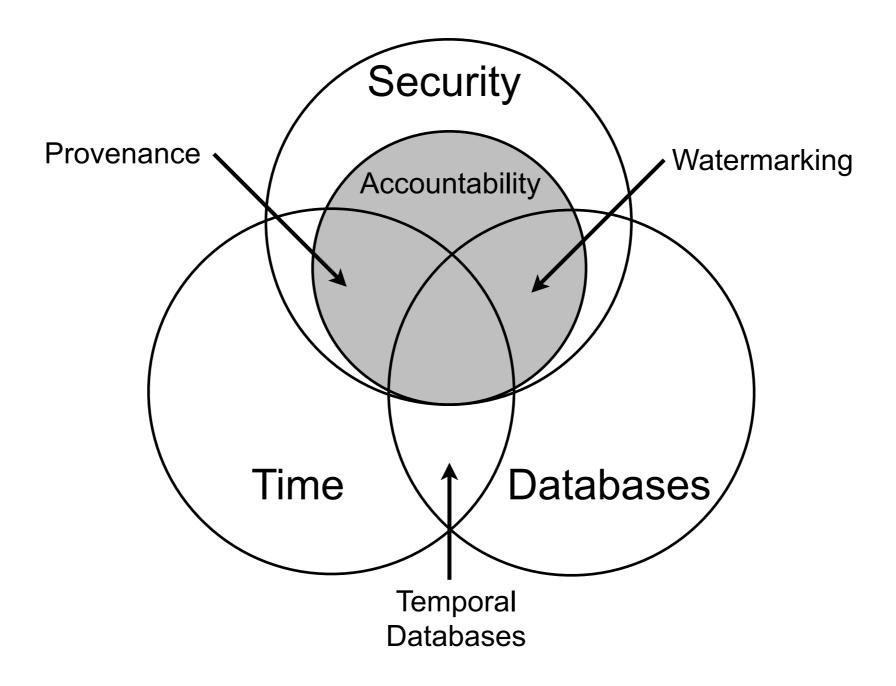


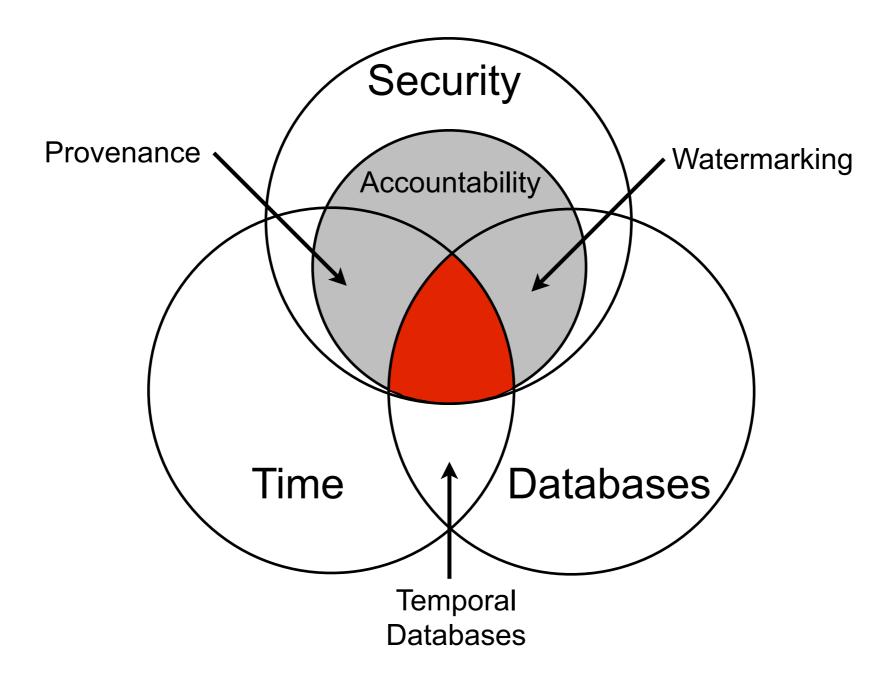




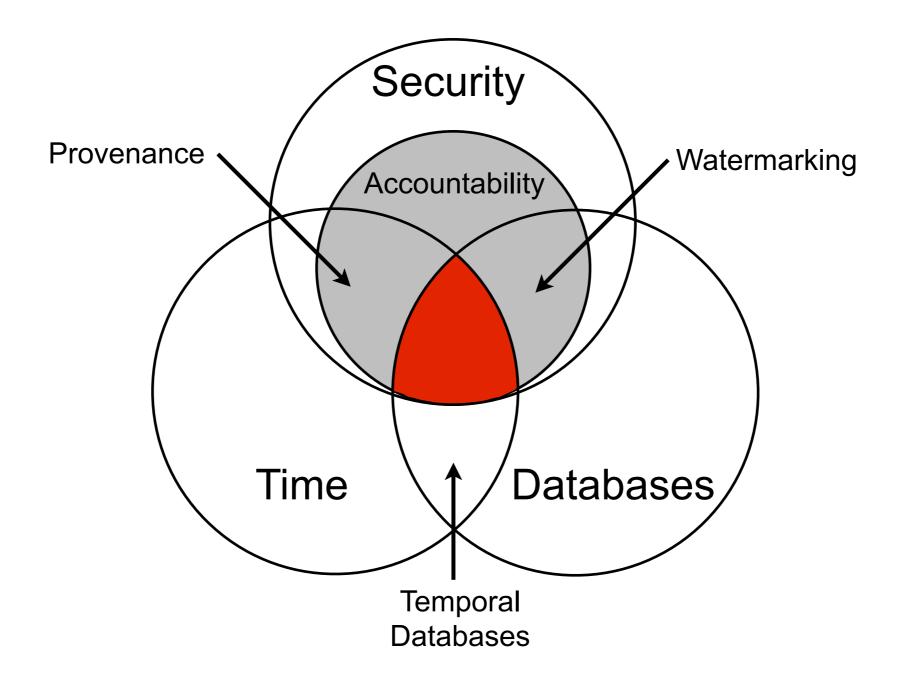








# Accountability ∩ Databases ∩ Time



Temporal concepts are found throughout this area of interest.

### Outline

- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations

### Outline

- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations

# Approach

#### Approach

- Continuous assurance technology
  - provides technology-enabled auditing
  - produces audit results close to occurrence of relevant events
  - achieves meaningful operationalization of information accountability.

#### Approach

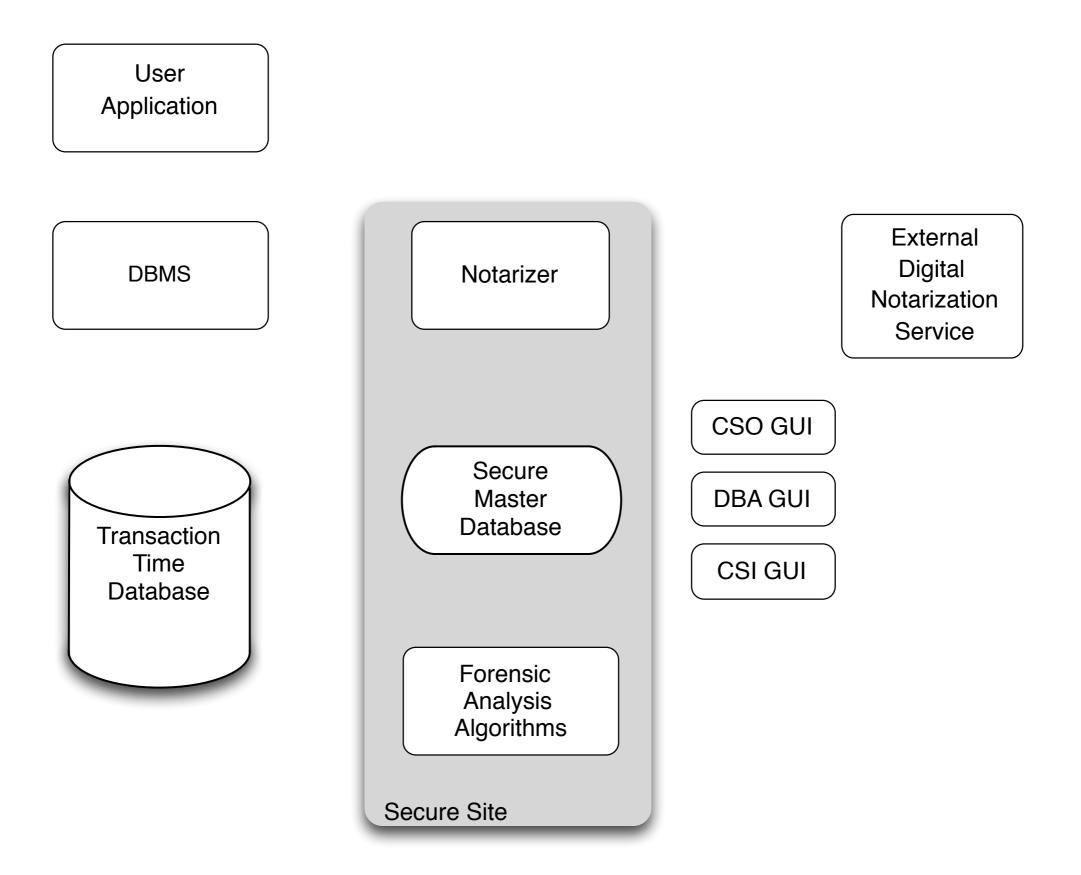
- Continuous assurance technology
  - provides technology-enabled auditing
  - produces audit results close to occurrence of relevant events
  - achieves meaningful operationalization of information accountability.
- Cryptographic hashing captures state of database as it evolves.

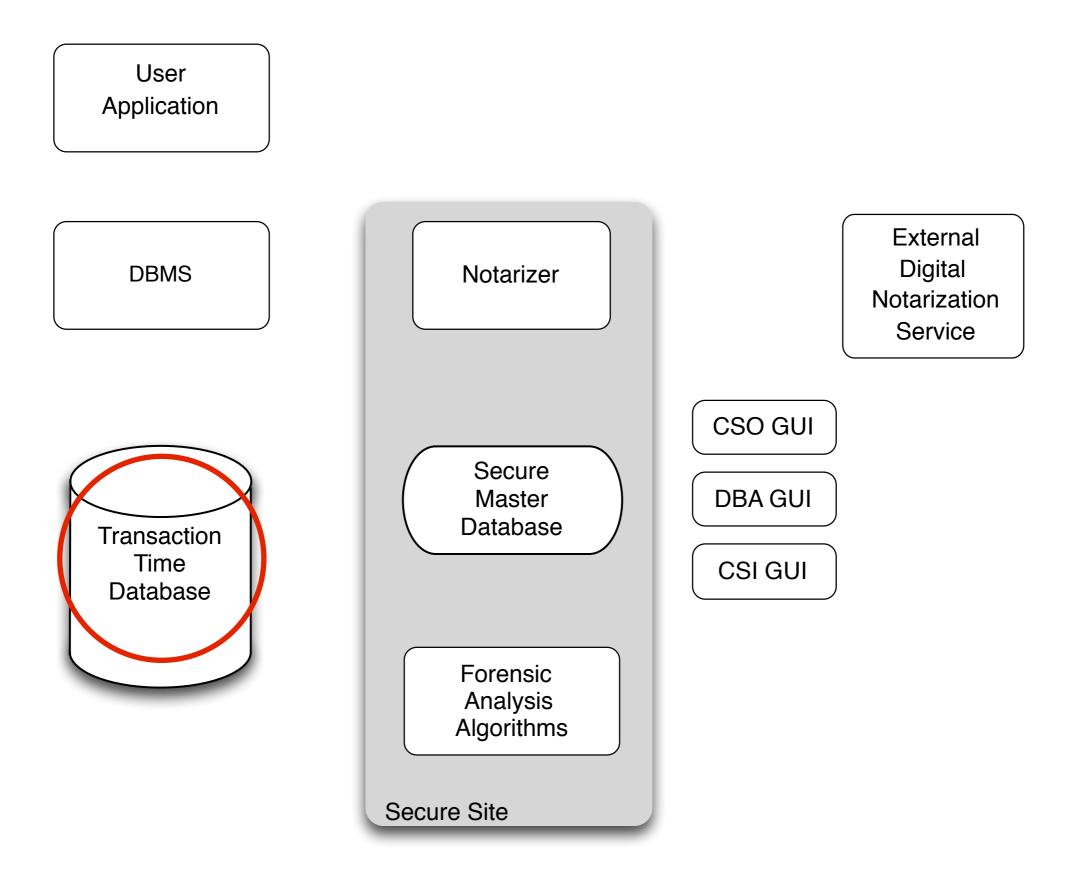
- Trusted computing base (TCB)
  - Correctly booted and running hardware, OS and DBMS
  - TCB runs correctly until intrusion
- A trusted external digital notarization service (EDNS)
- The adversary could be
  - Inside/outside intruders who gain full control of the whole TCB and logs
  - Malware such as virus, bugs, power surge
- Regret Interval: minimum time before someone can reverse the change
  - Determined by the specific application

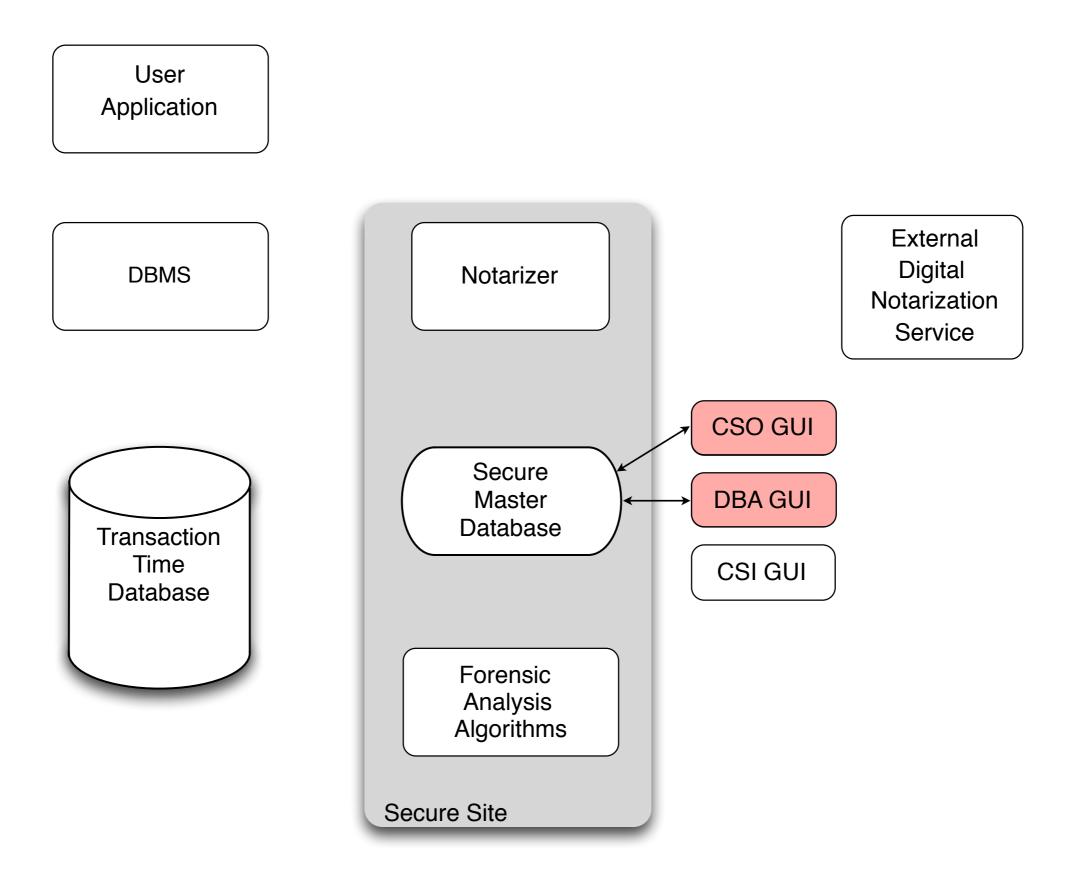
- Trusted computing base (TCB)
  - Correctly booted and running hardware, OS and DBMS
  - TCB runs correctly until intrusion
- A trusted external digital notarization service (EDNS)
- The adversary could be
  - Inside/outside intruders who gain full control of the whole TCB and logs
  - Malware such as virus, bugs, power surge
- Regret Interval: minimum time before someone can reverse the change
  - Determined by the specific application

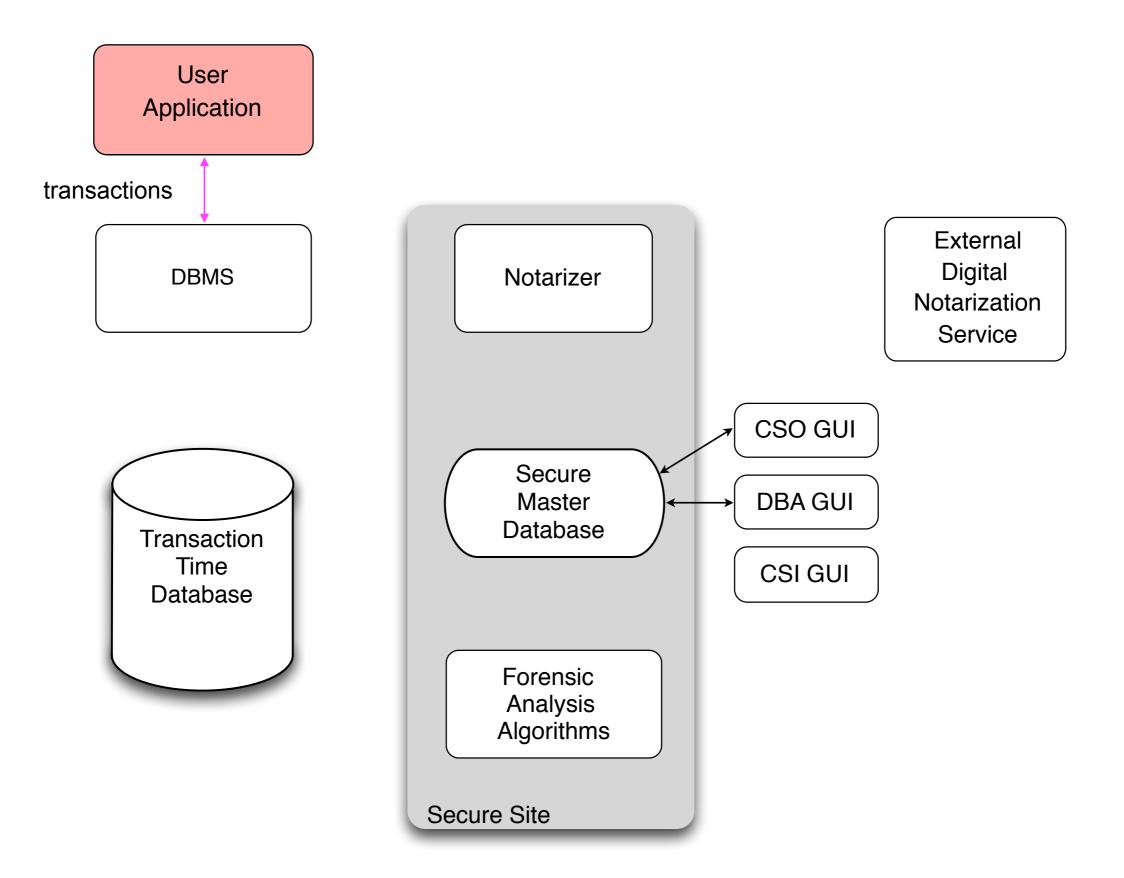
- Trusted computing base (TCB)
  - Correctly booted and running hardware, OS and DBMS
  - TCB runs correctly until intrusion
- A trusted external digital notarization service (EDNS)
- The adversary could be
  - Inside/outside intruders who gain full control of the whole TCB and logs
  - Malware such as virus, bugs, power surge
- Regret Interval: minimum time before someone can reverse the change
  - Determined by the specific application

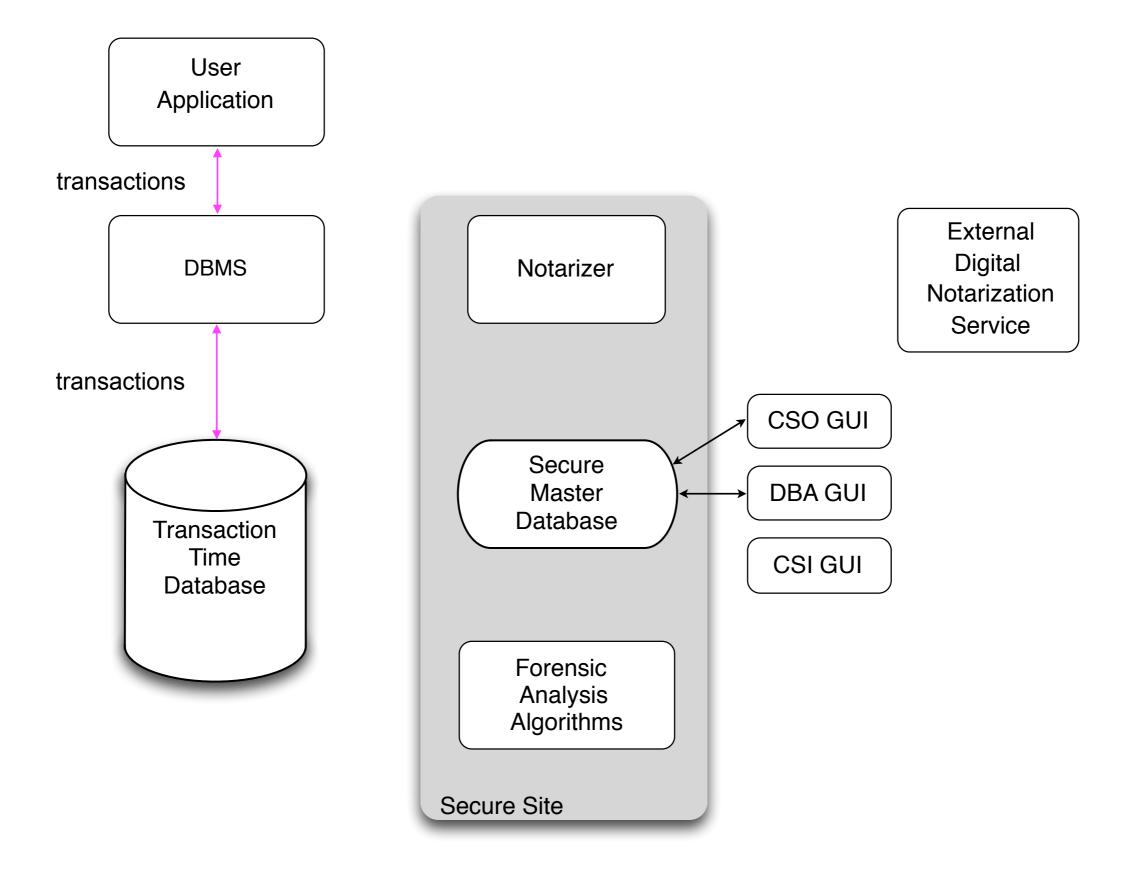
- Trusted computing base (TCB)
  - Correctly booted and running hardware, OS and DBMS
  - TCB runs correctly until intrusion
- A trusted external digital notarization service (EDNS)
- The adversary could be
  - Inside/outside intruders who gain full control of the whole TCB and logs
  - Malware such as virus, bugs, power surge
- Regret Interval: minimum time before someone can reverse the change
  - Determined by the specific application

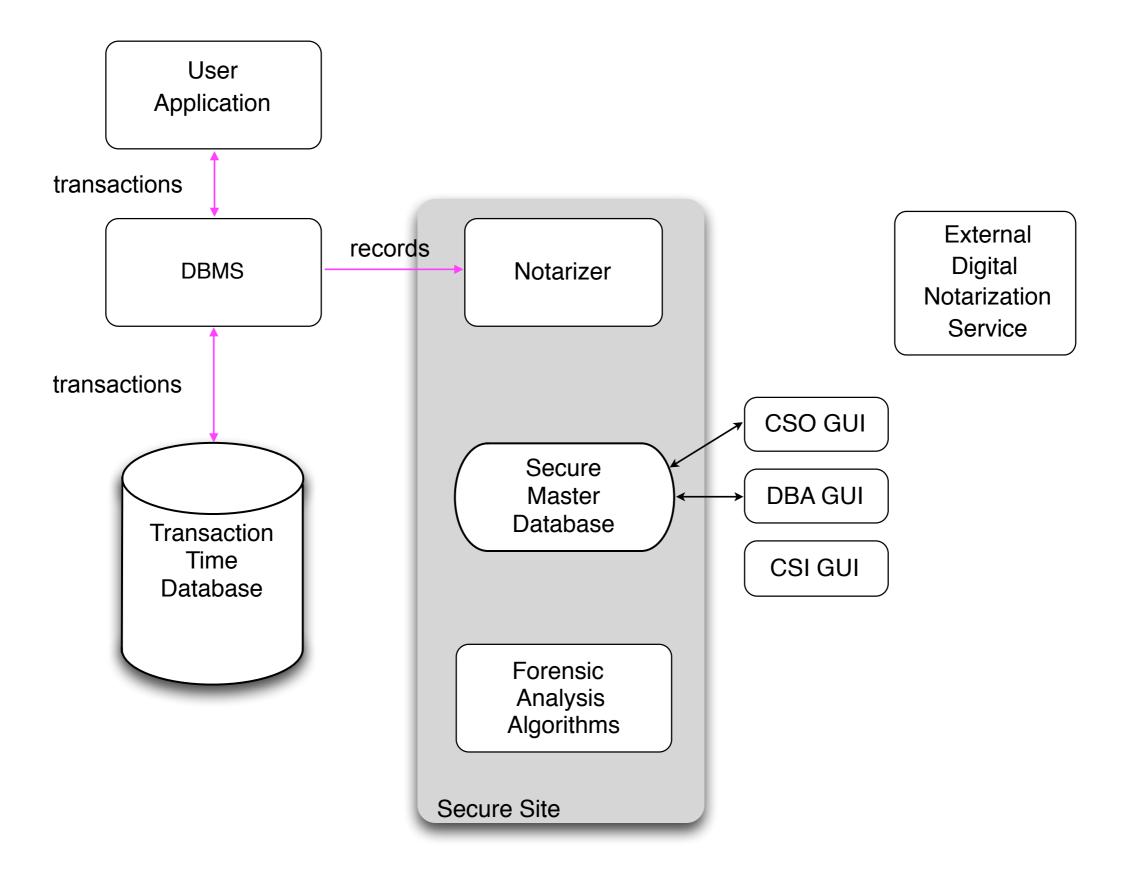




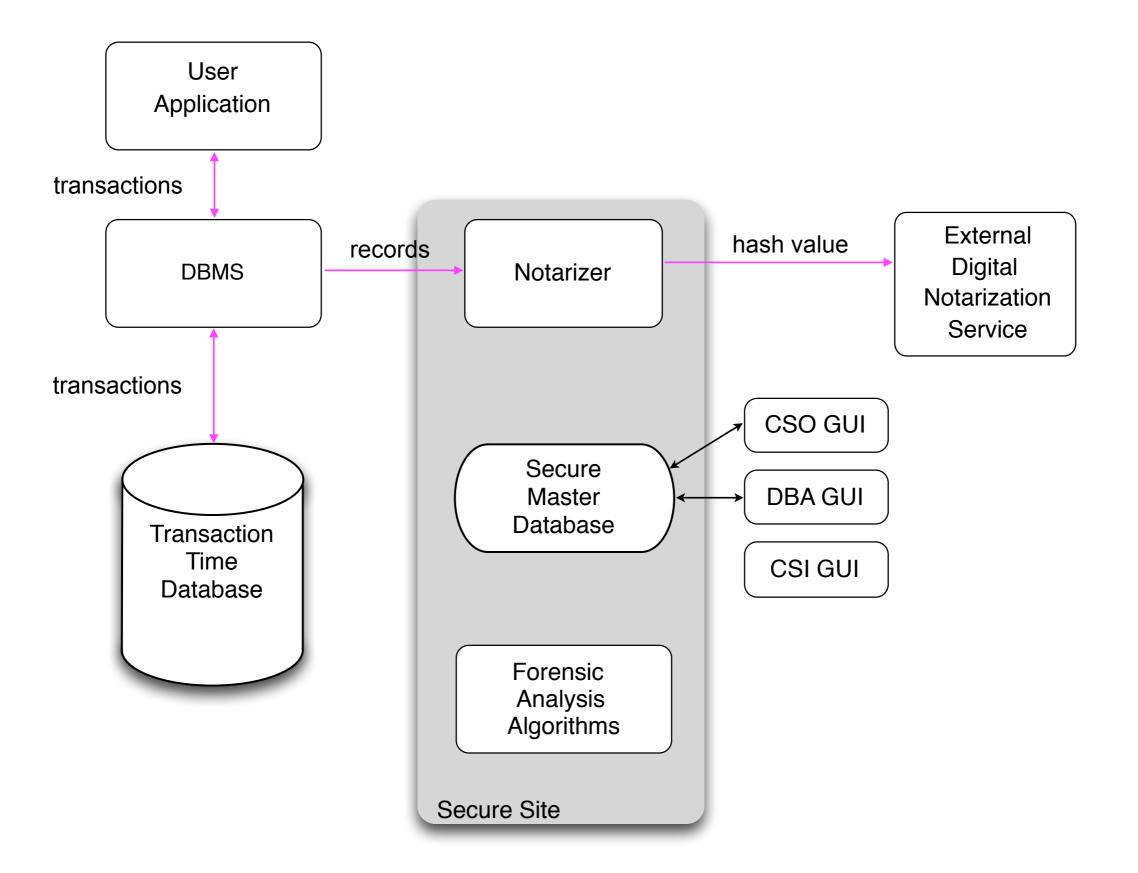




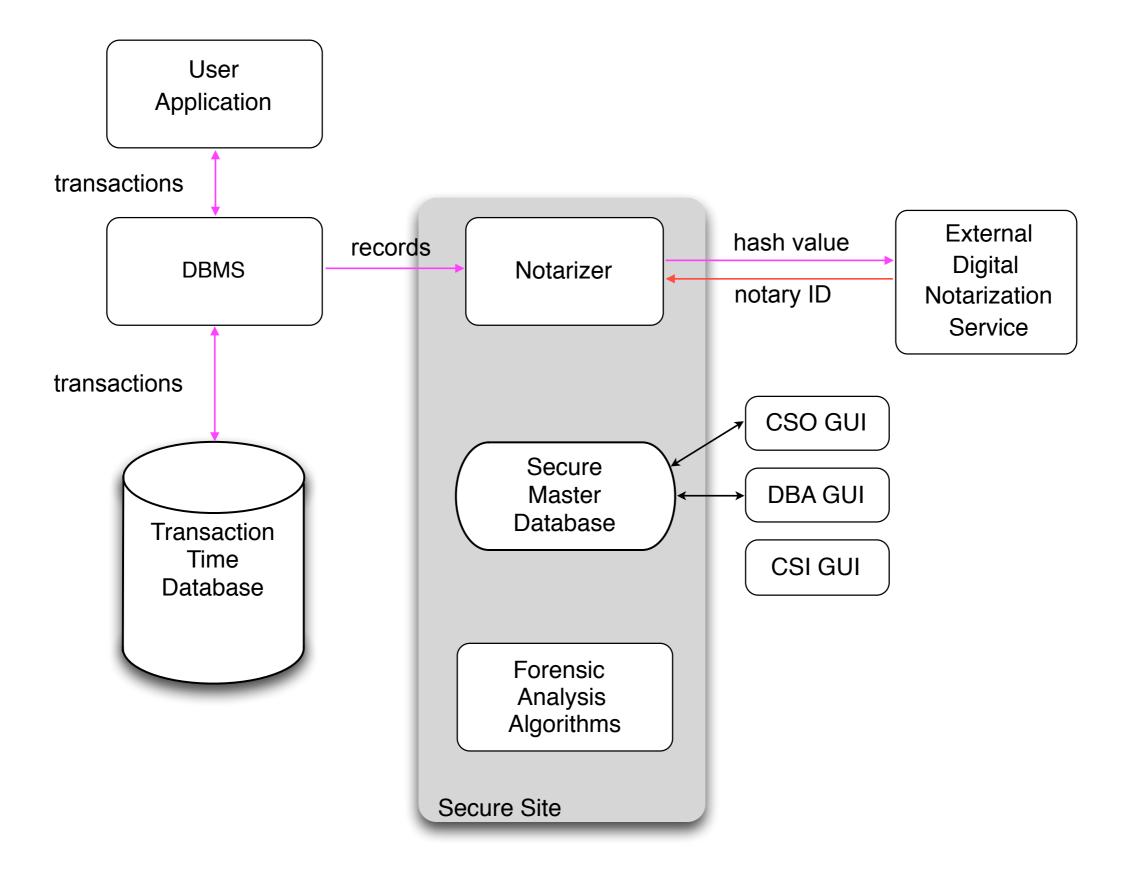




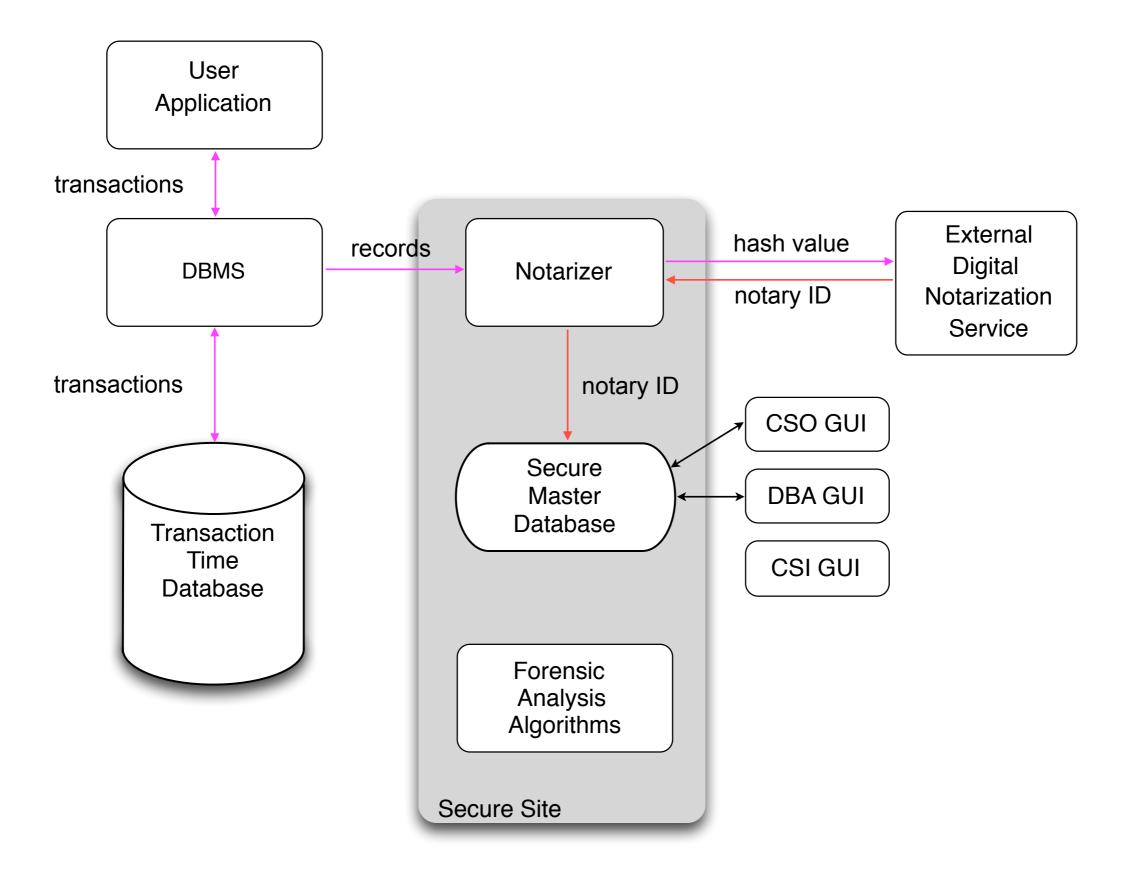
# **Total Chain Computation Phase**

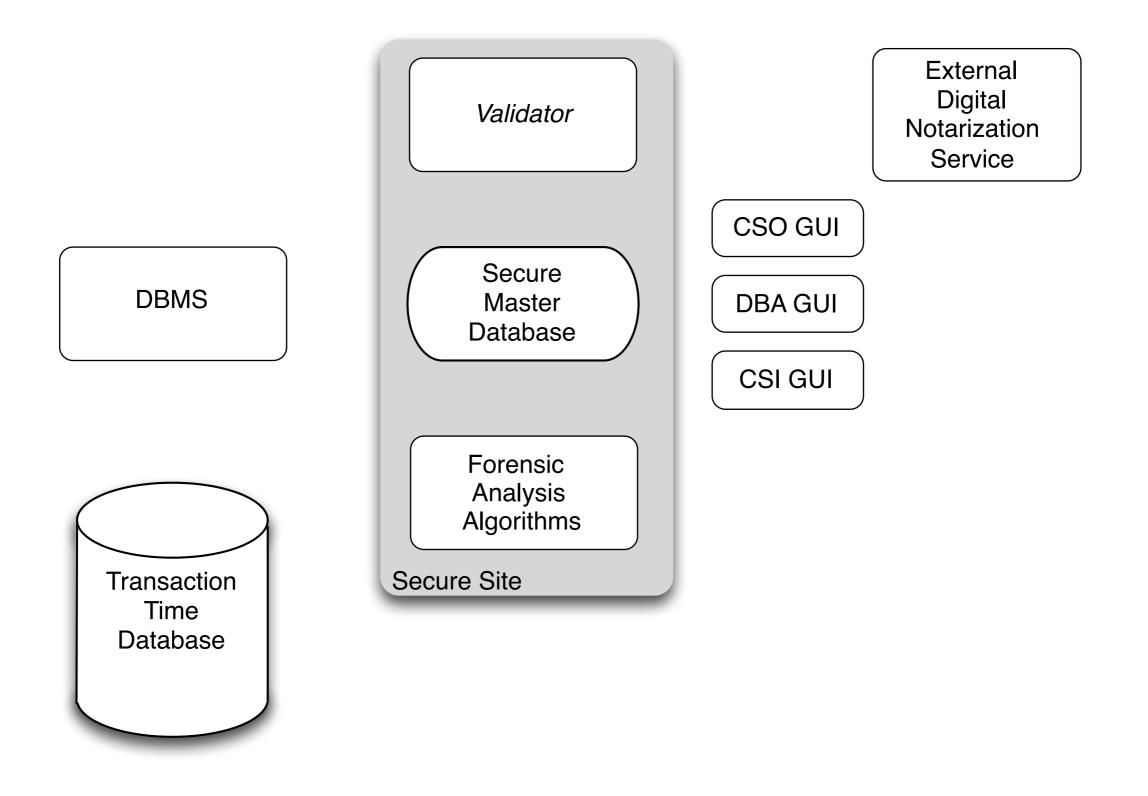


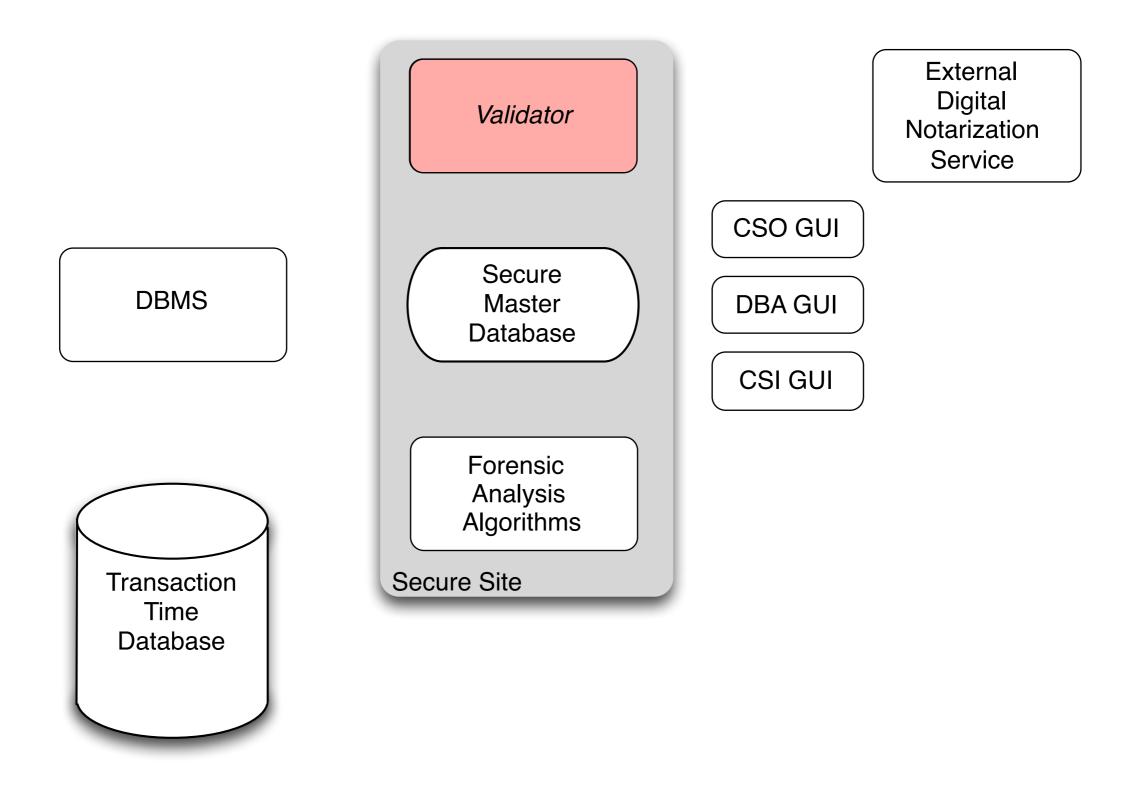
# **Total Chain Computation Phase**

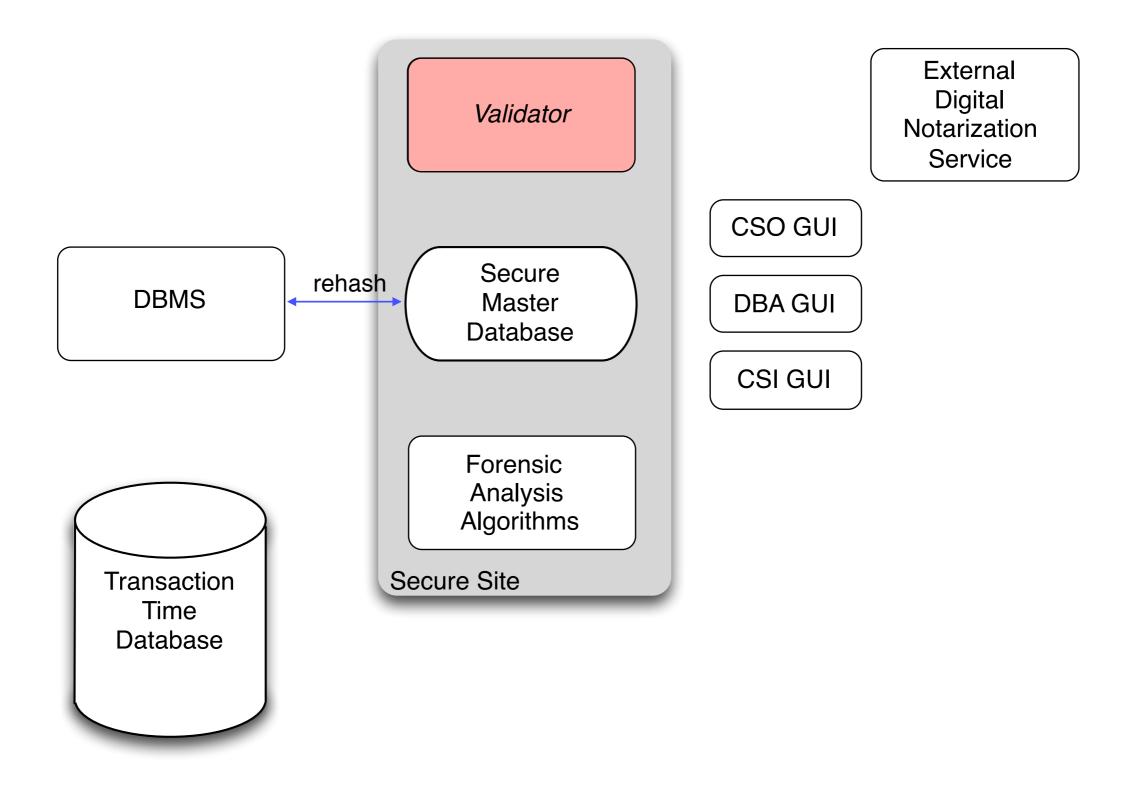


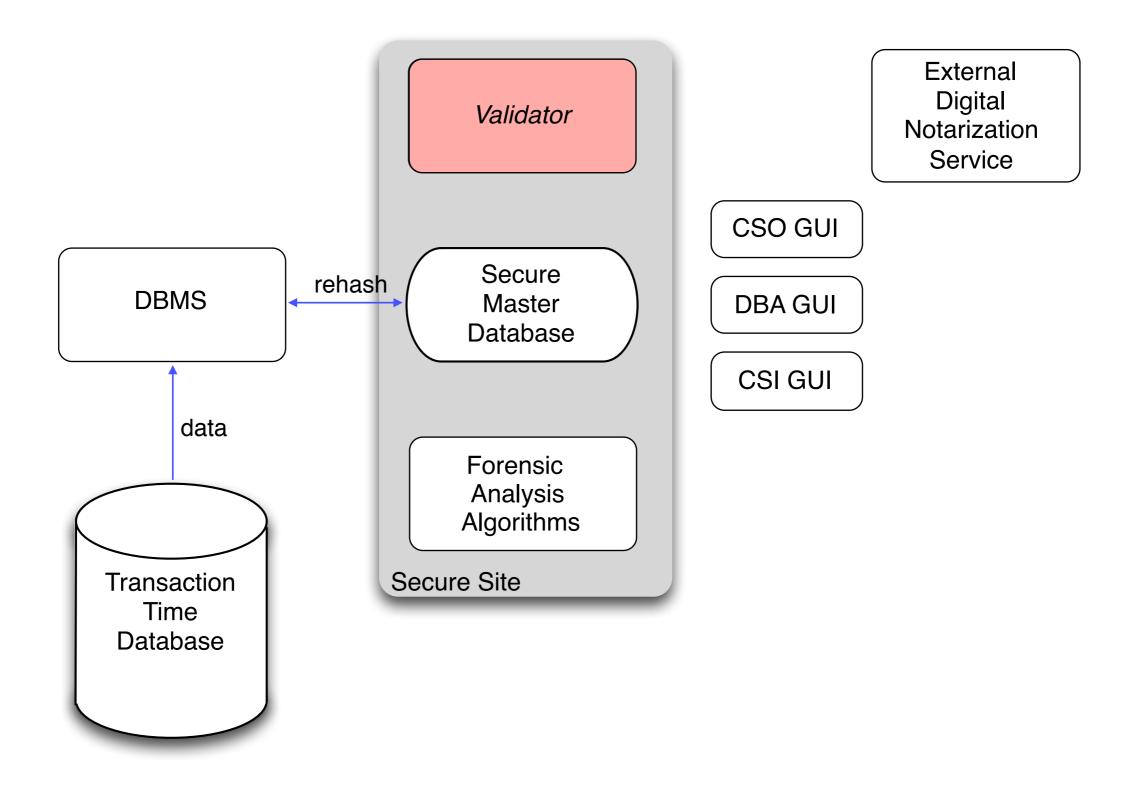
# **Total Chain Computation Phase**

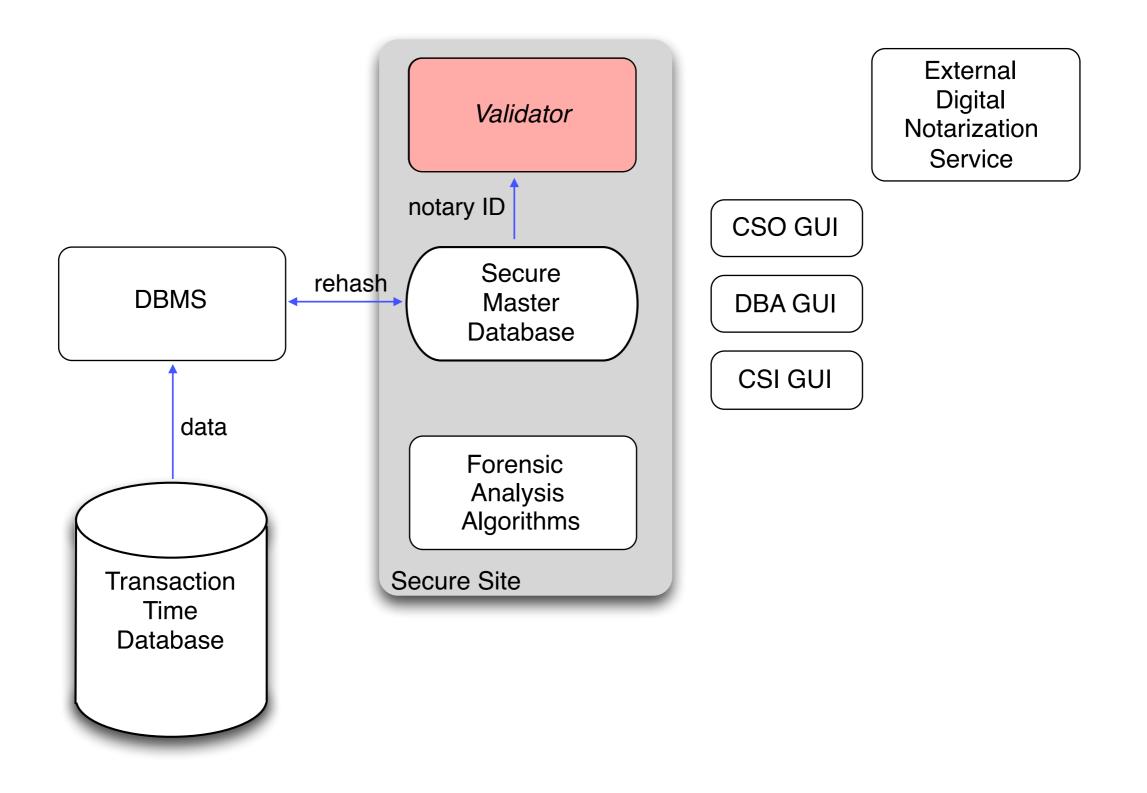


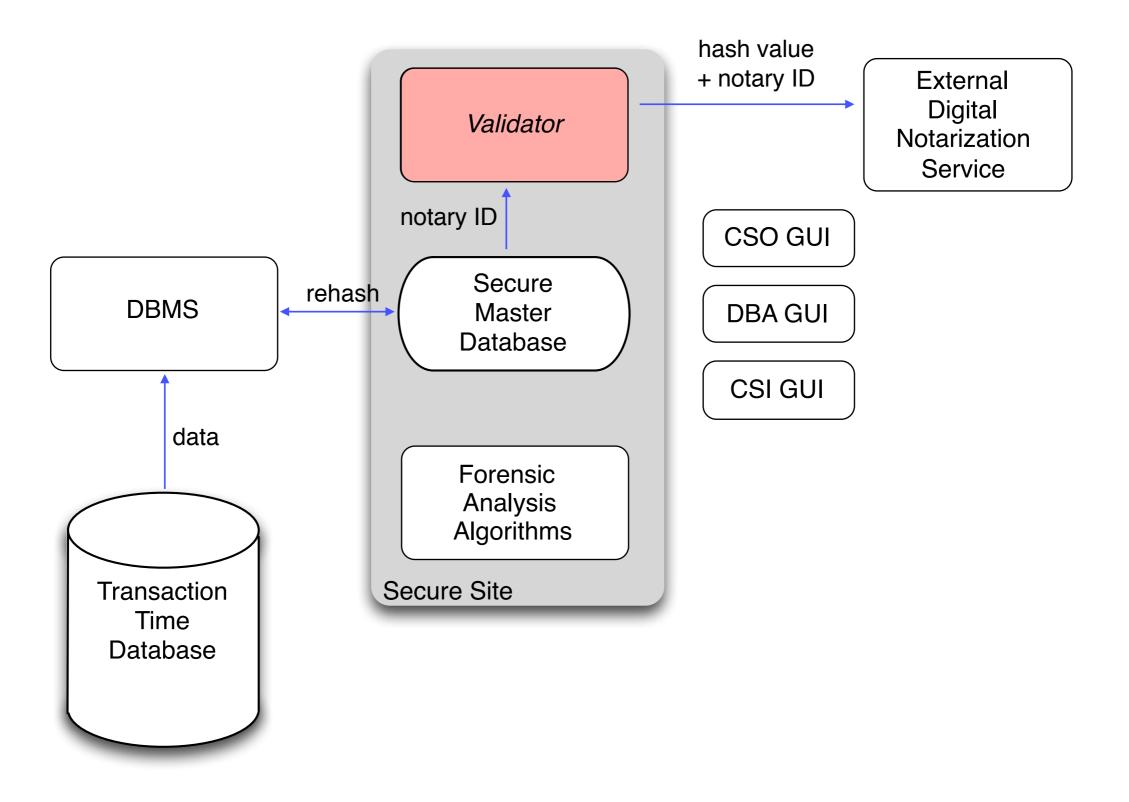


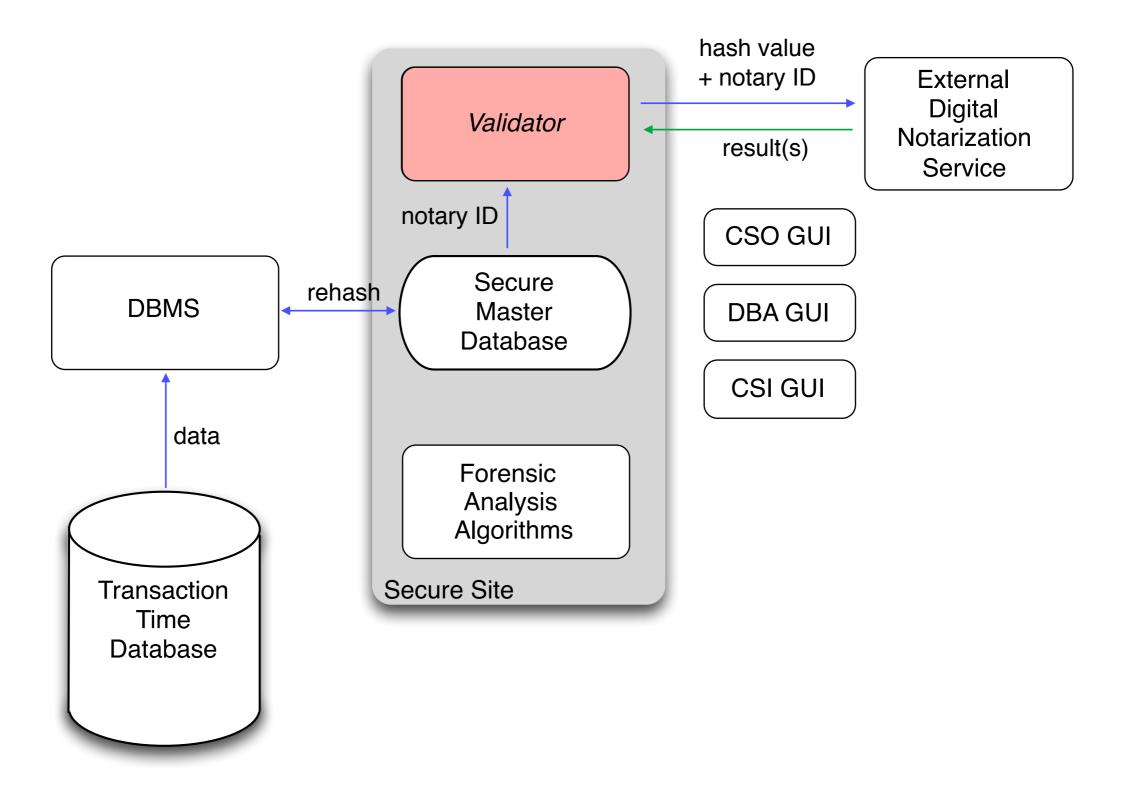


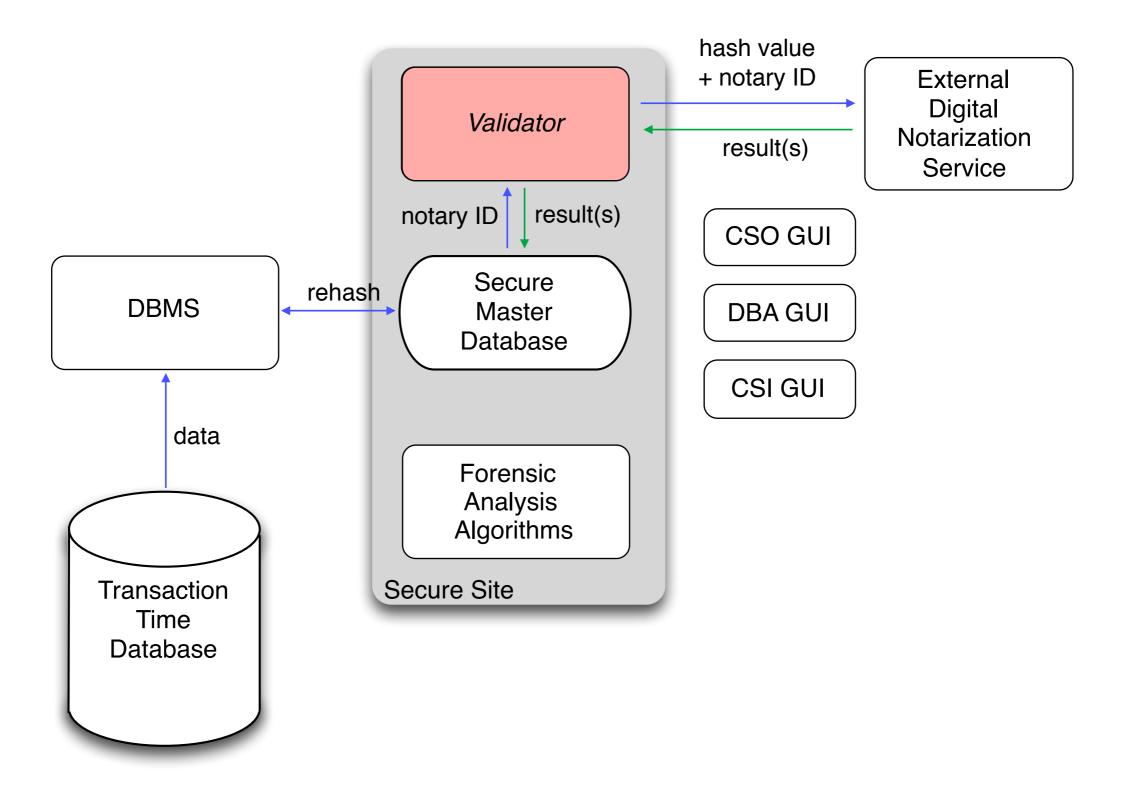


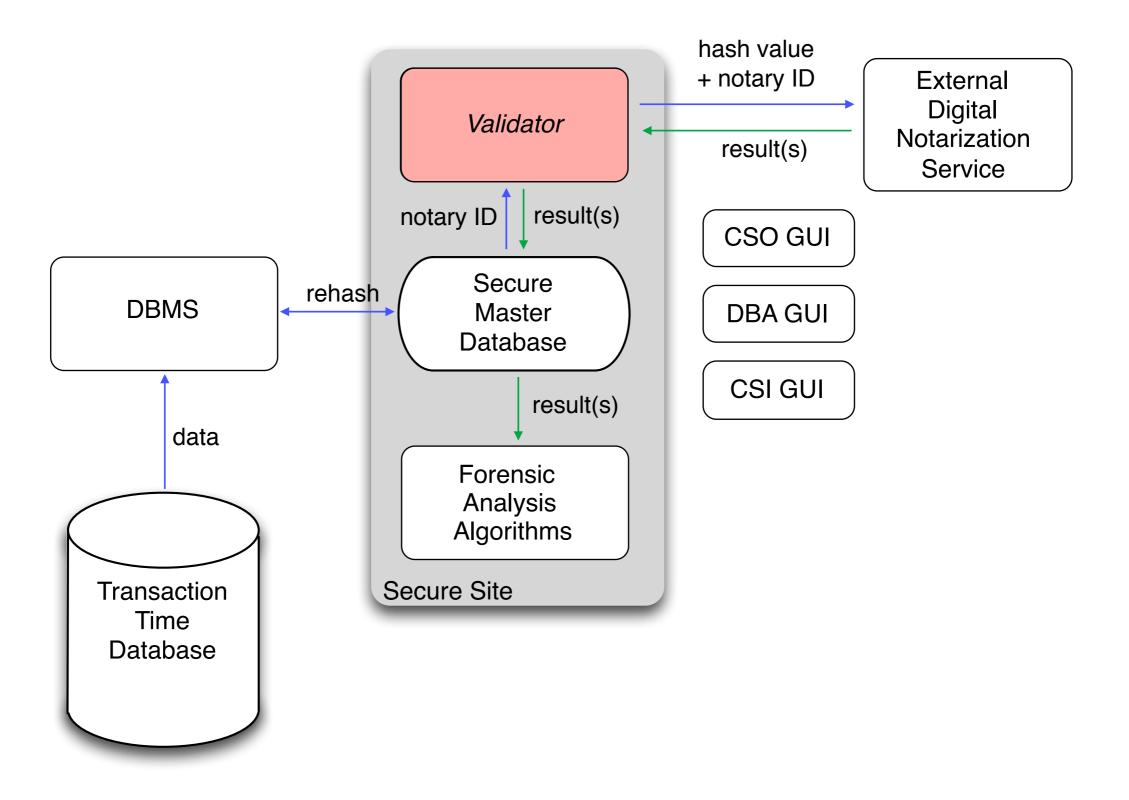


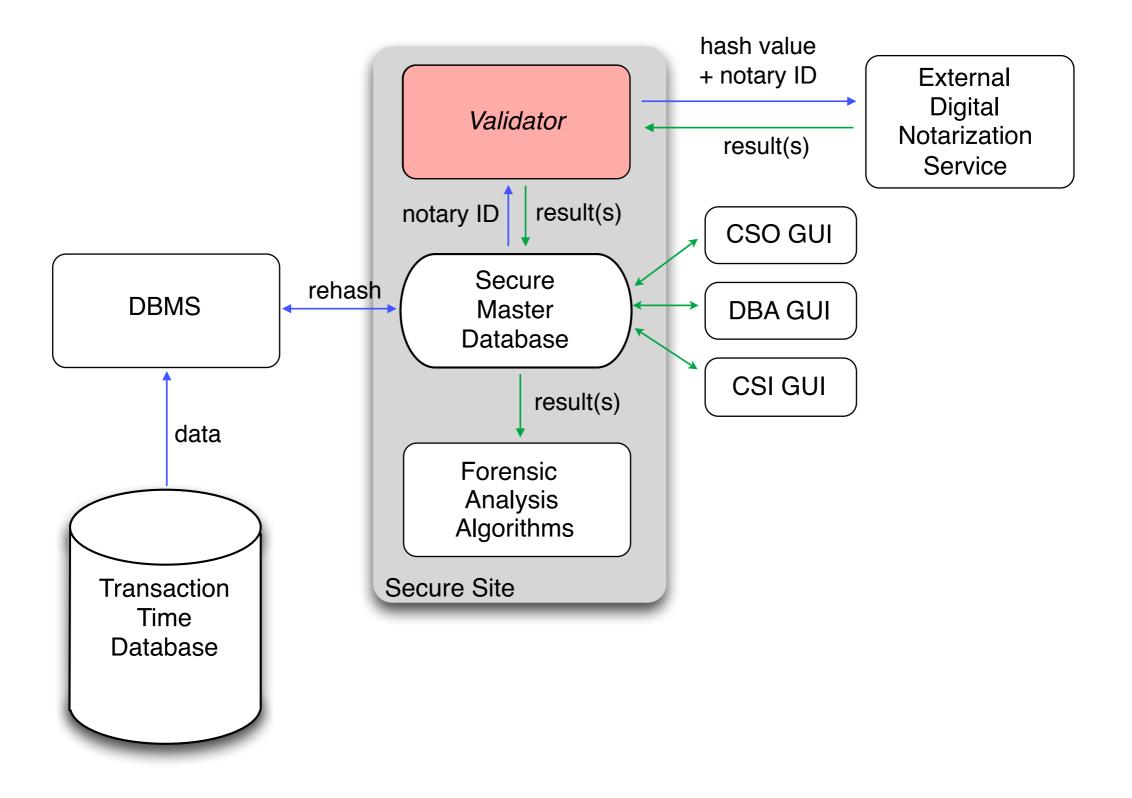


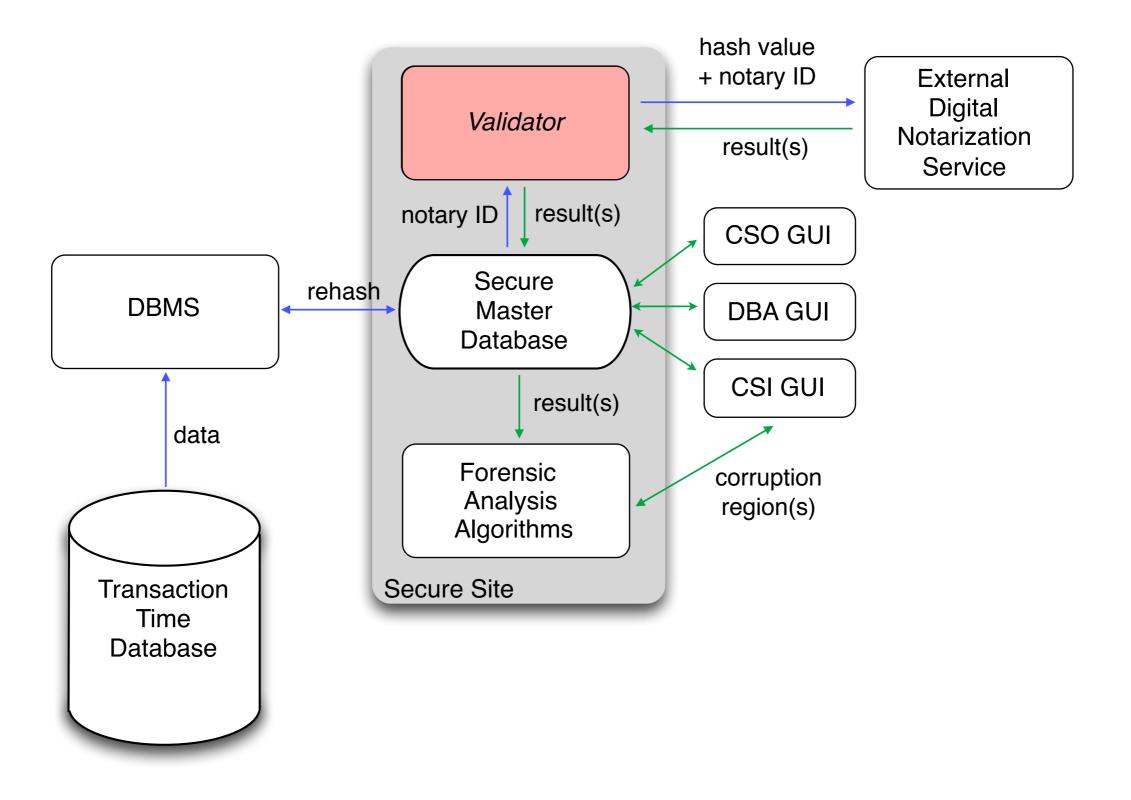






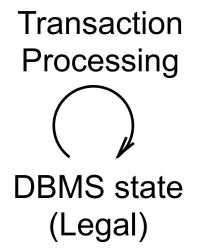


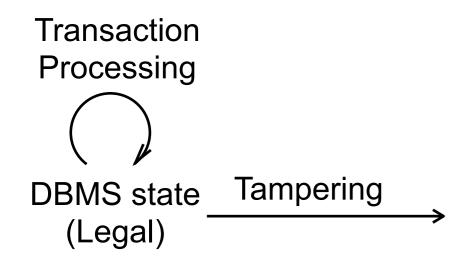


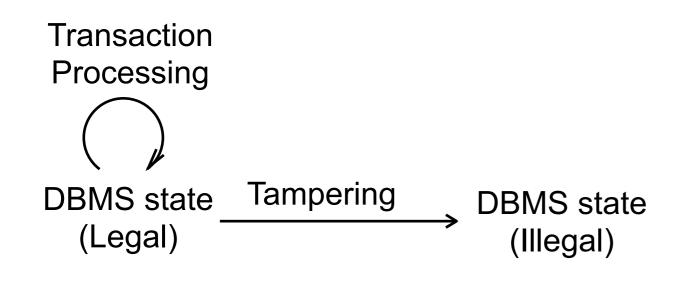


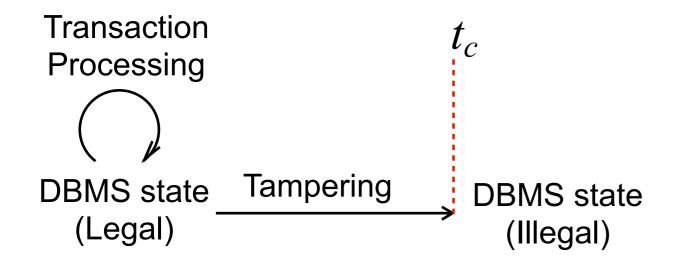
# Outline

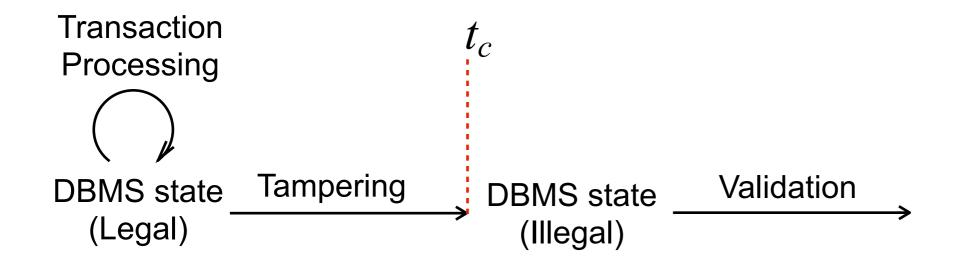
- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations

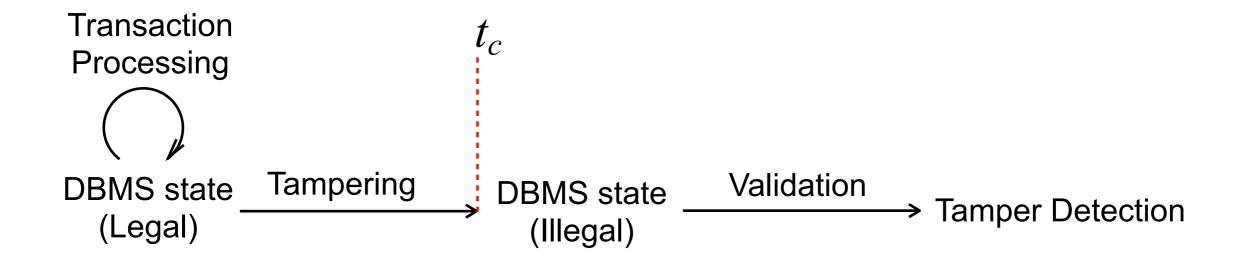


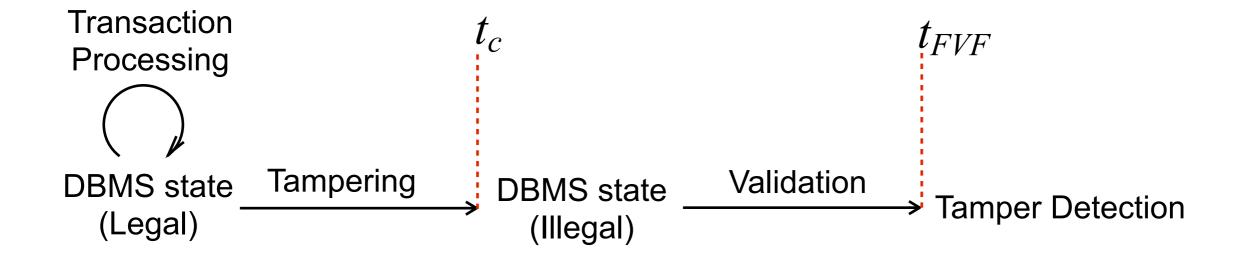


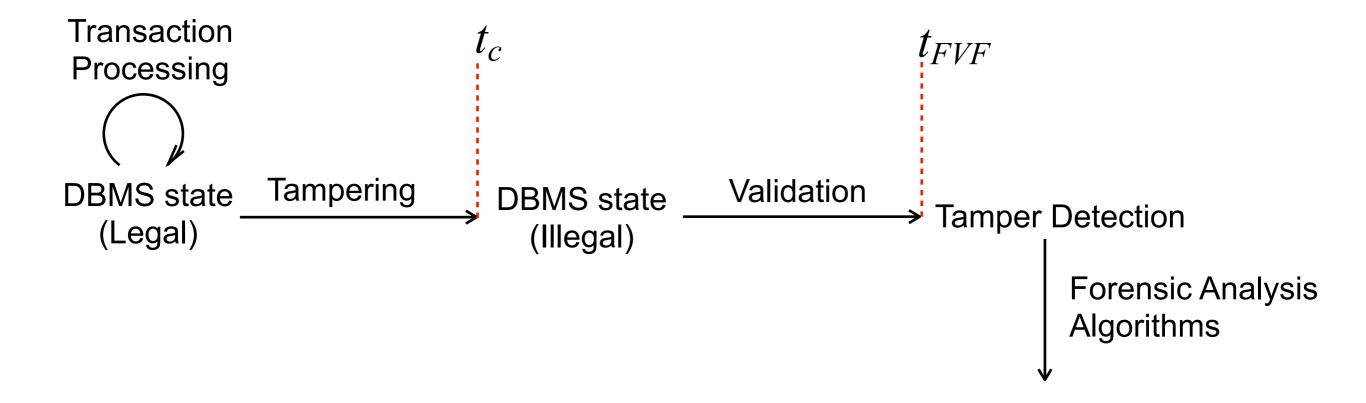


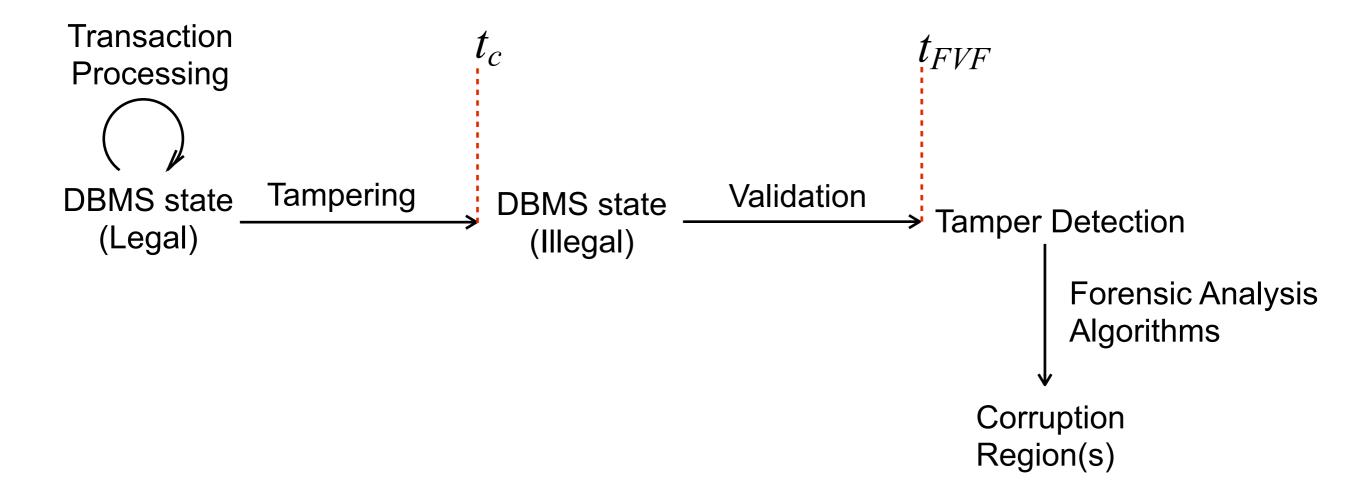


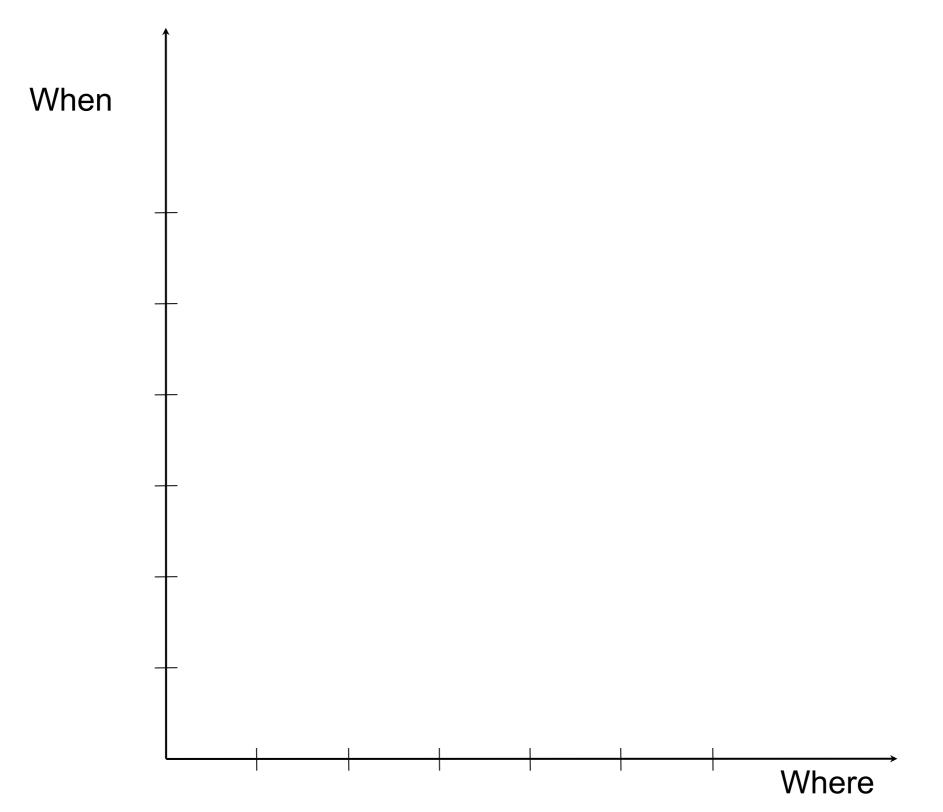


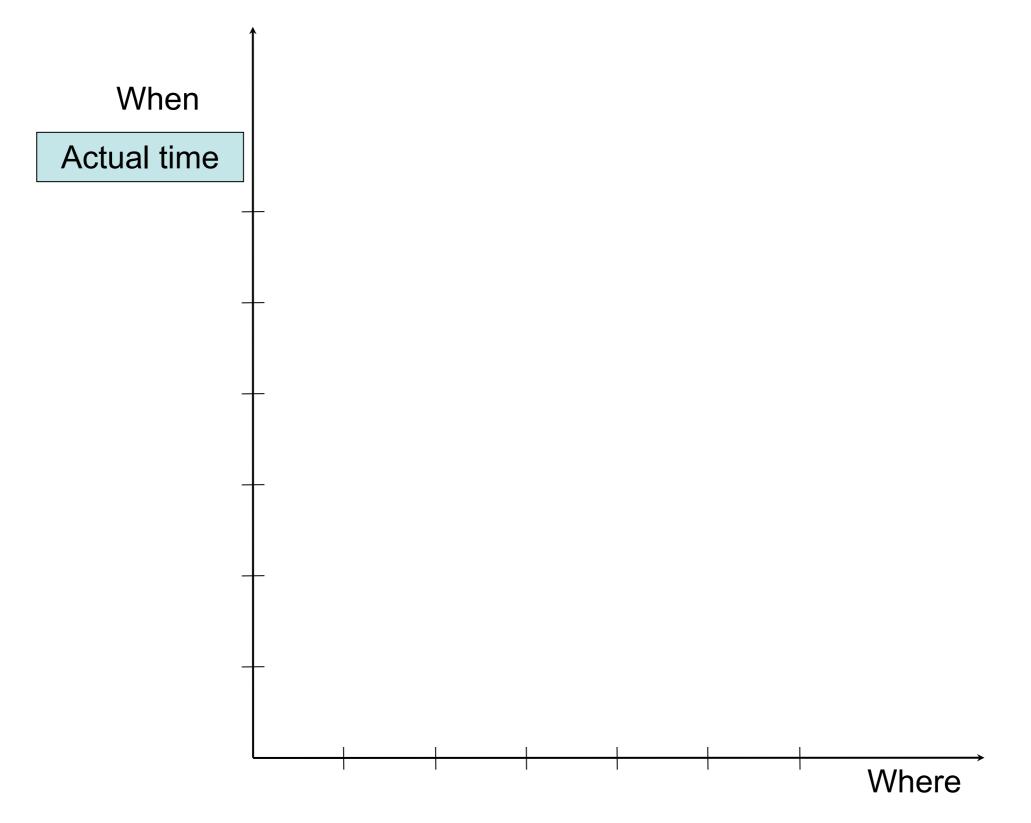


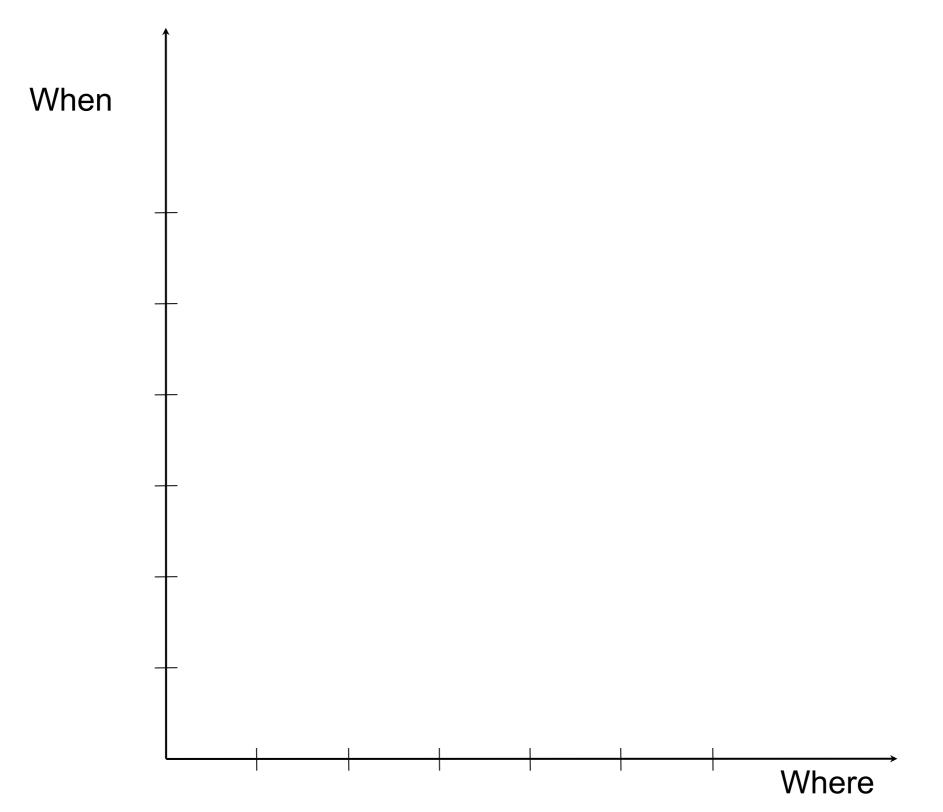


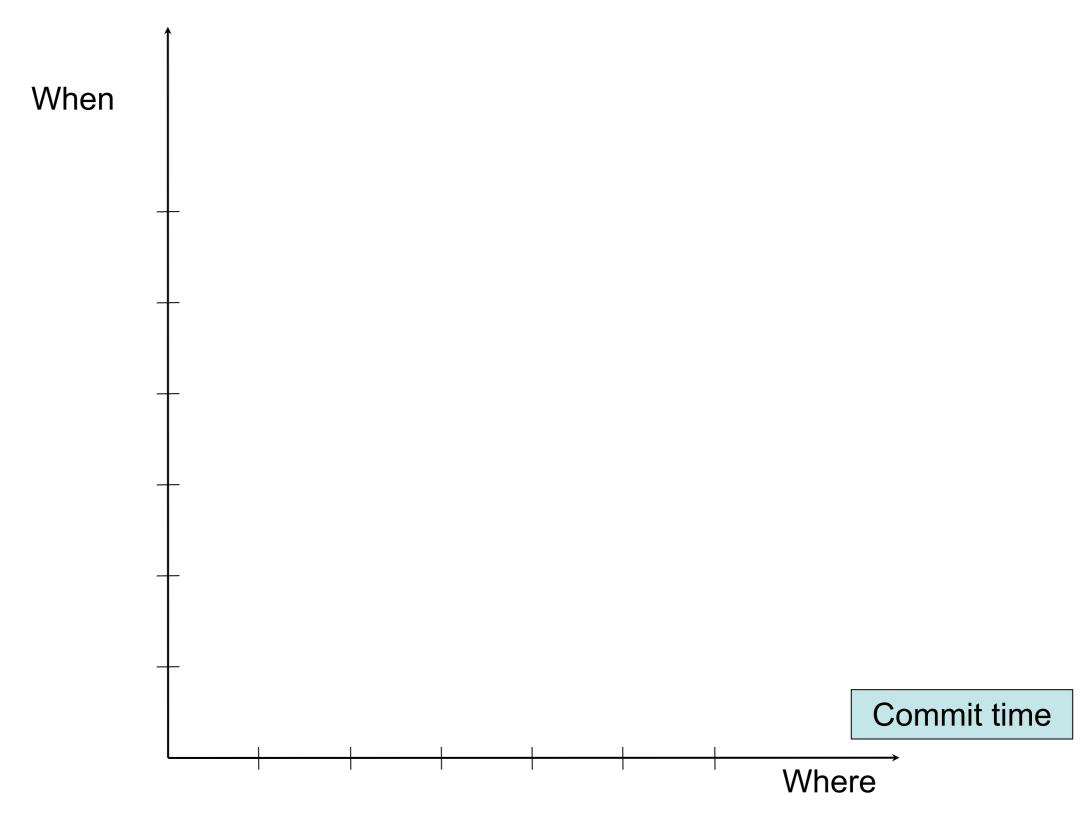


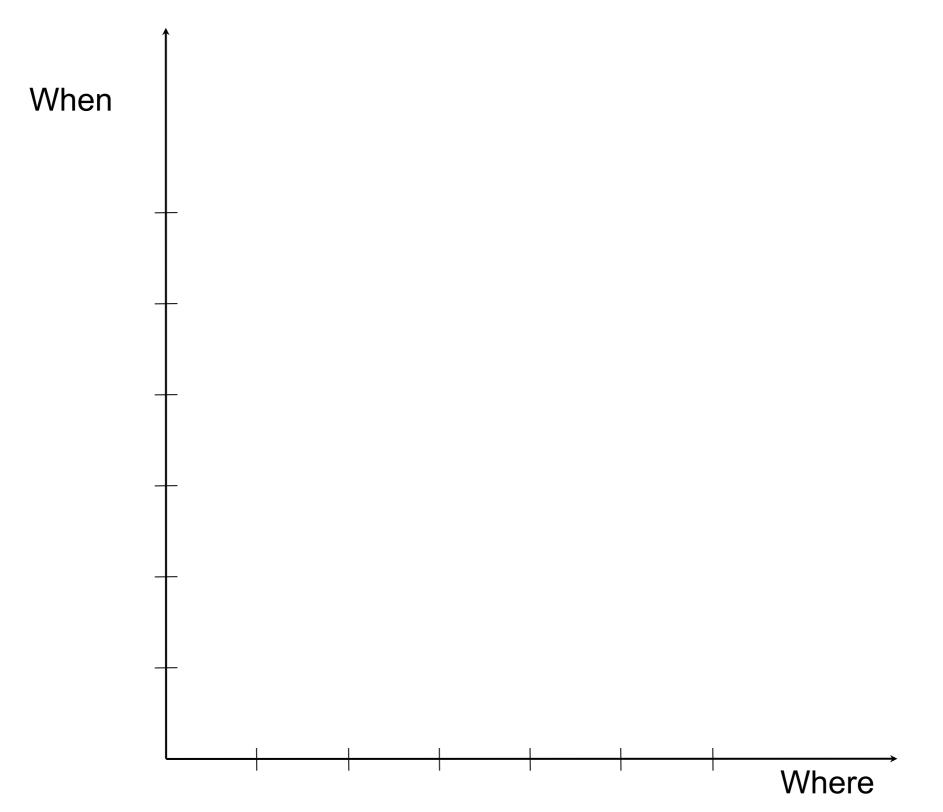


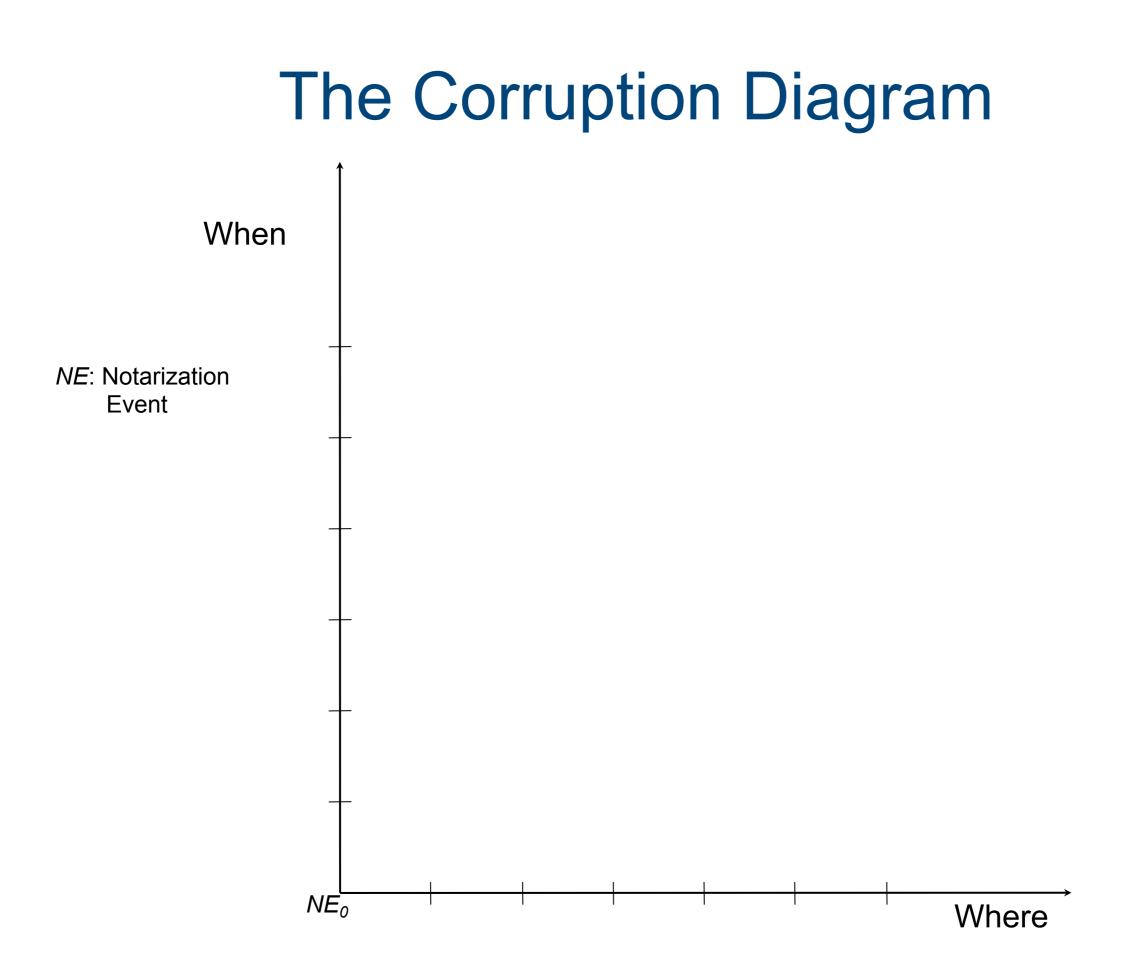


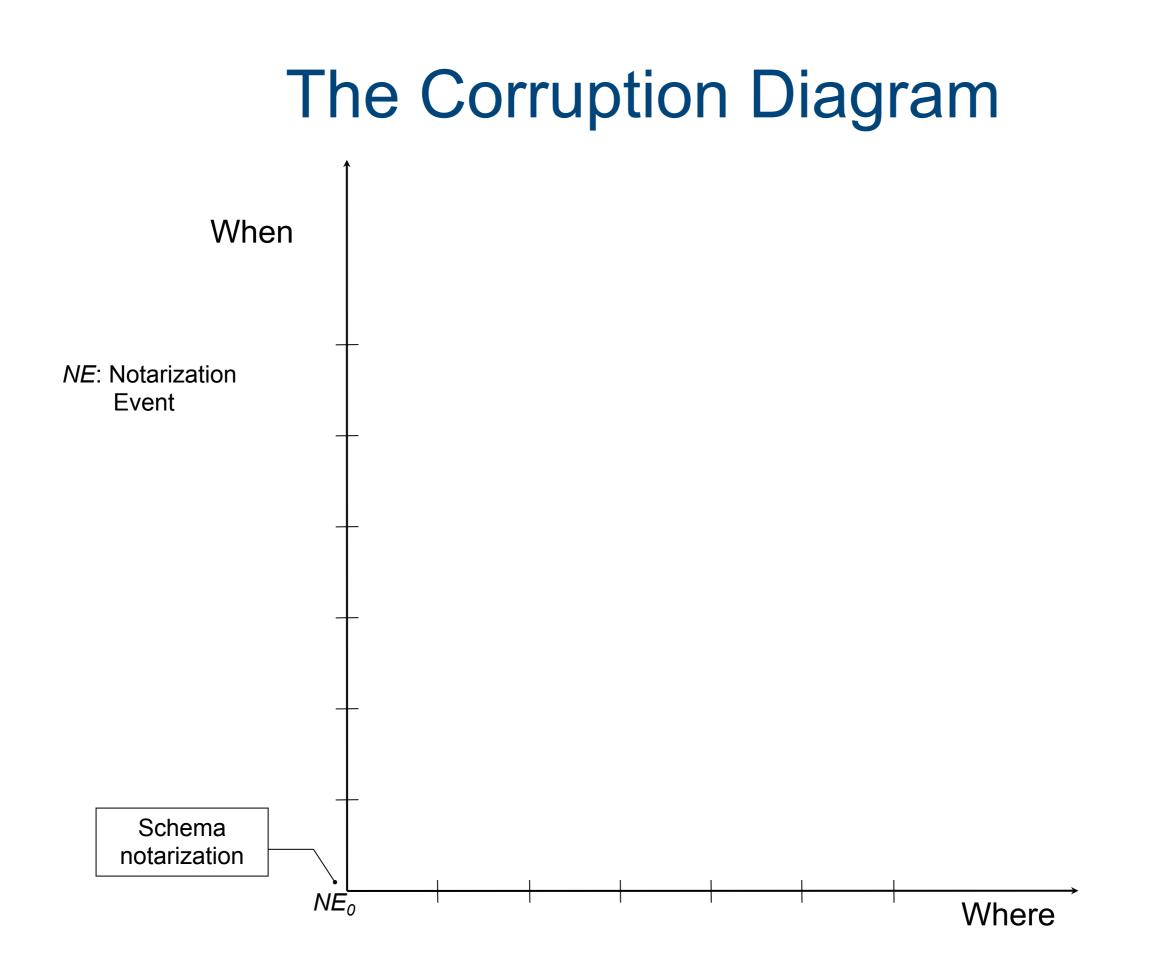


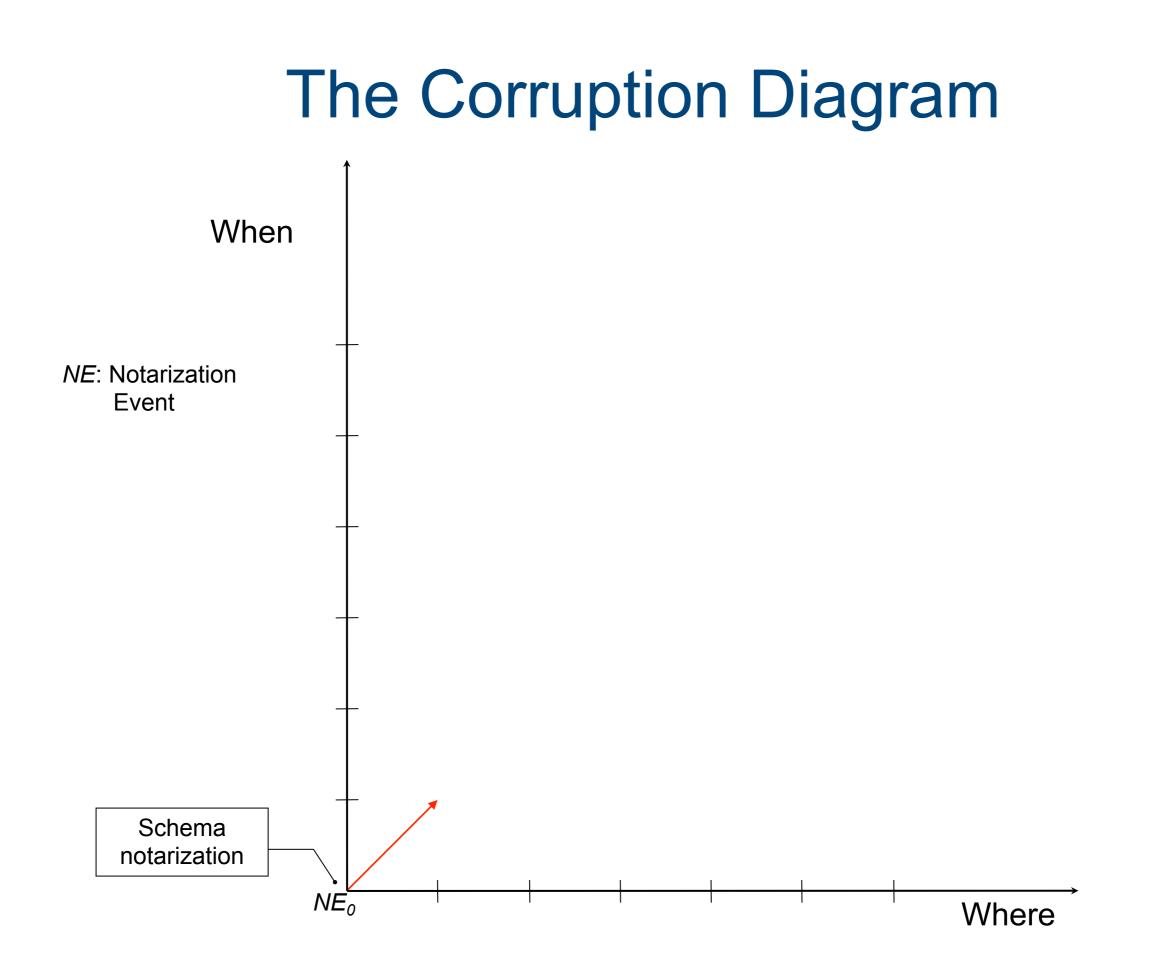


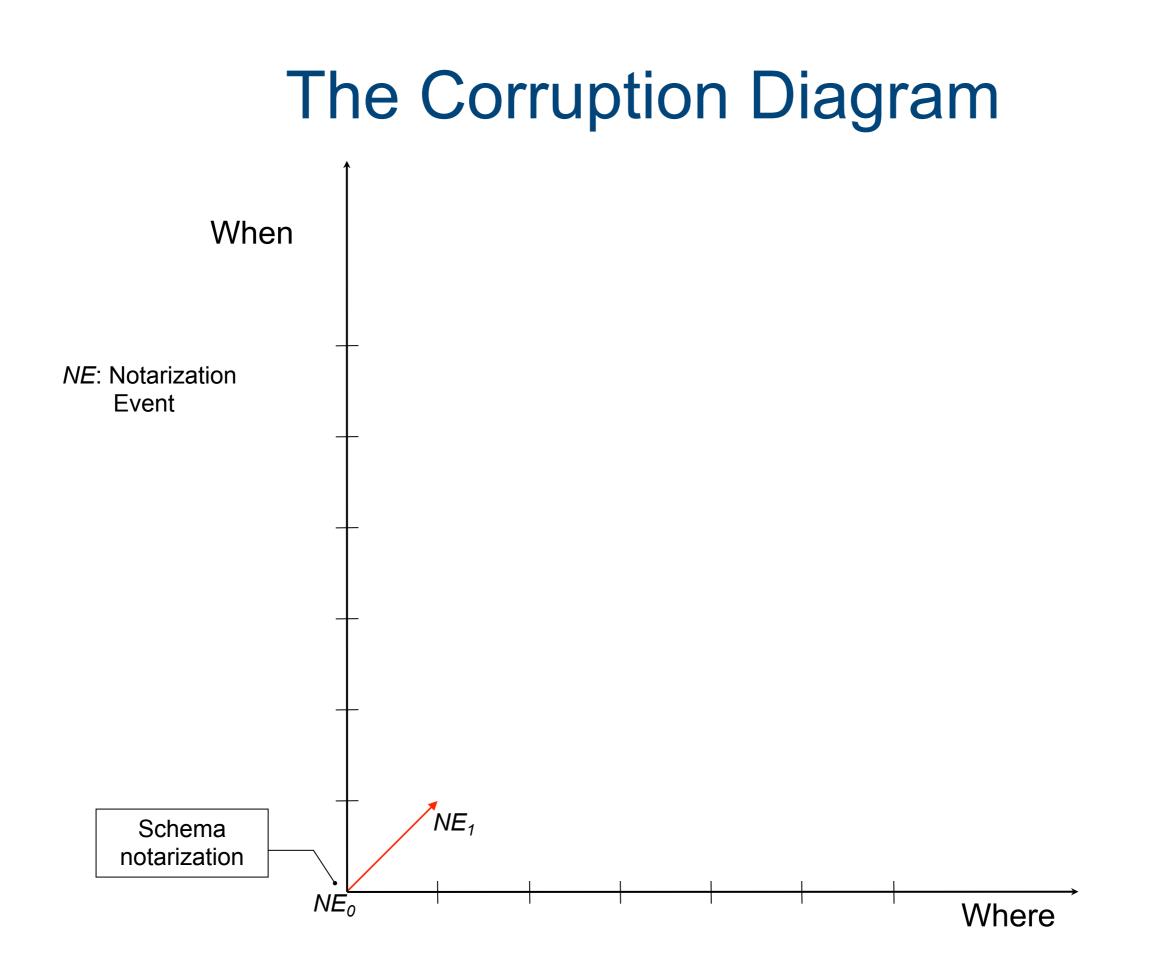


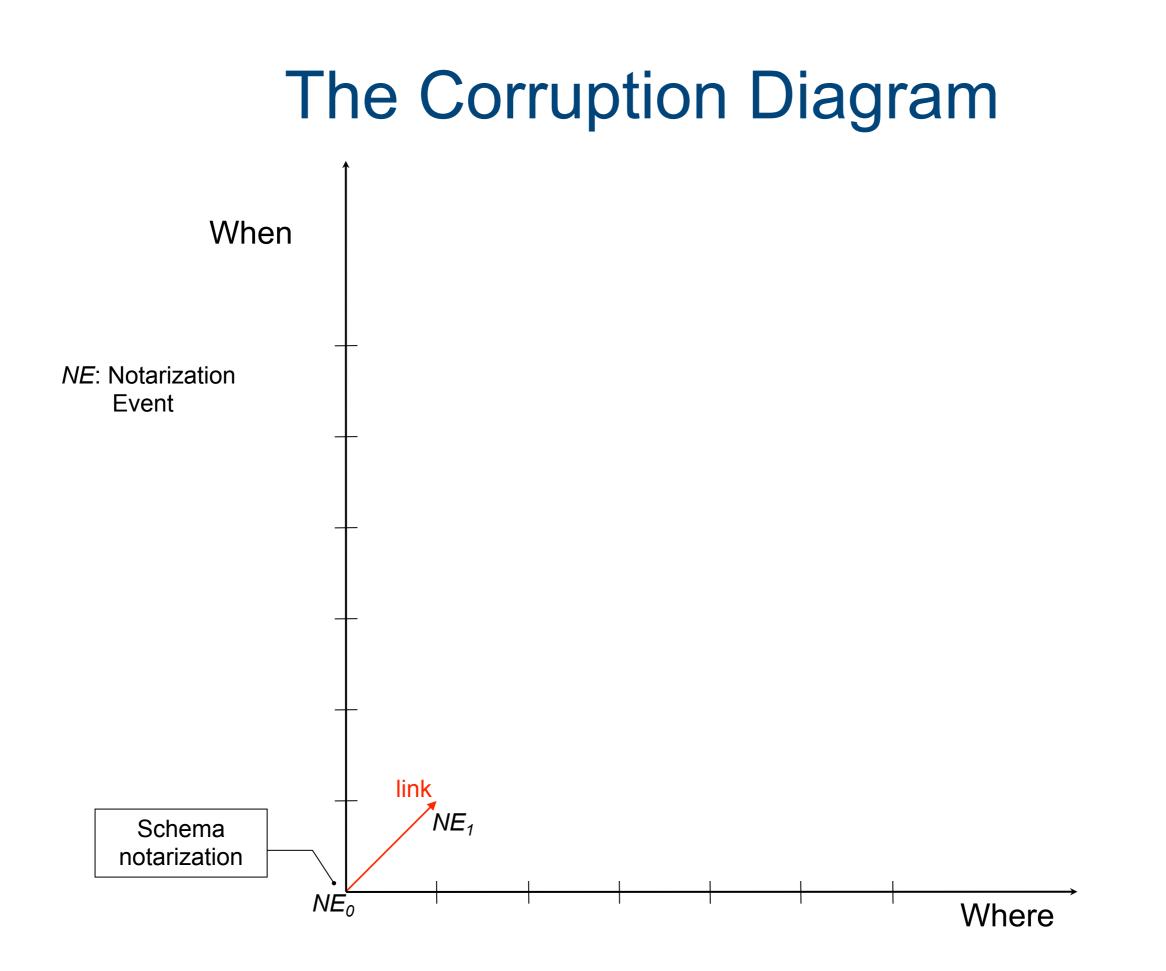


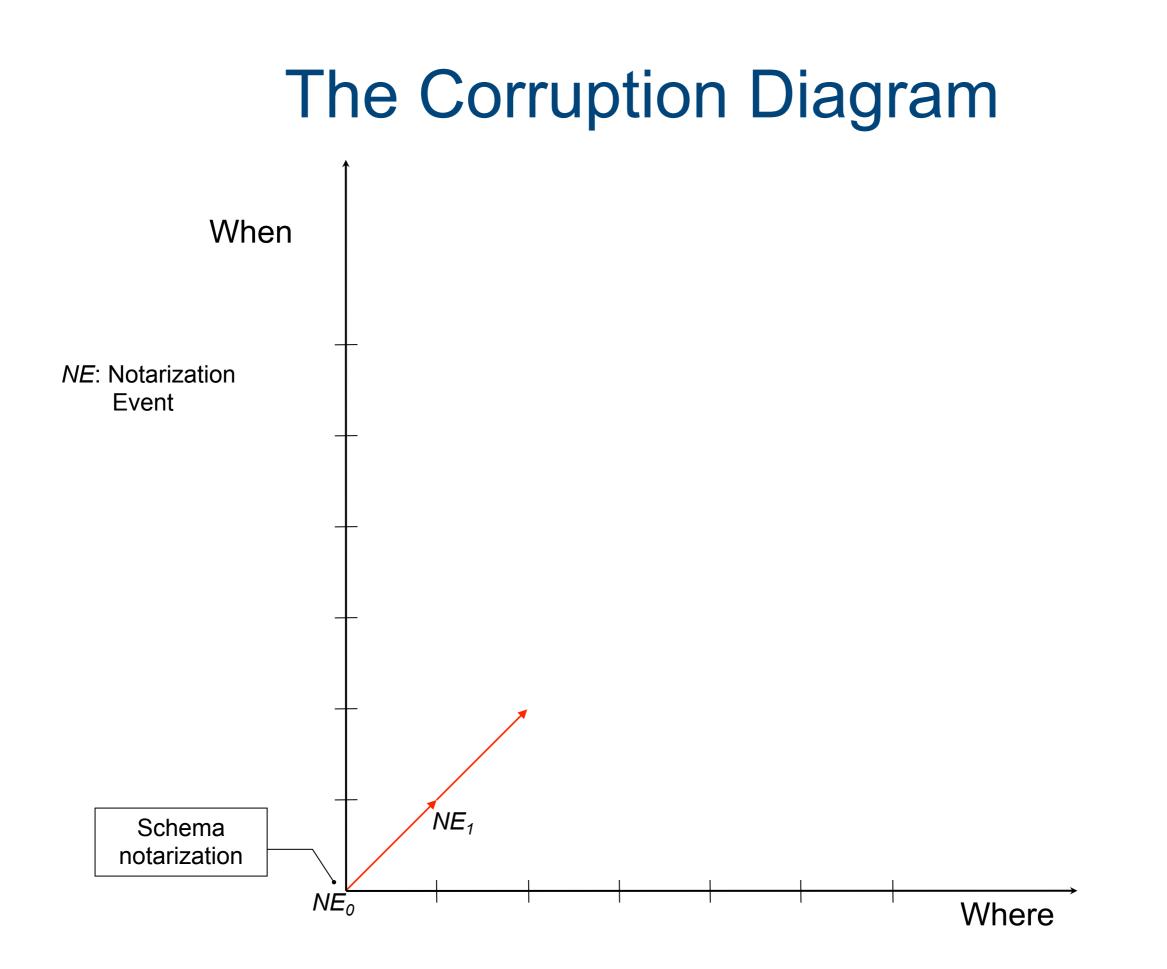


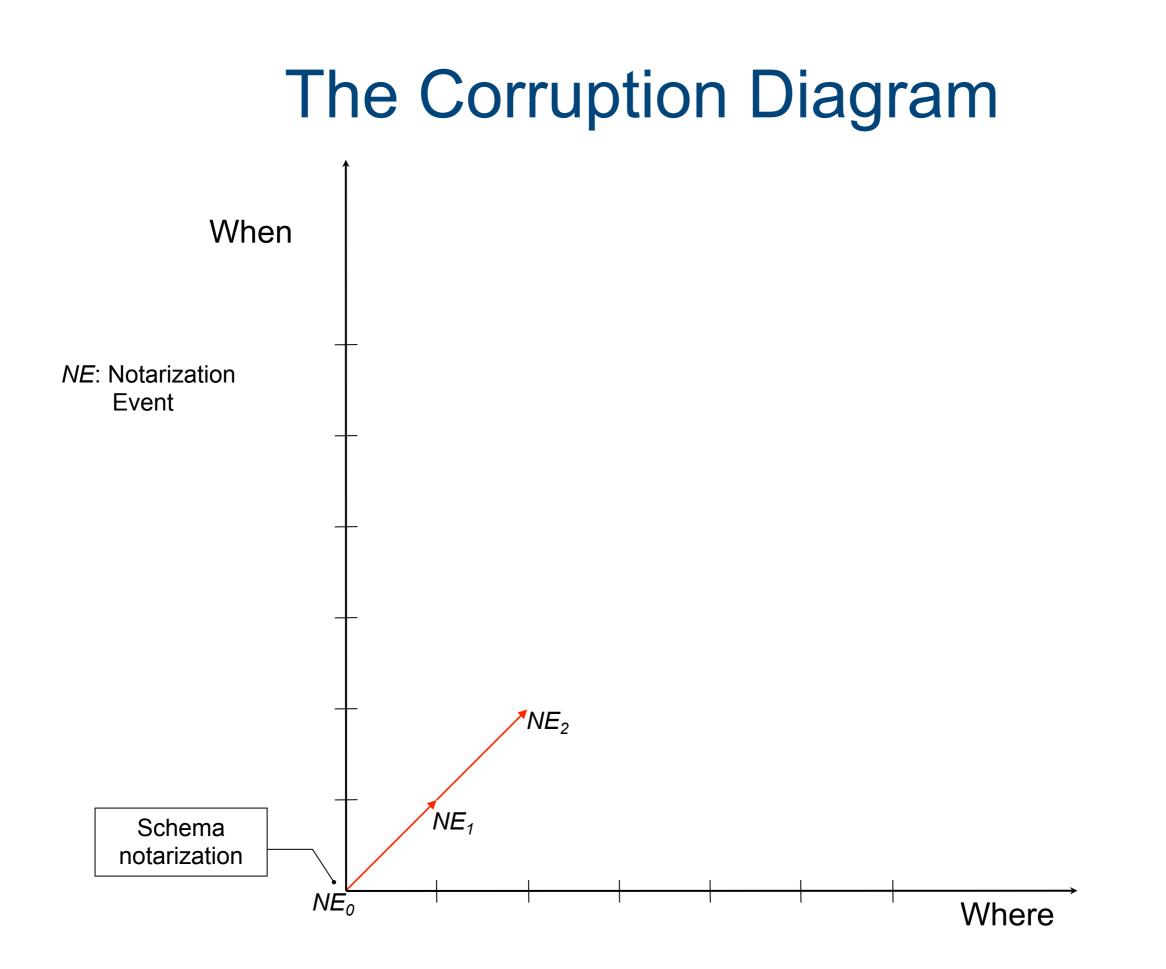


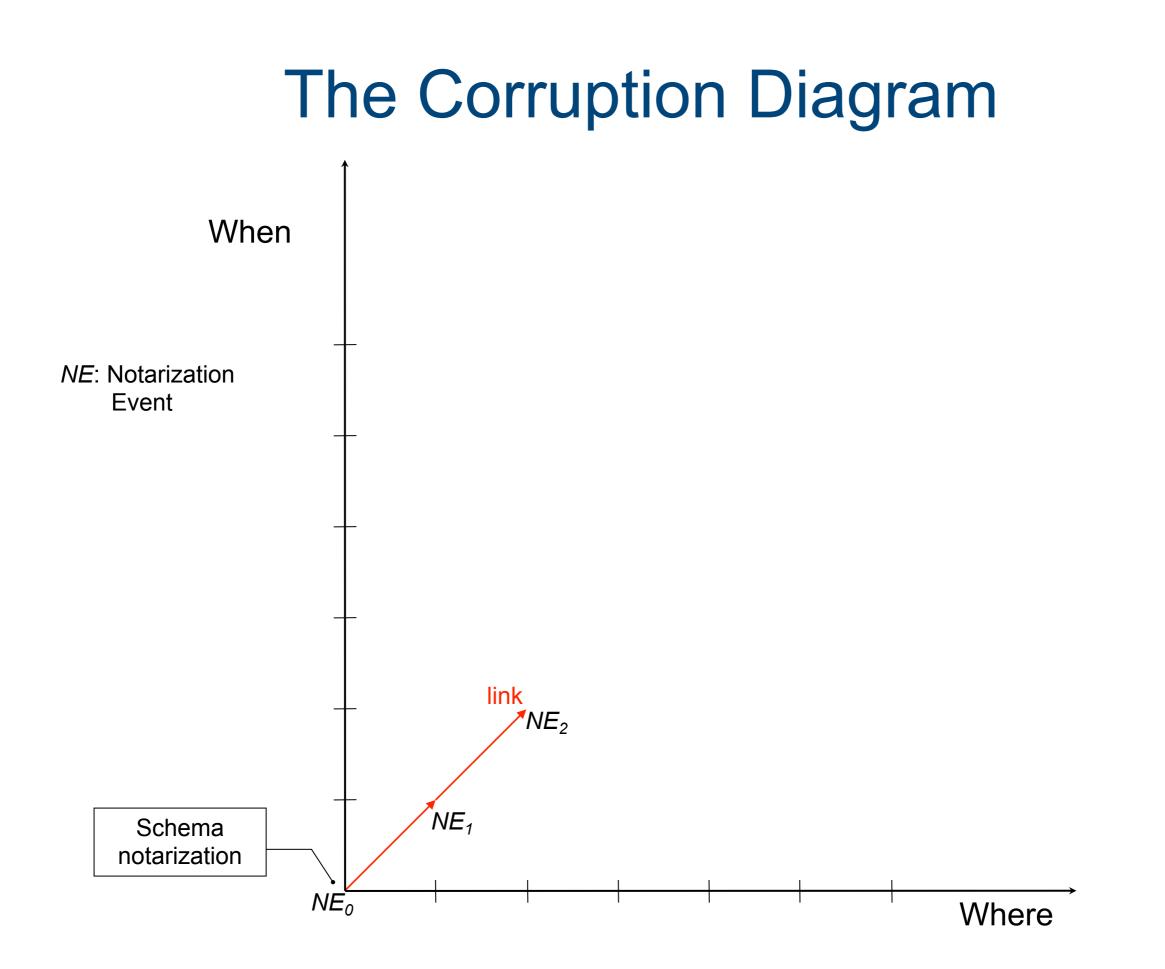


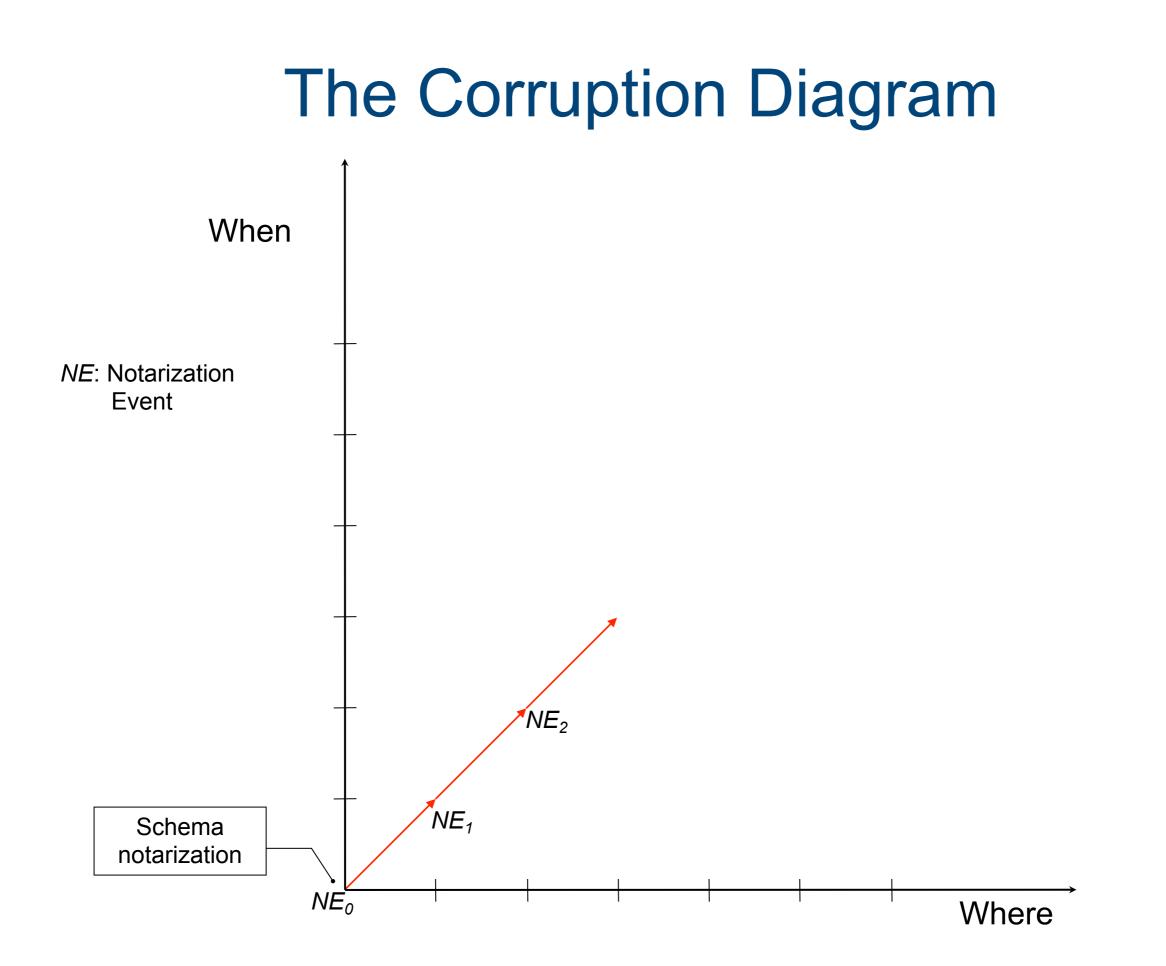


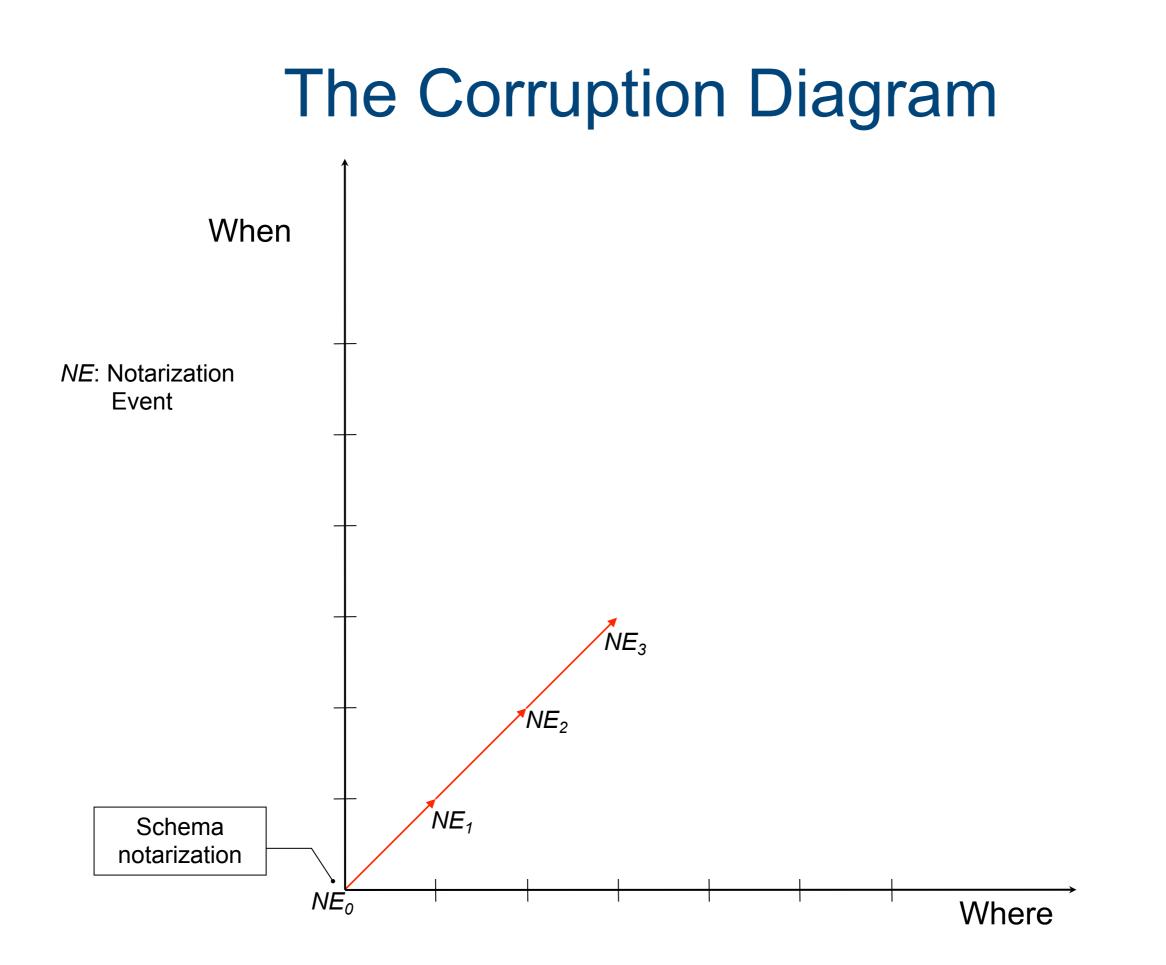


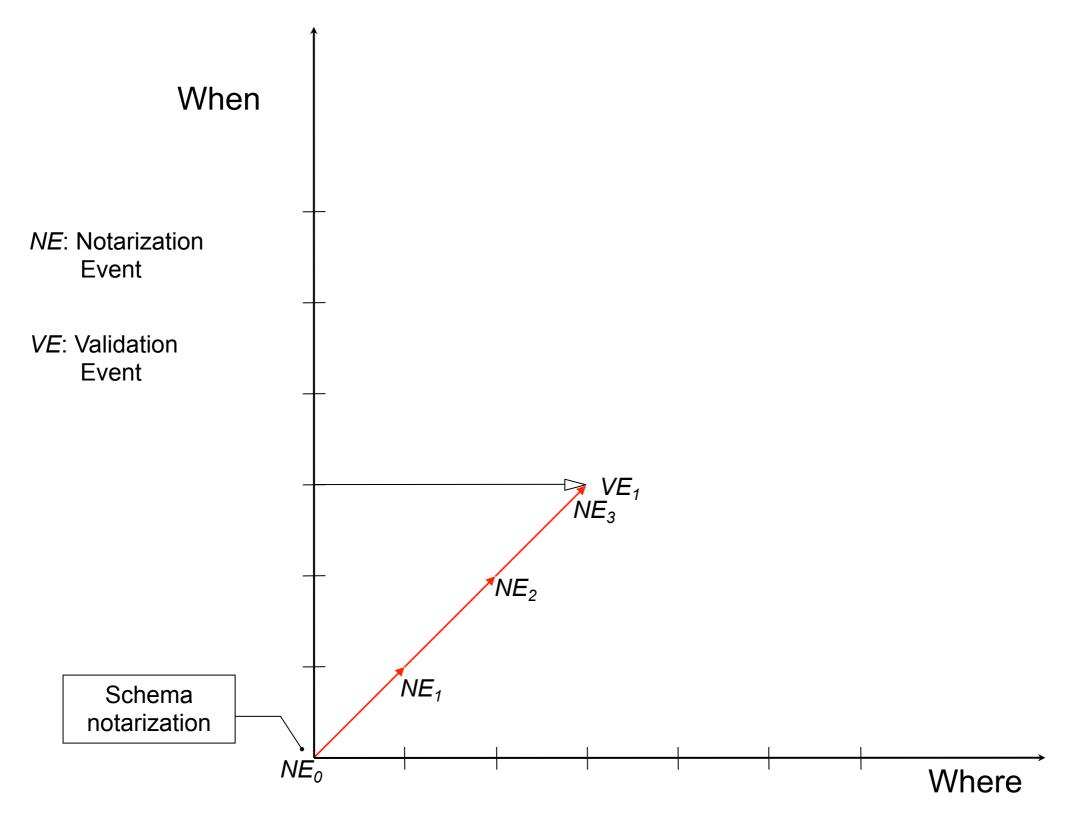


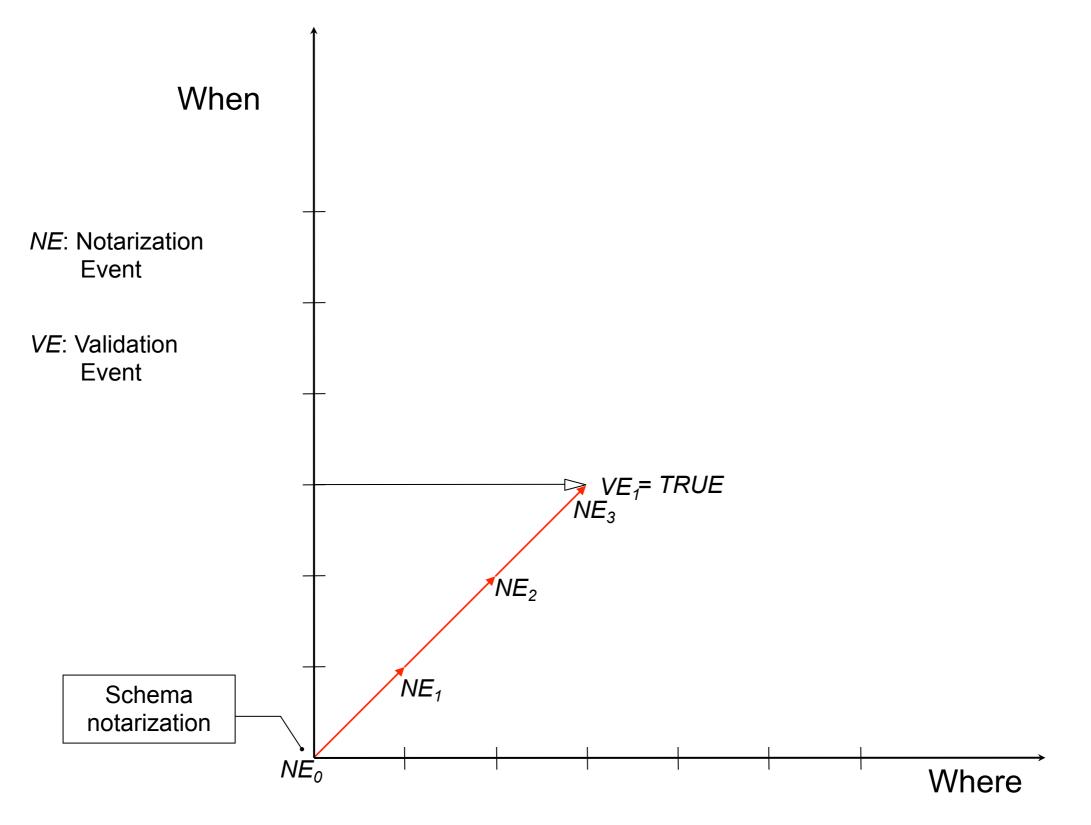


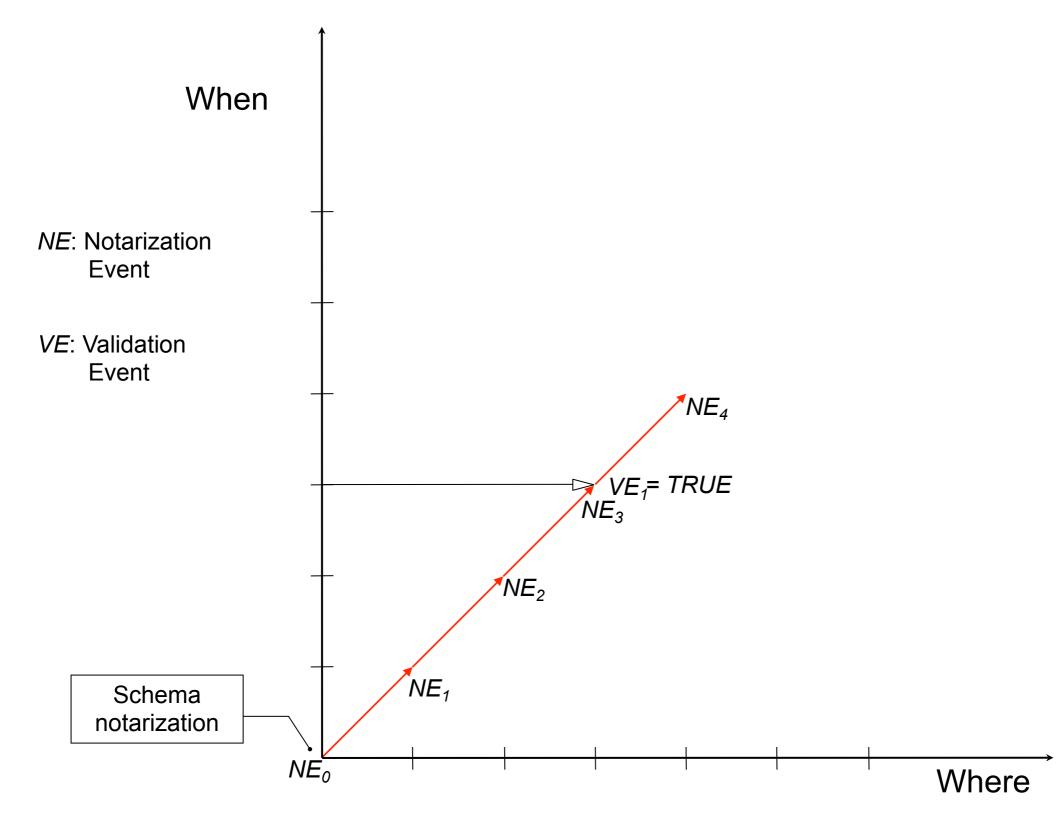


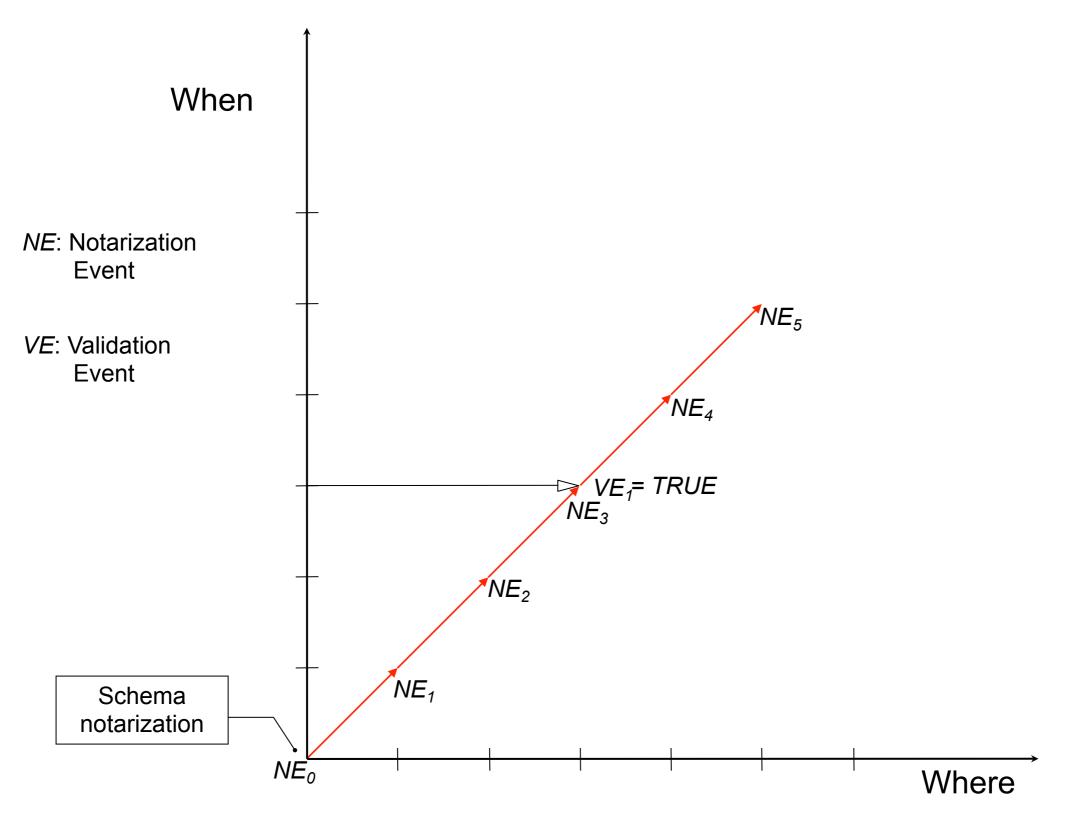


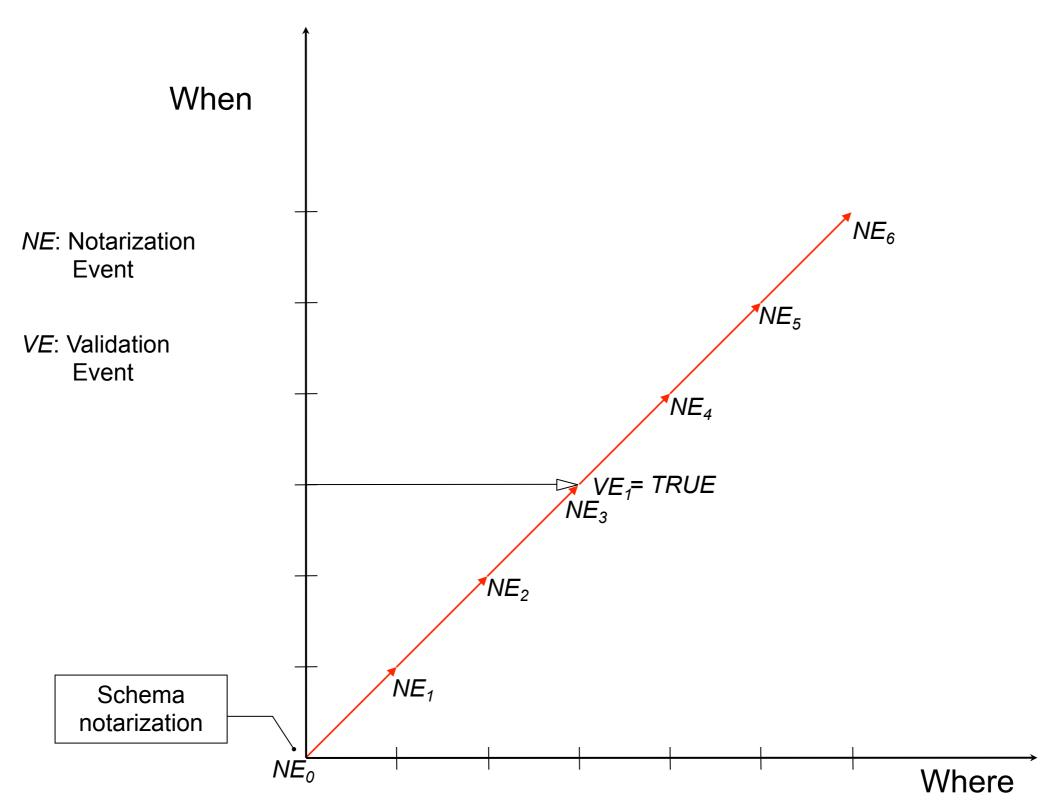


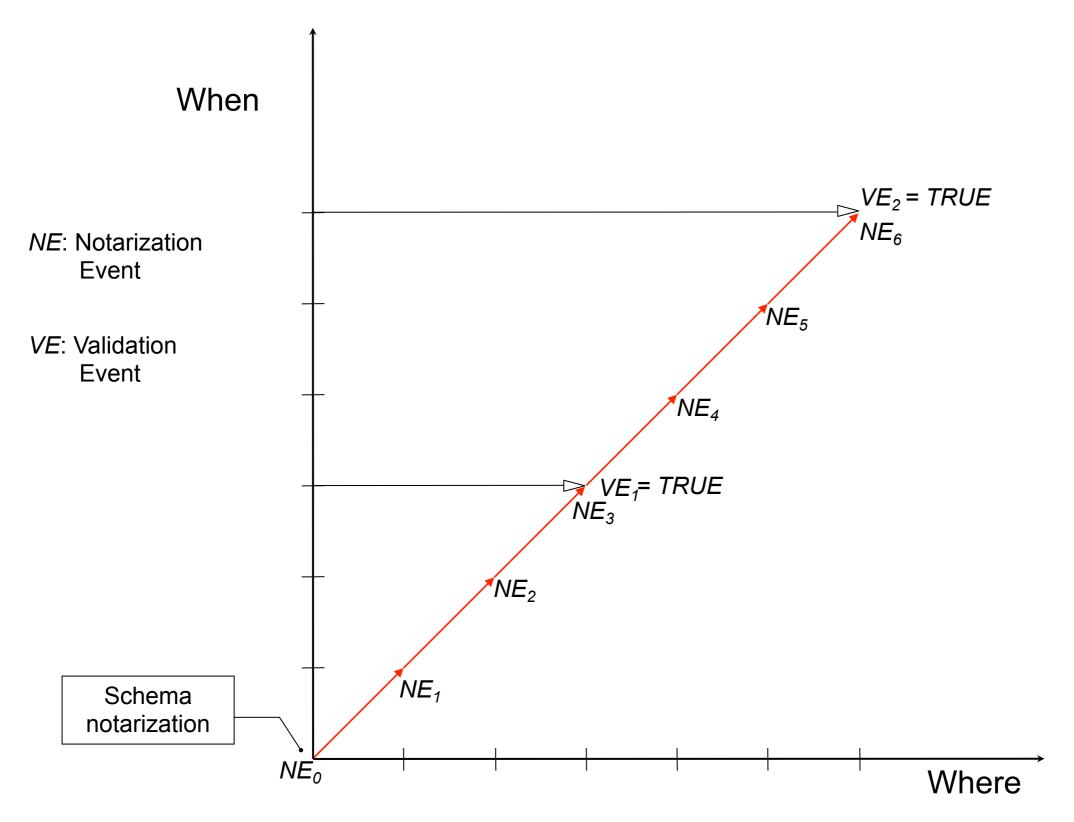


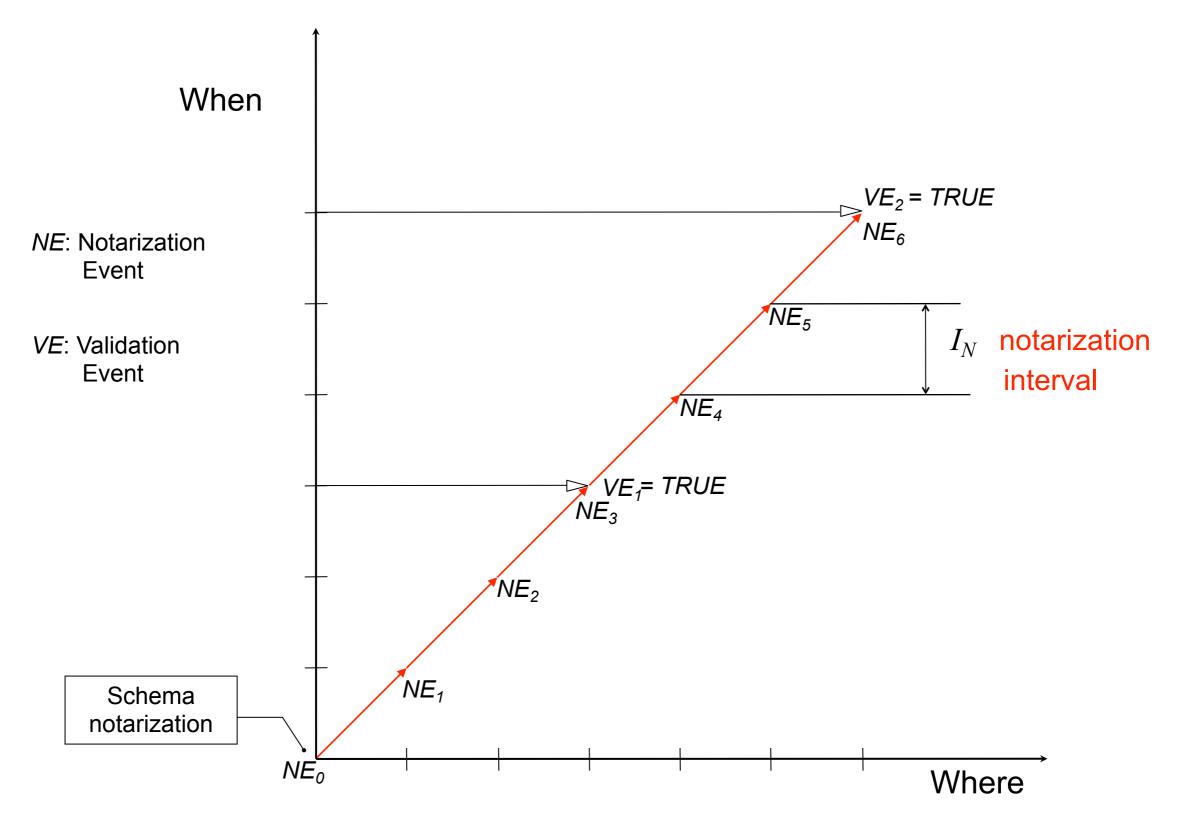


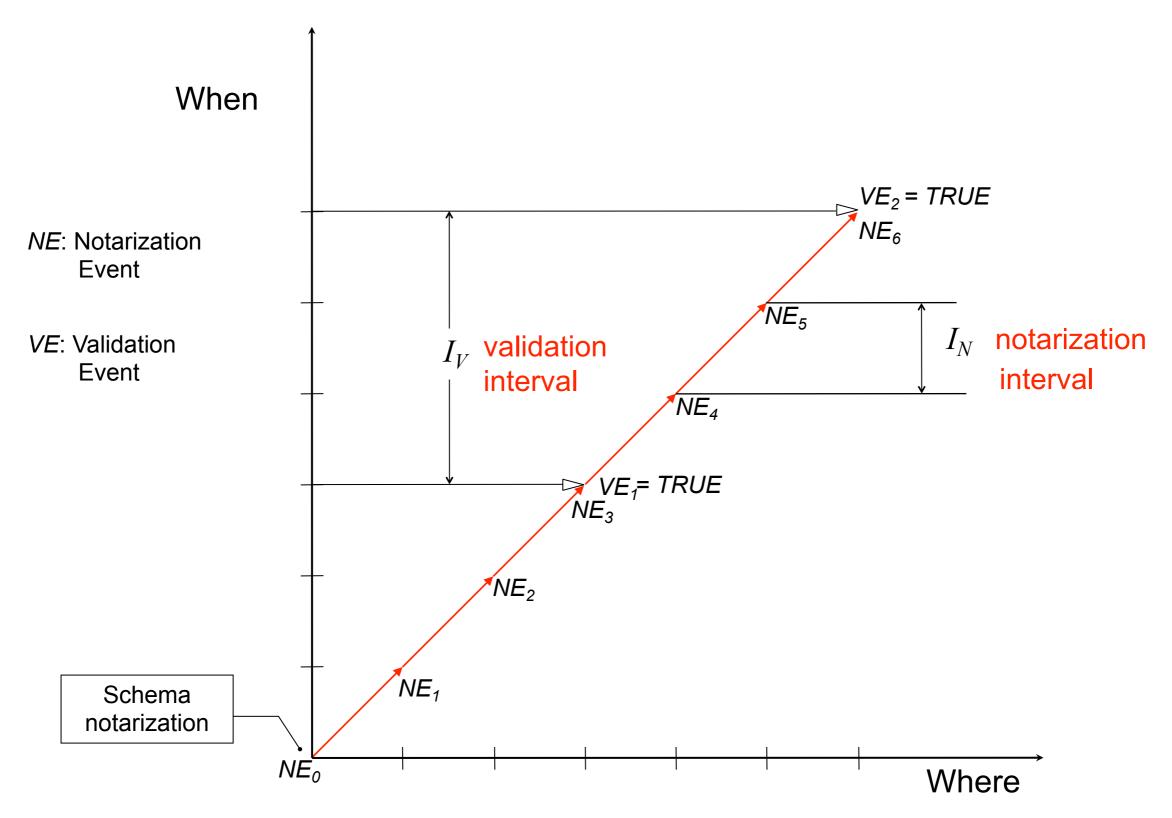


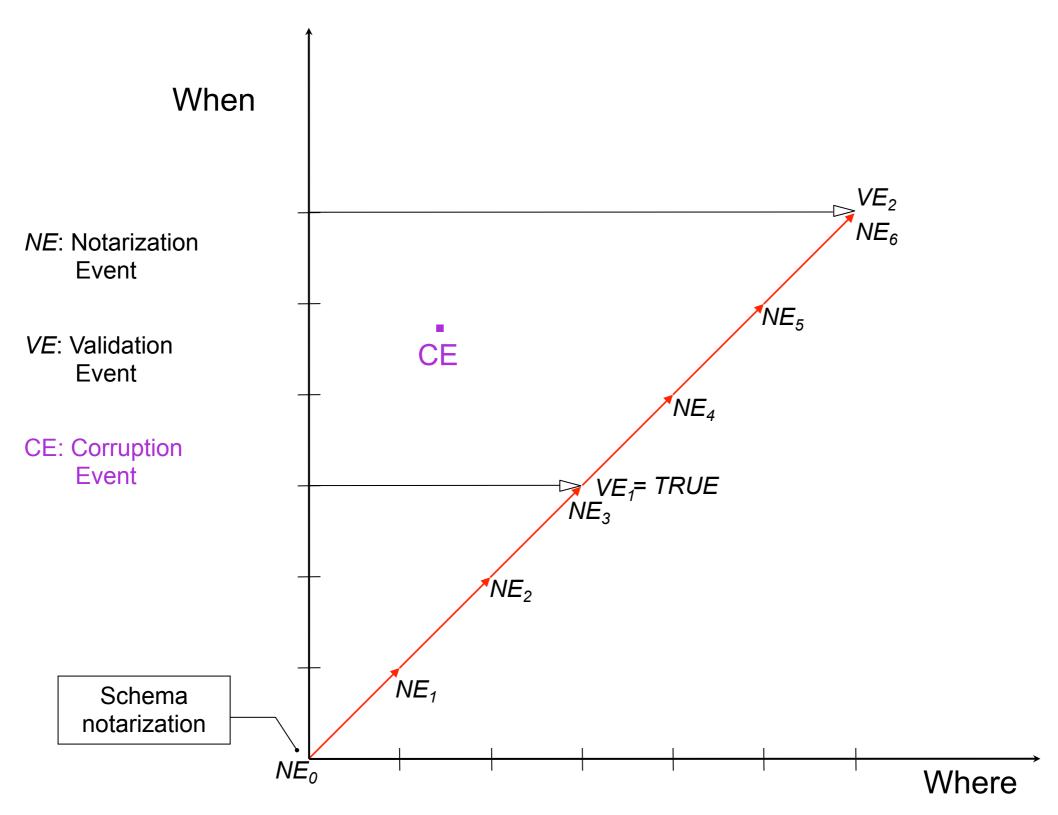


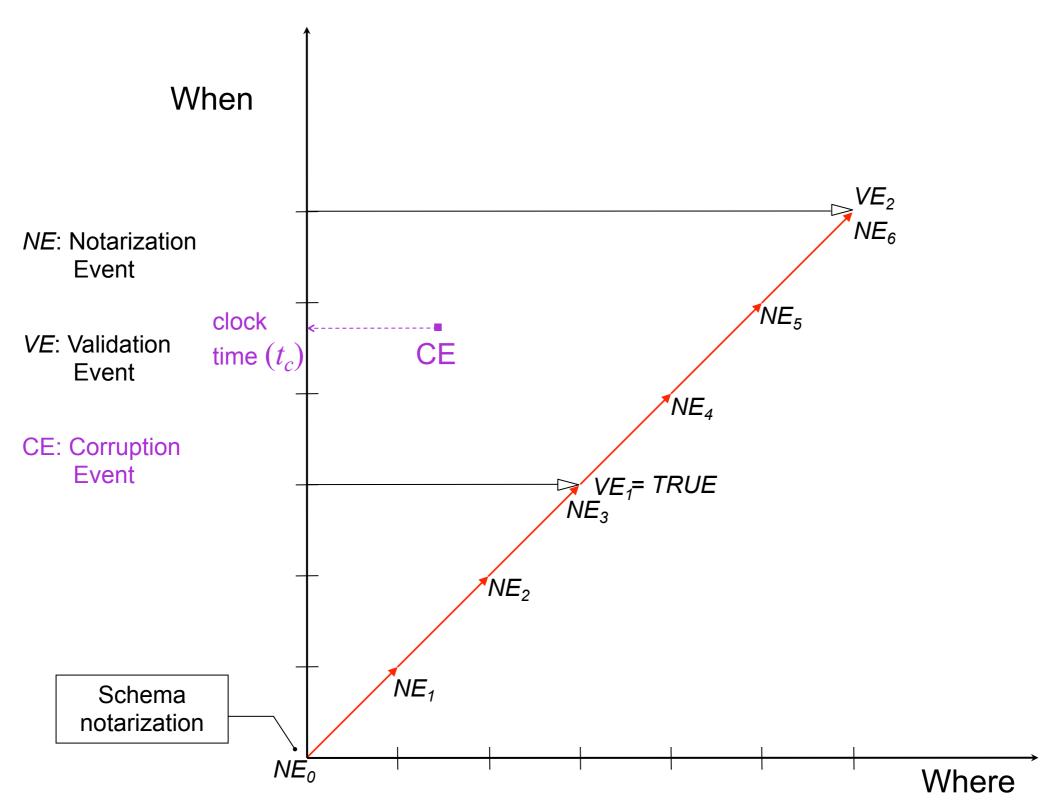


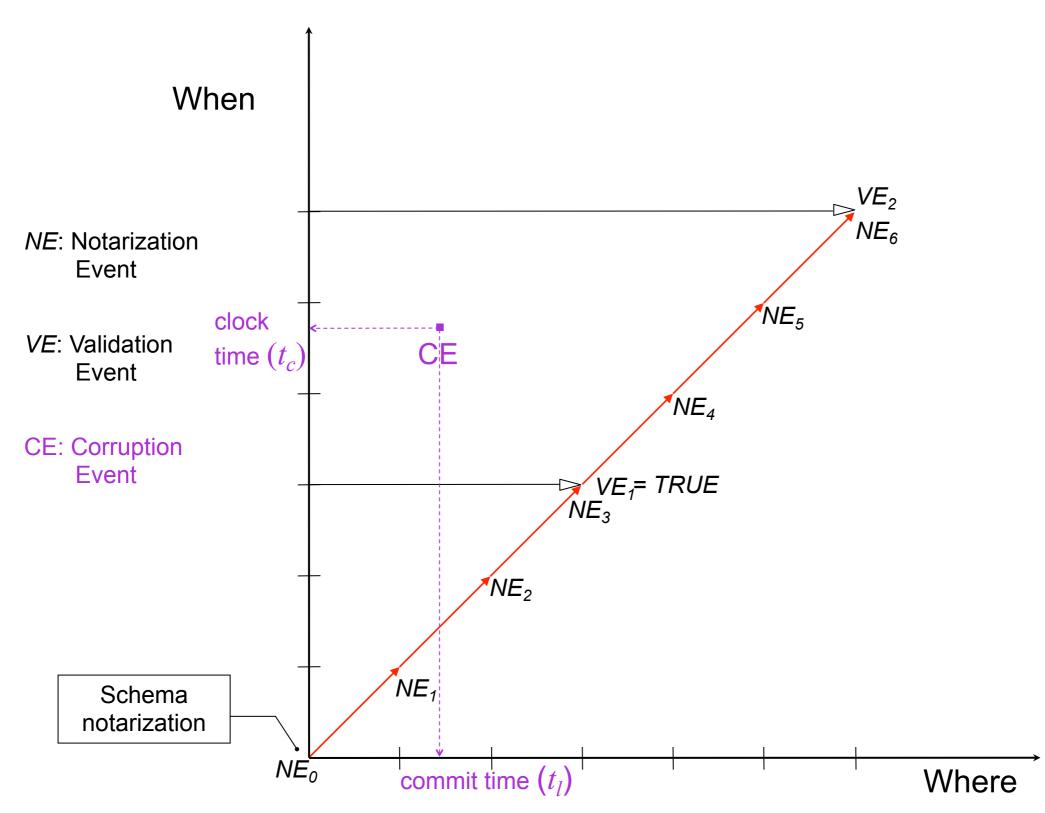


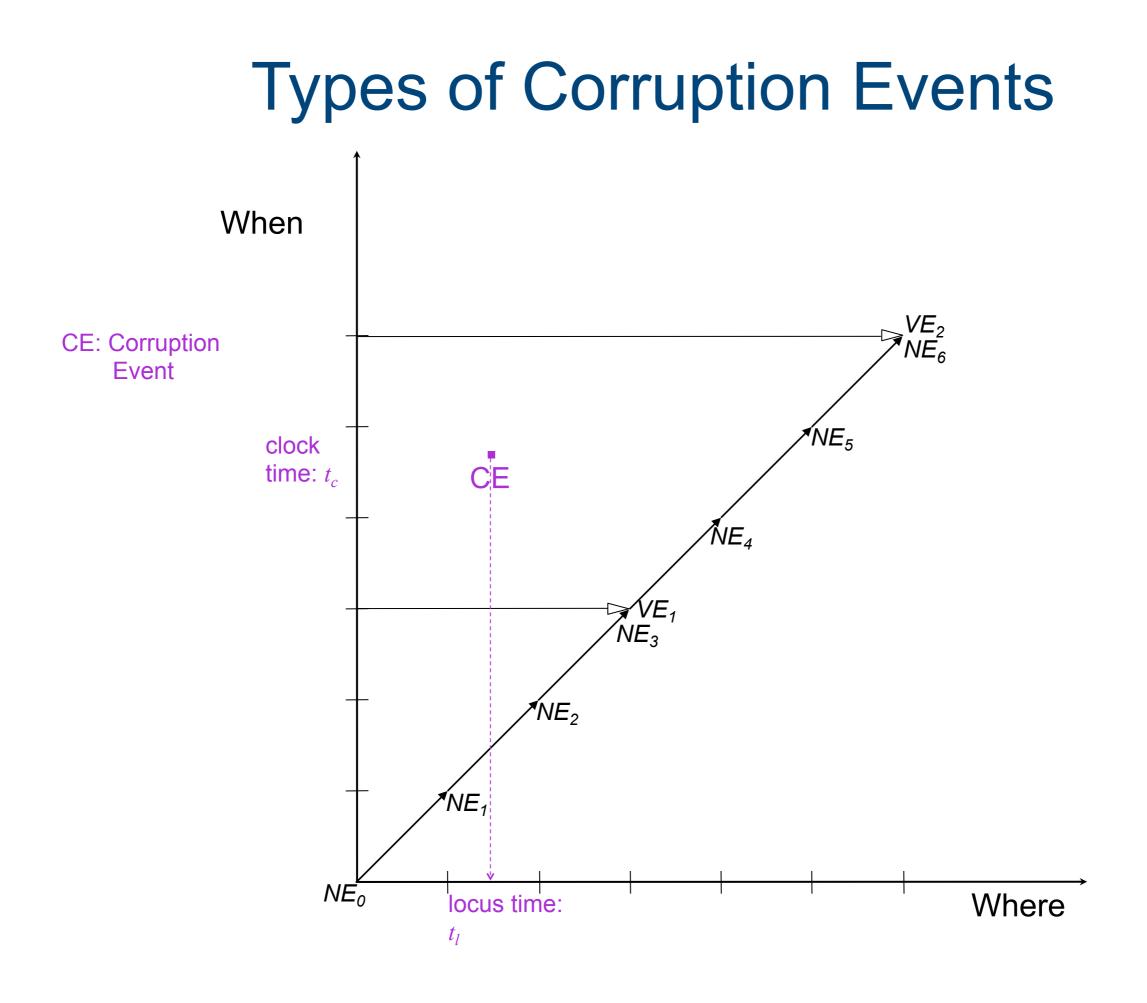






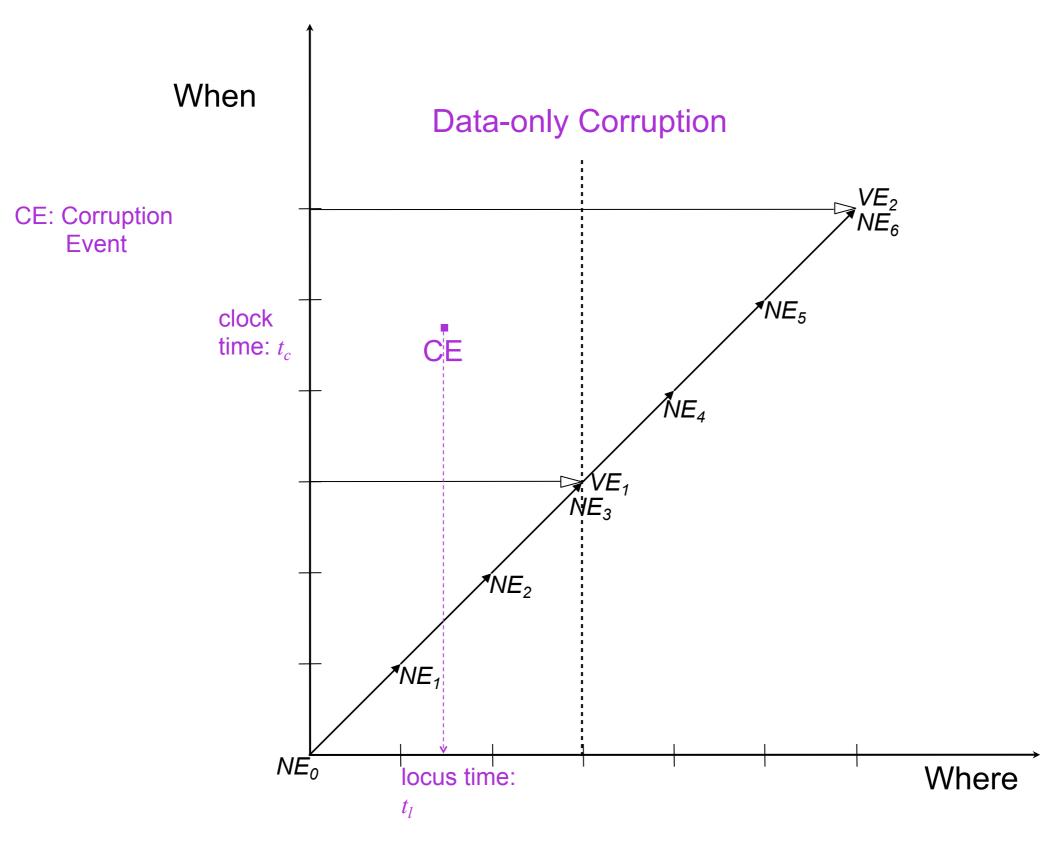


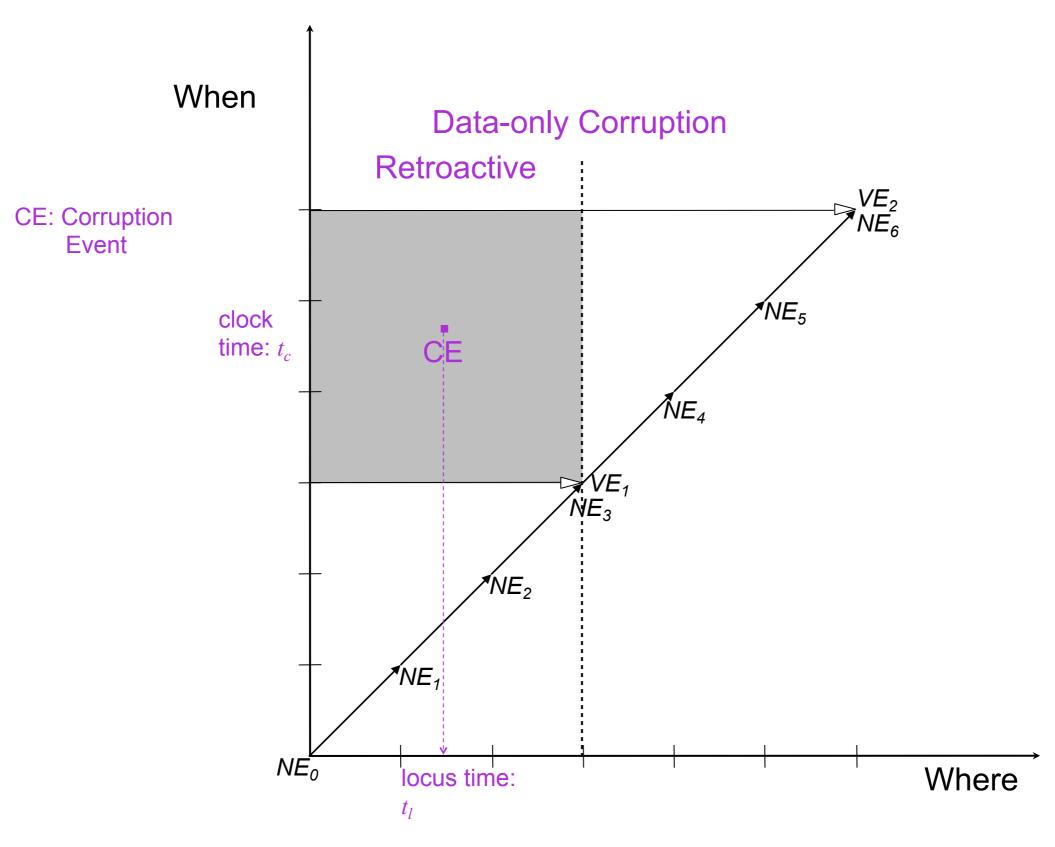


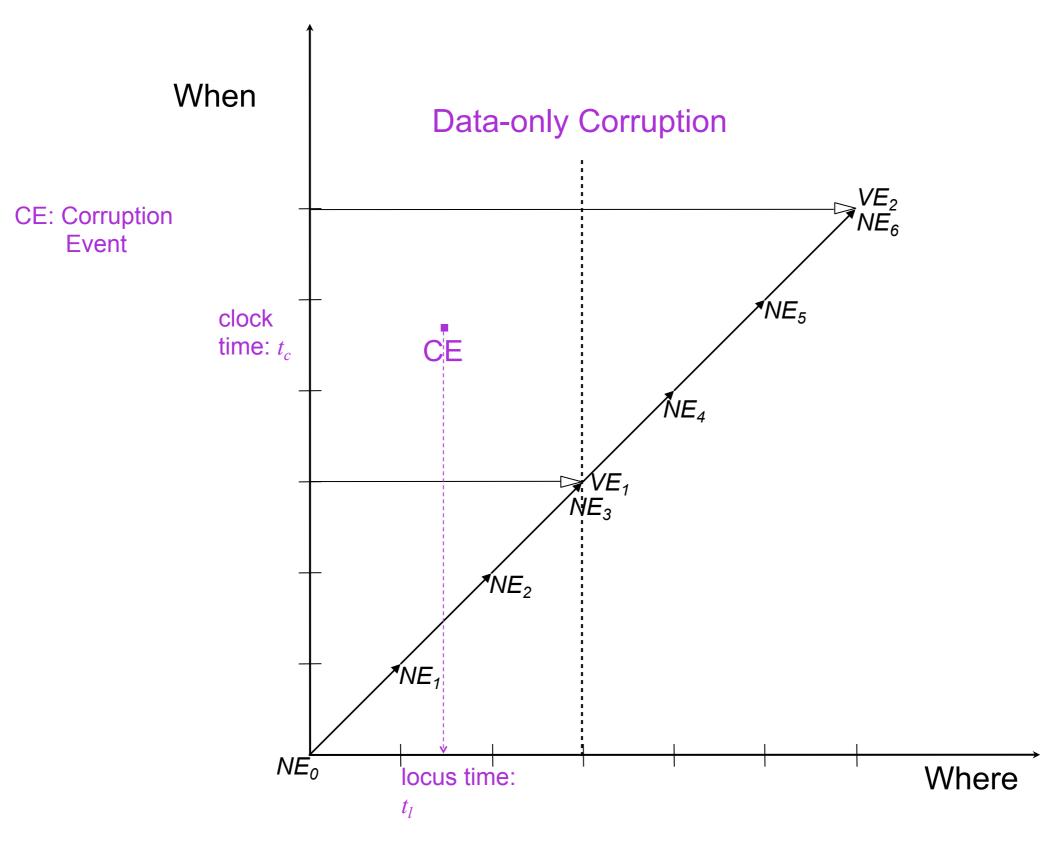


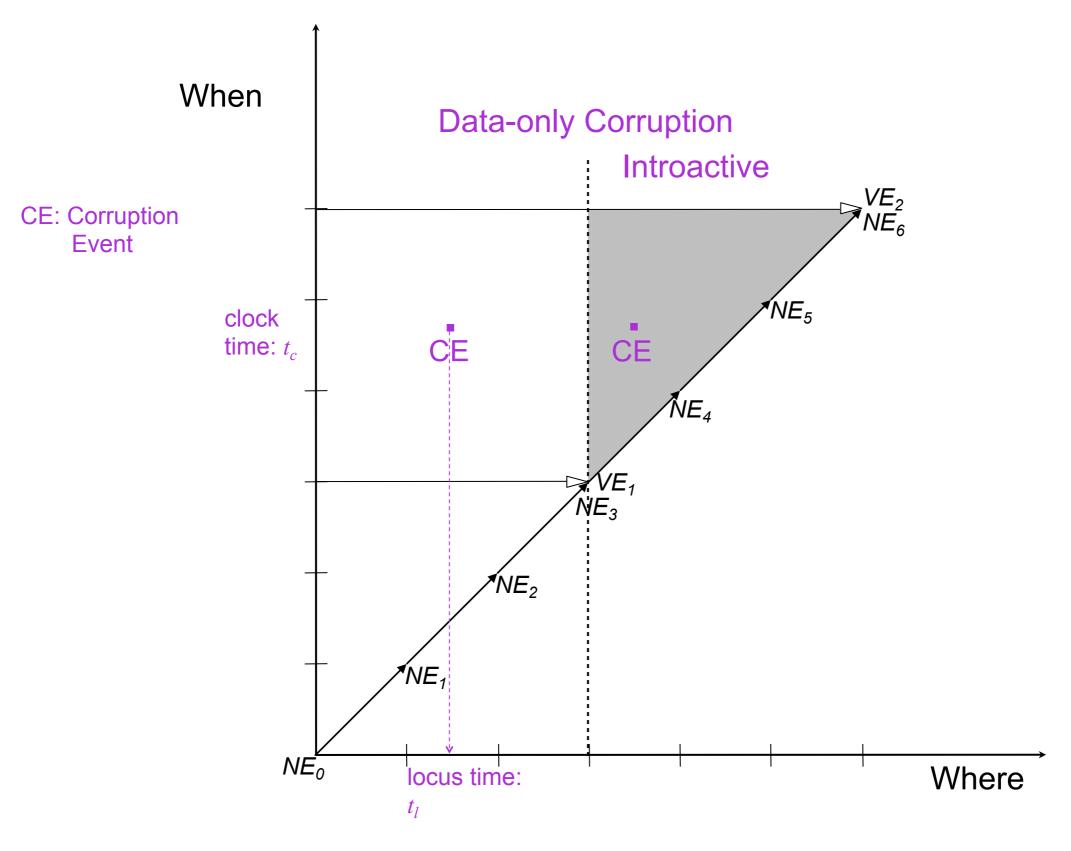
#### **Types of Corruption Events** When **Data-only Corruption** $VE_2$ $NE_6$ **CE:** Corruption **Event** NE<sub>5</sub> clock CE time: $t_c$ ŇE₄ NE<sub>3</sub> NE<sub>2</sub> NE<sub>1</sub> NE<sub>0</sub> locus time: Where

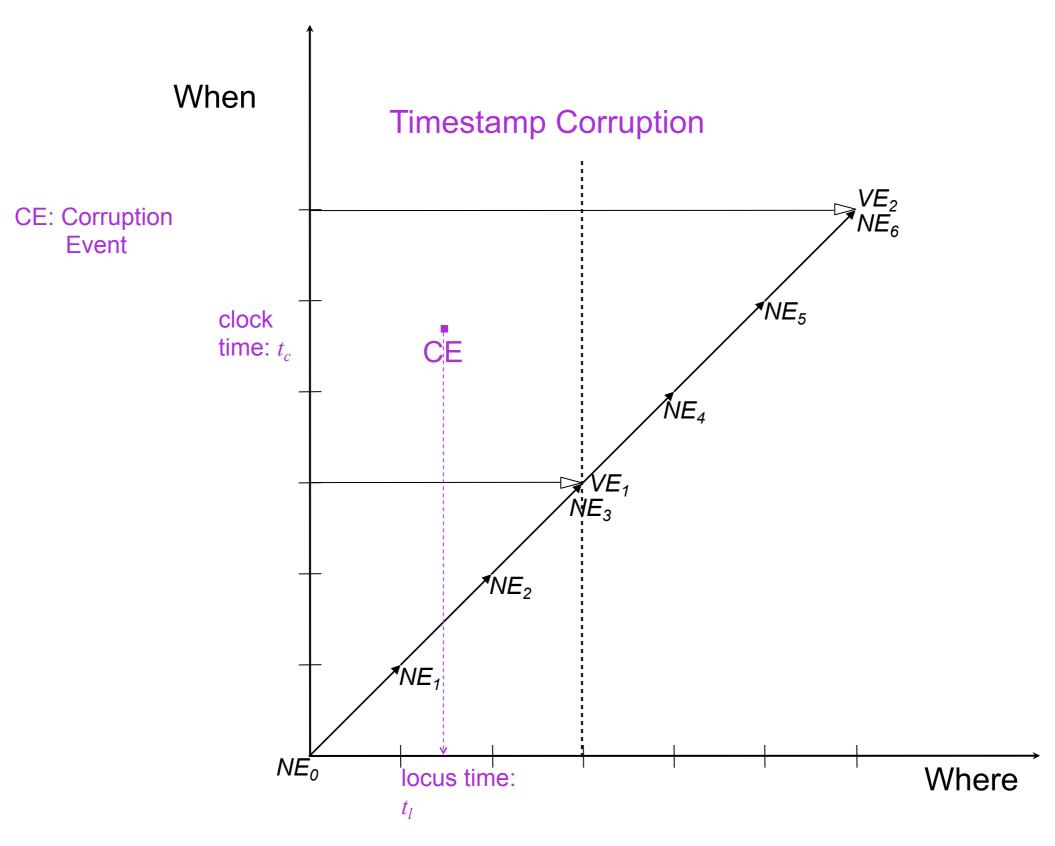
 $t_l$ 

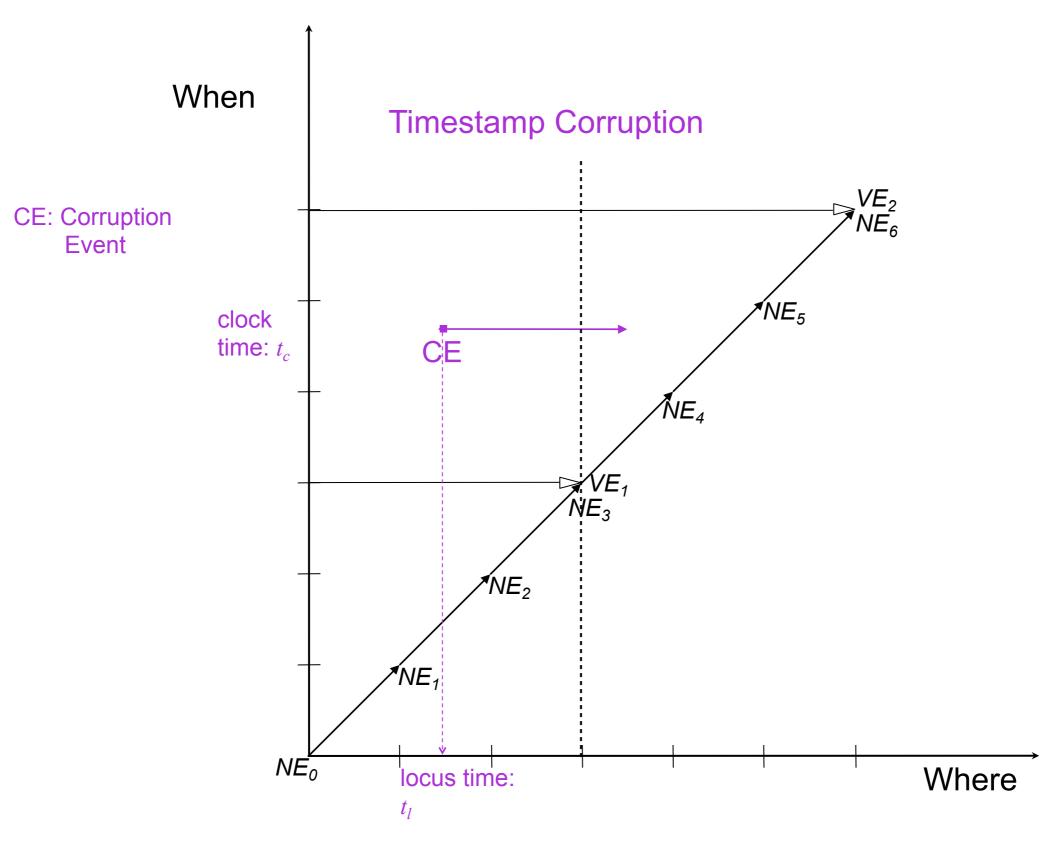


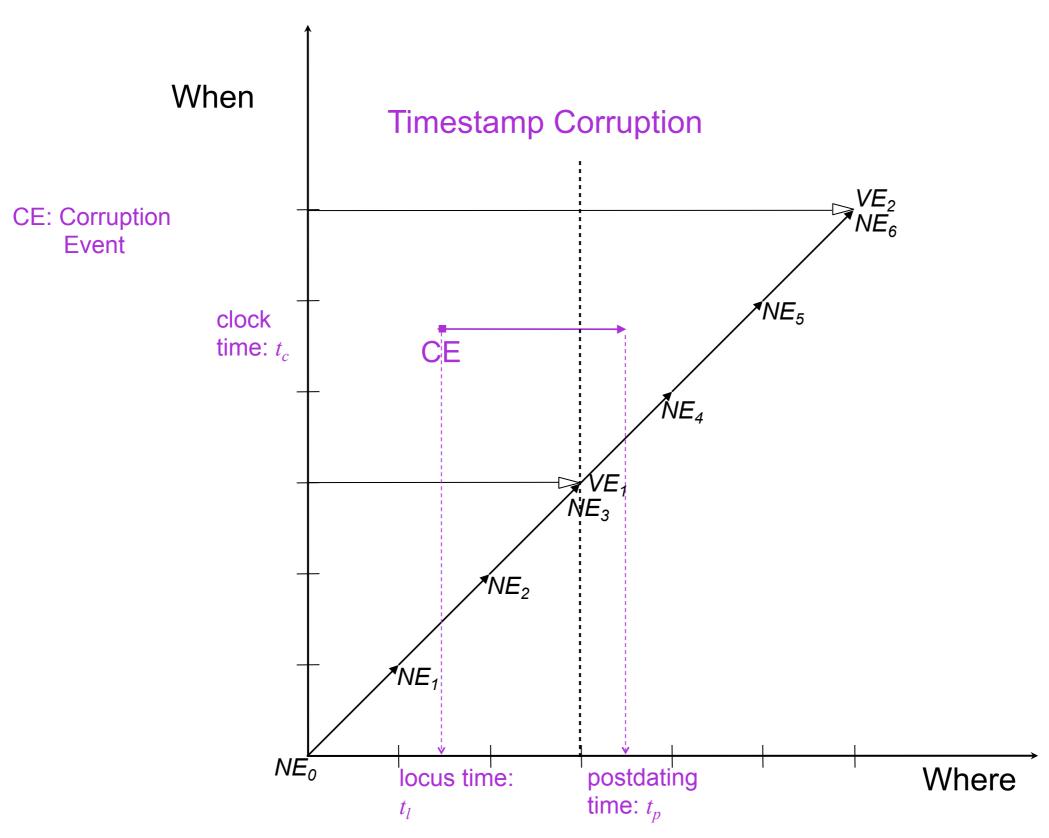




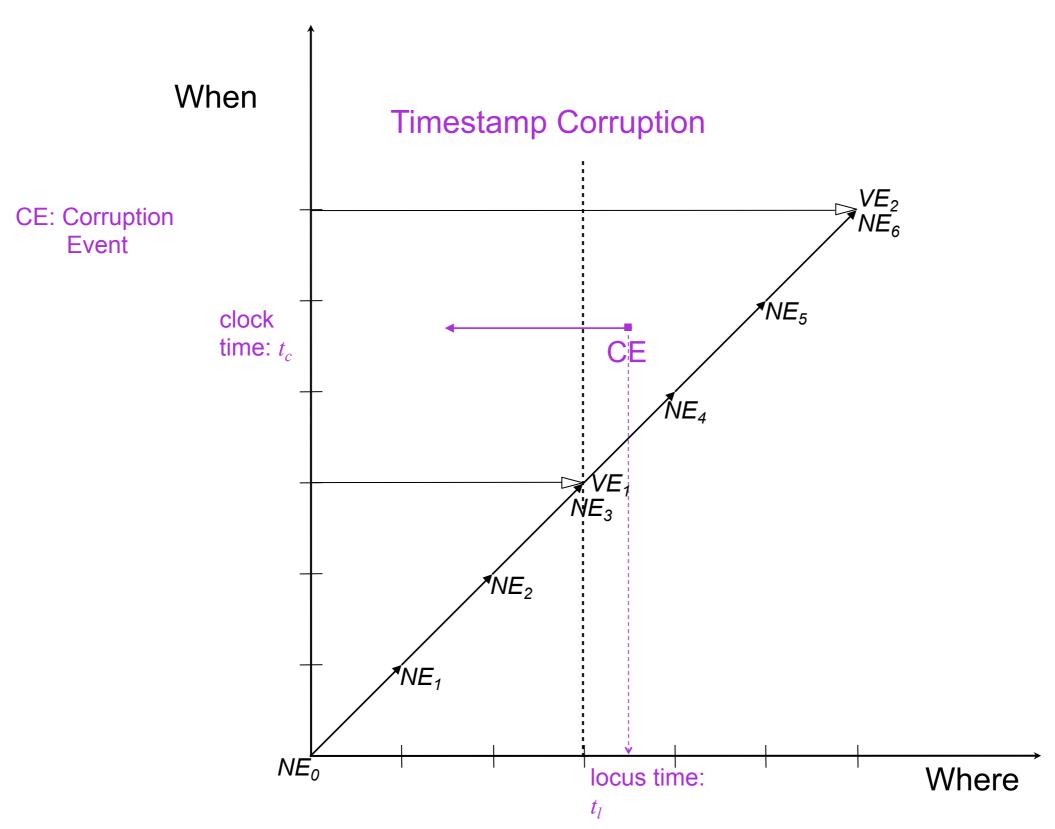


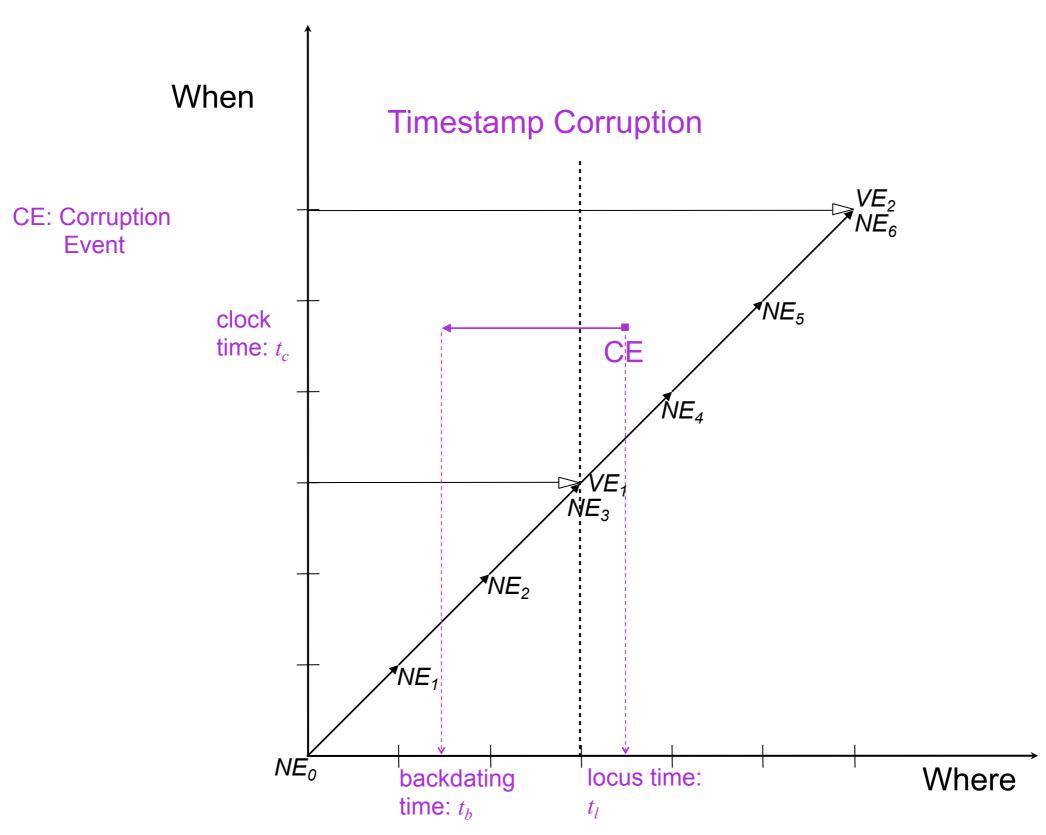






#### **Types of Corruption Events** When **Timestamp Corruption** $VE_2$ $NE_6$ **CE:** Corruption **Event** NE<sub>5</sub> clock CE time: $t_c$ NE₄ ¥VE₁ NE₃ NE<sub>2</sub> NE<sub>1</sub> NE<sub>0</sub> locus time: Where $t_l$





• If a corruption is detected, then the forensic analysis phase begins.

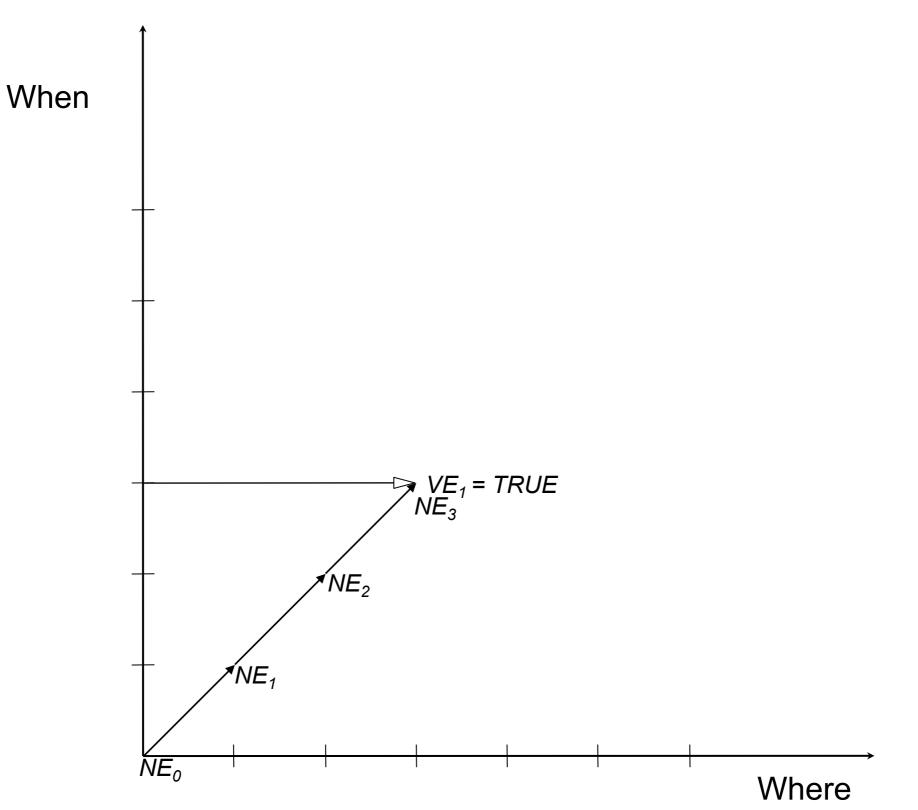
- If a corruption is detected, then the forensic analysis phase begins.
- A *forensic analysis algorithm* is run as directed by the Database Administrator.

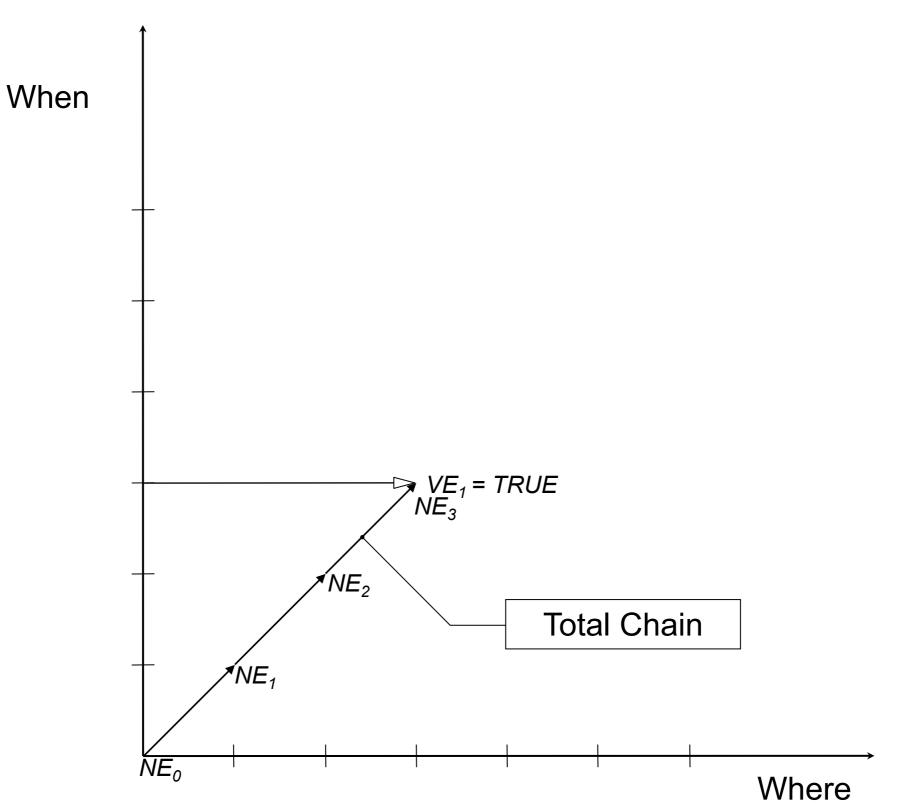
- If a corruption is detected, then the forensic analysis phase begins.
- A *forensic analysis algorithm* is run as directed by the Database Administrator.
- Attempt to ascertain a *corruption region*: the bounds on the uncertainty of the "where" and "when" of the corruption.

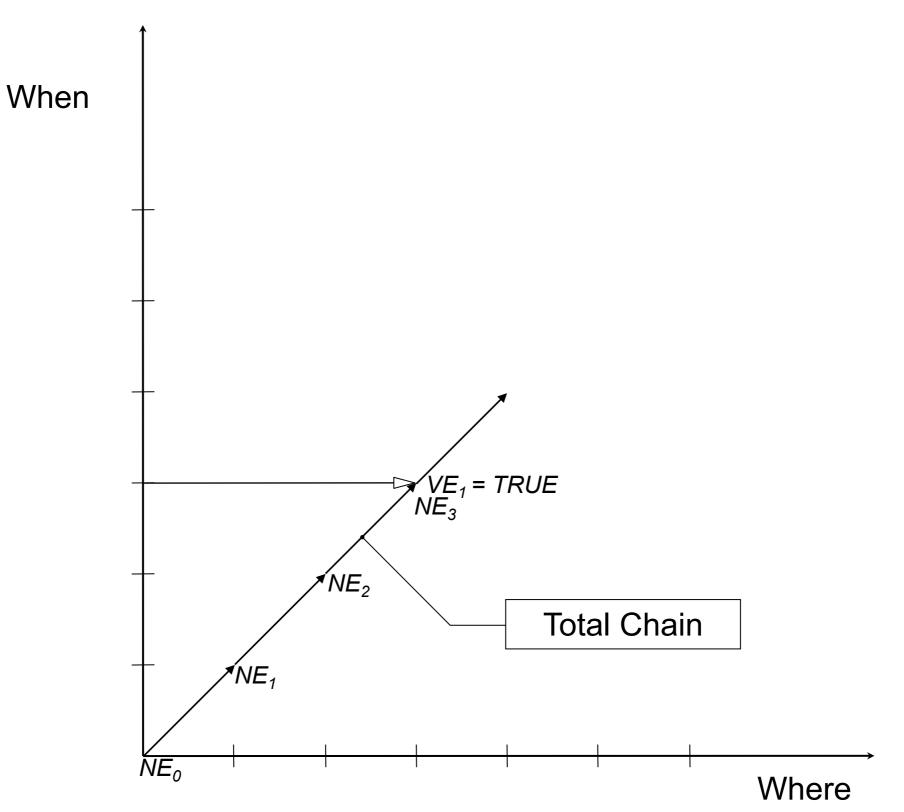
#### **Detection Resolution**

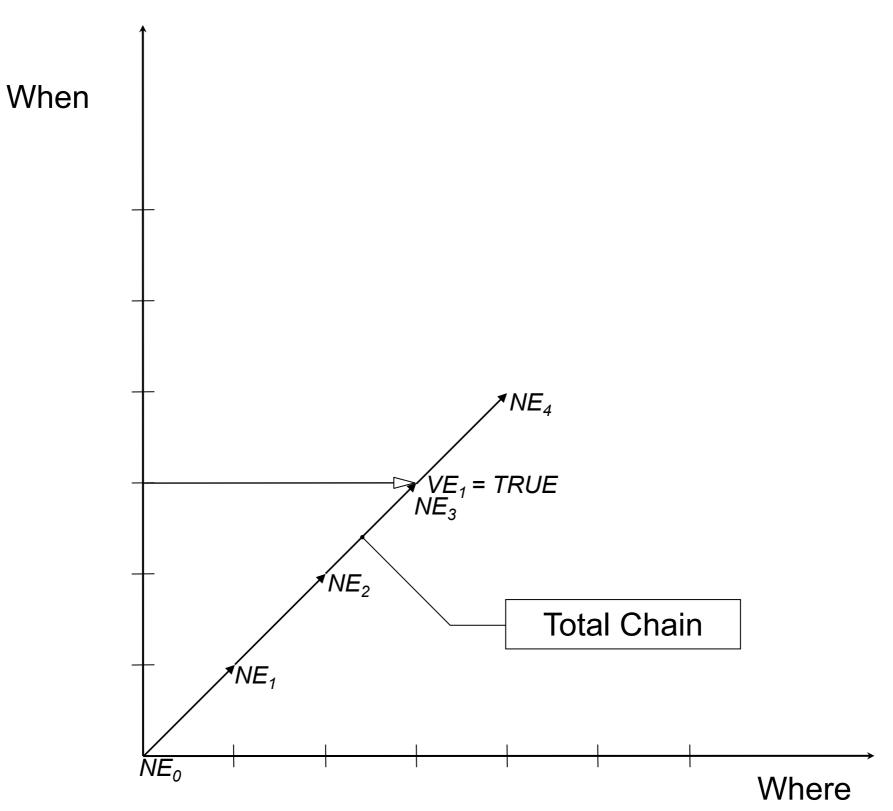
• Temporal Detection Resolution  $(R_t)$ : the finest granularity of temporal bounds uncertainty of a CE.

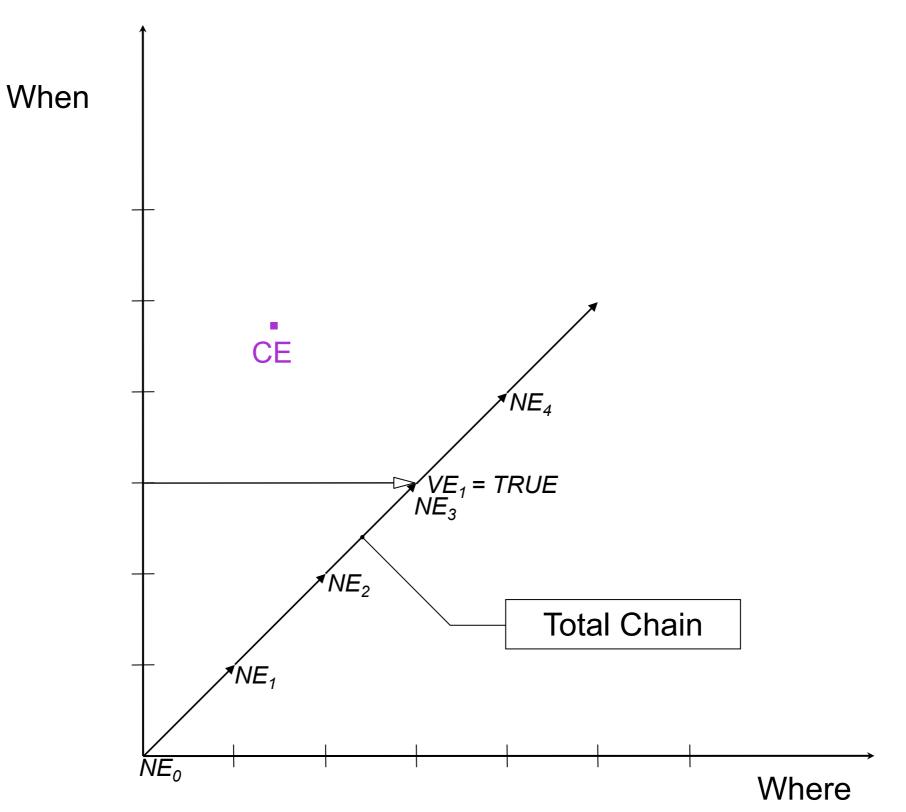
• Spatial Detection Resolution ( $R_s$ ): the finest granularity of spatial bounds uncertainty of a CE.

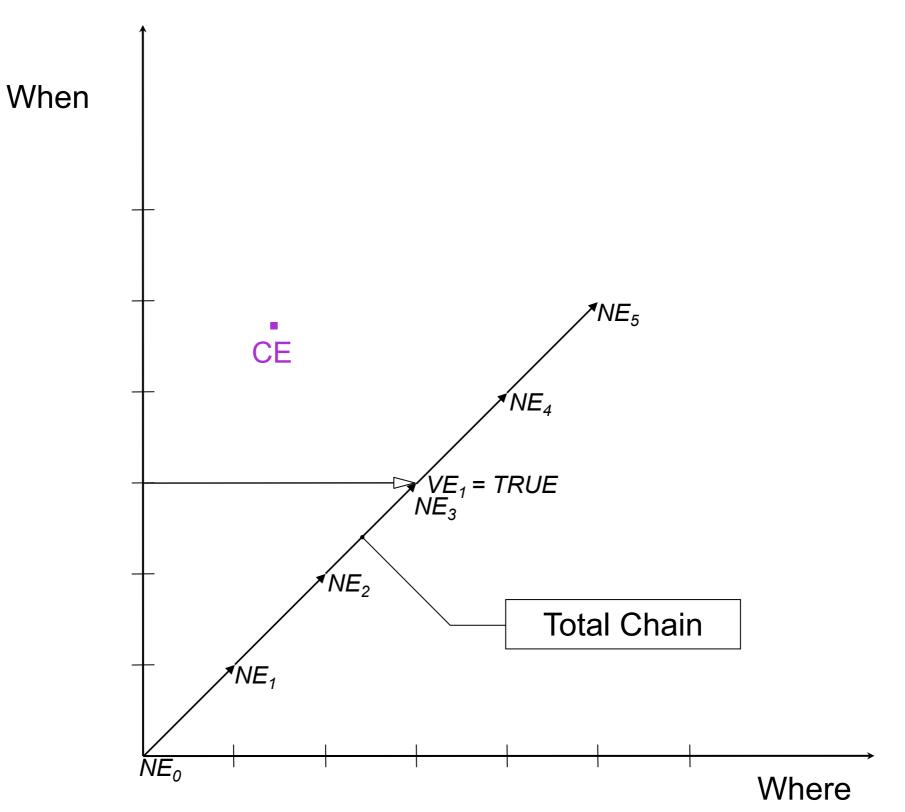


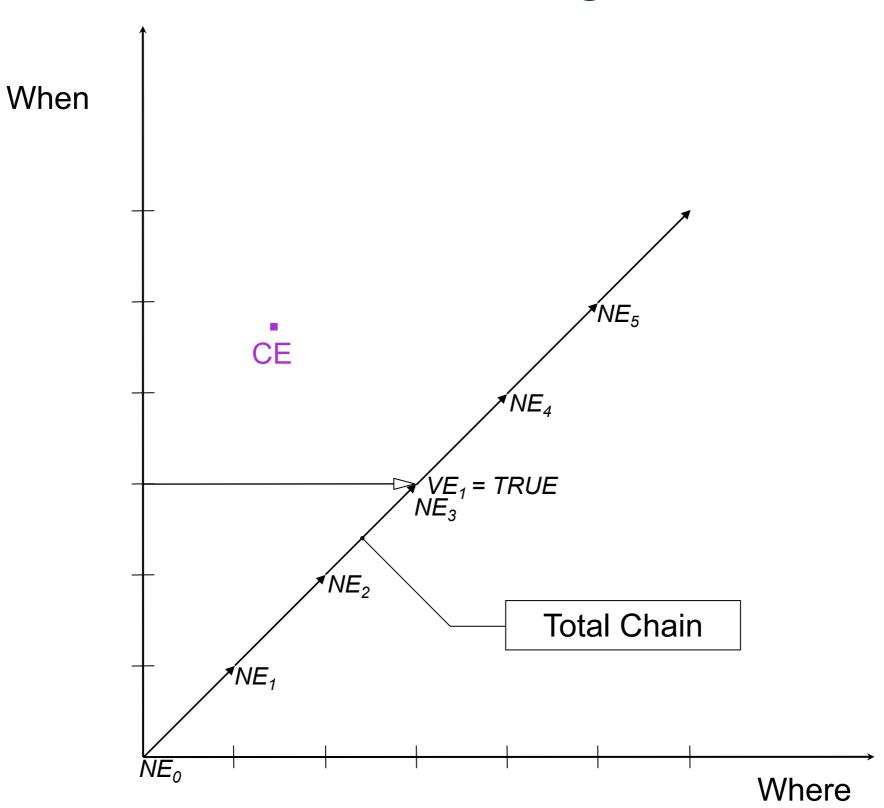


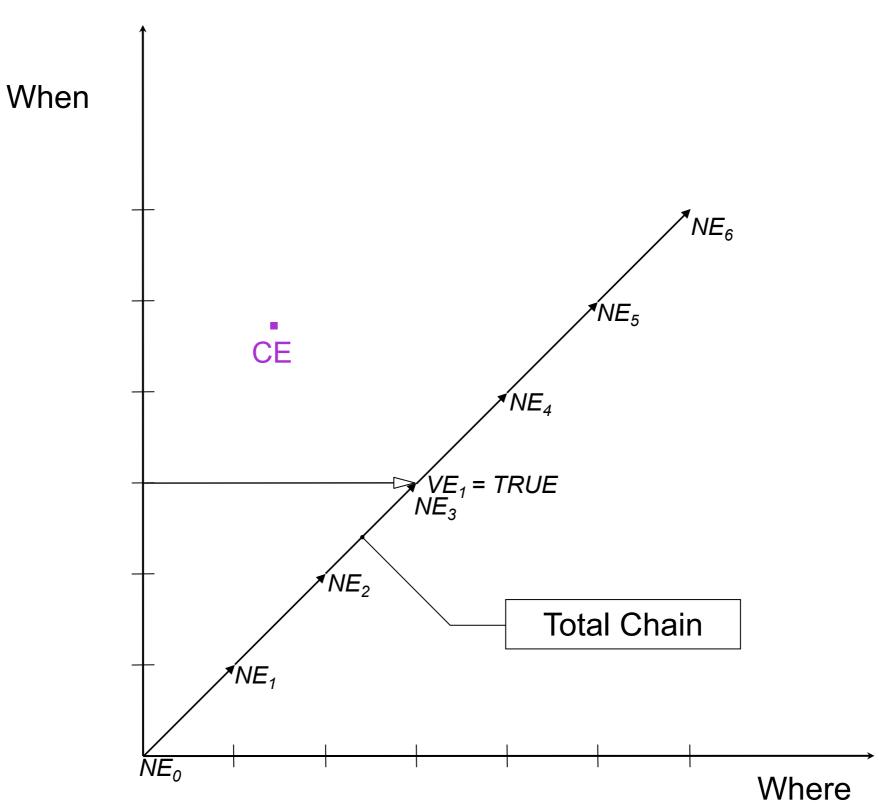


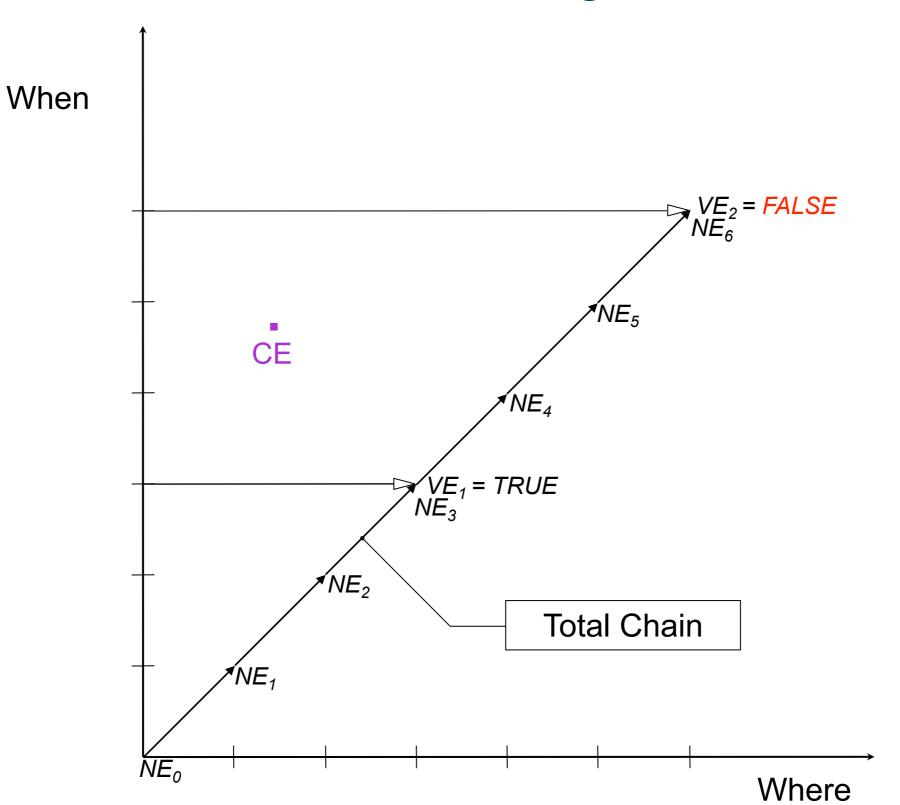


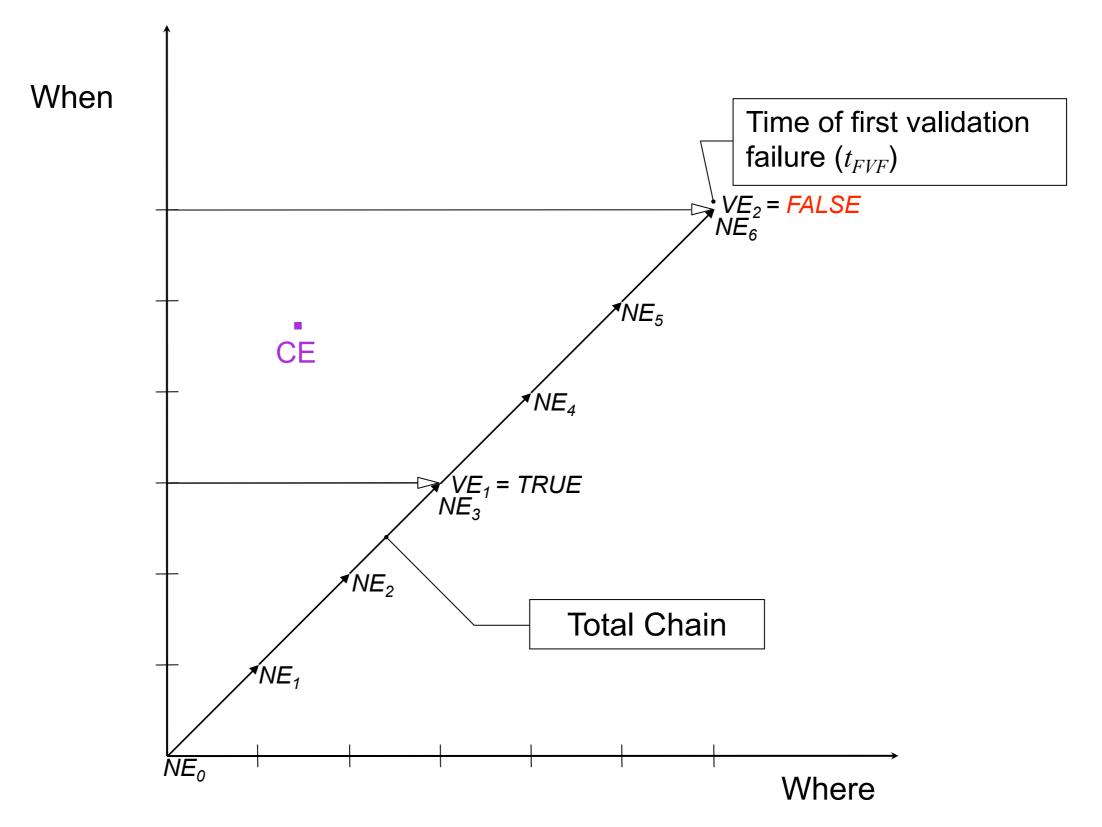


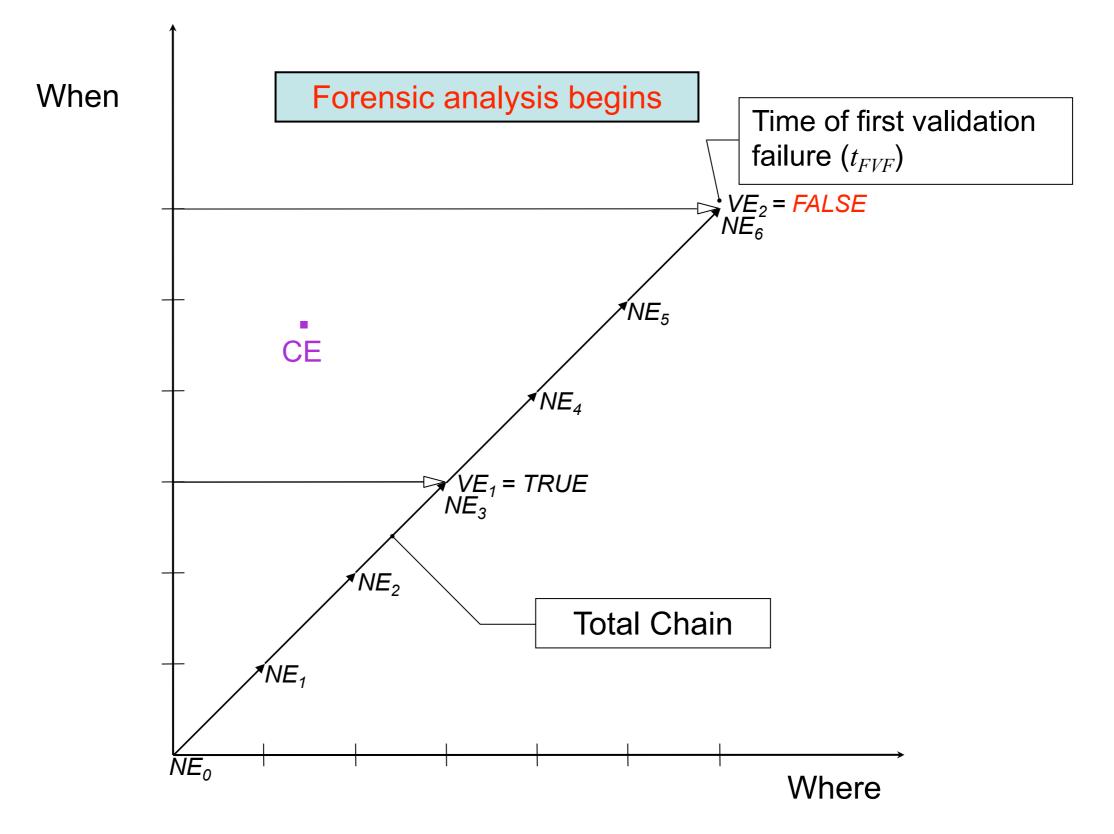


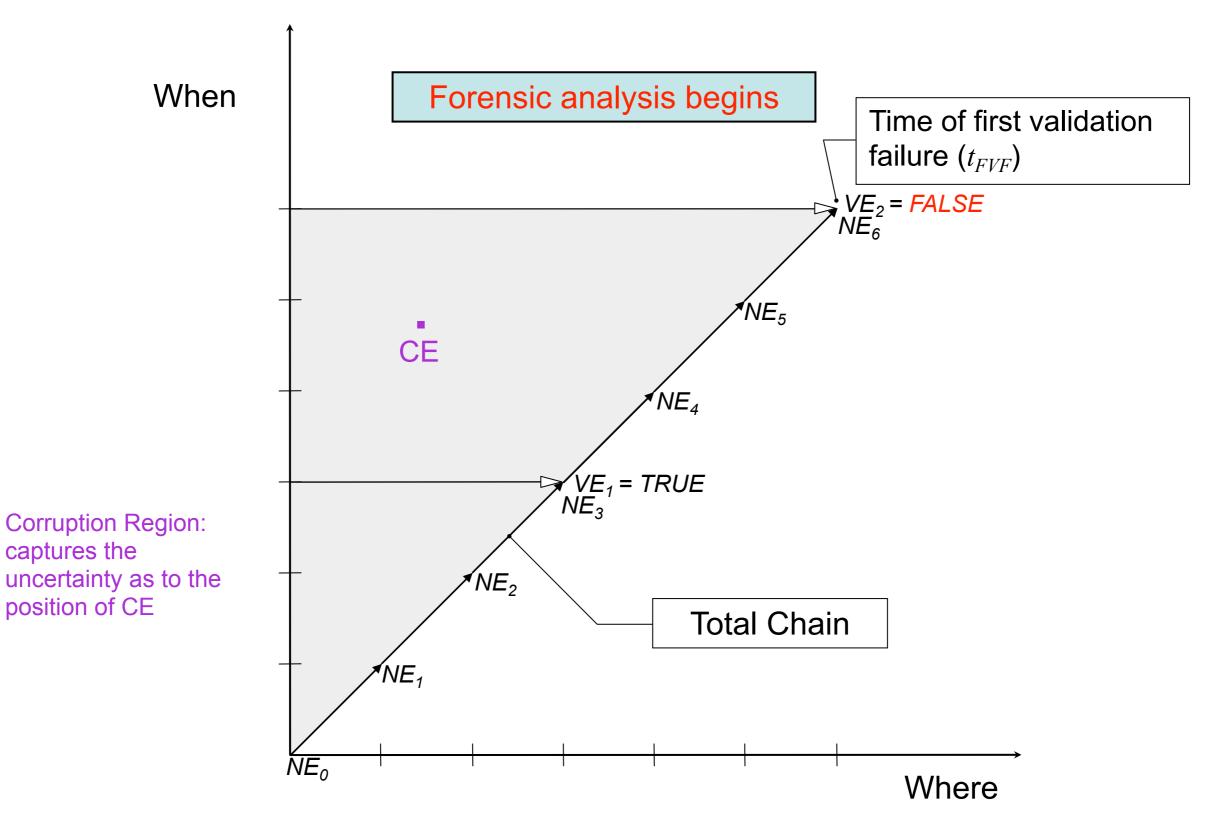


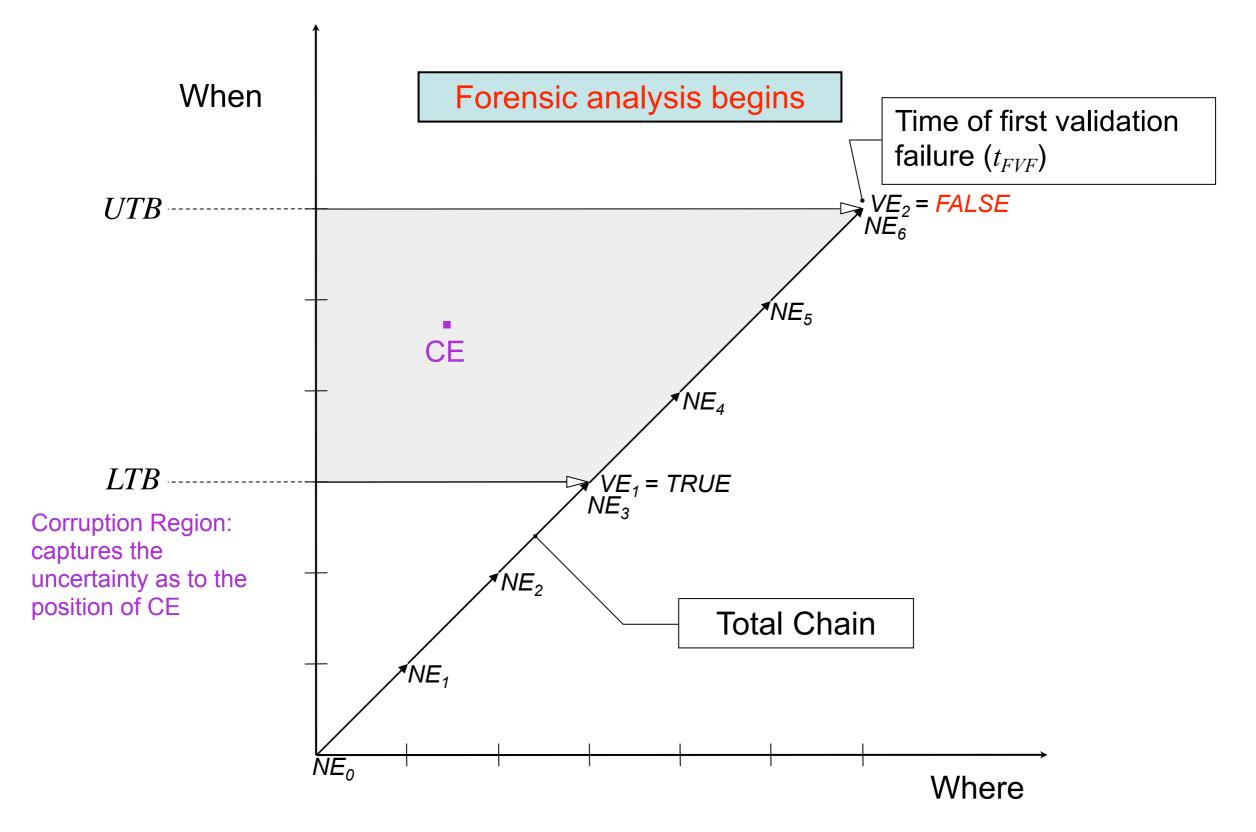


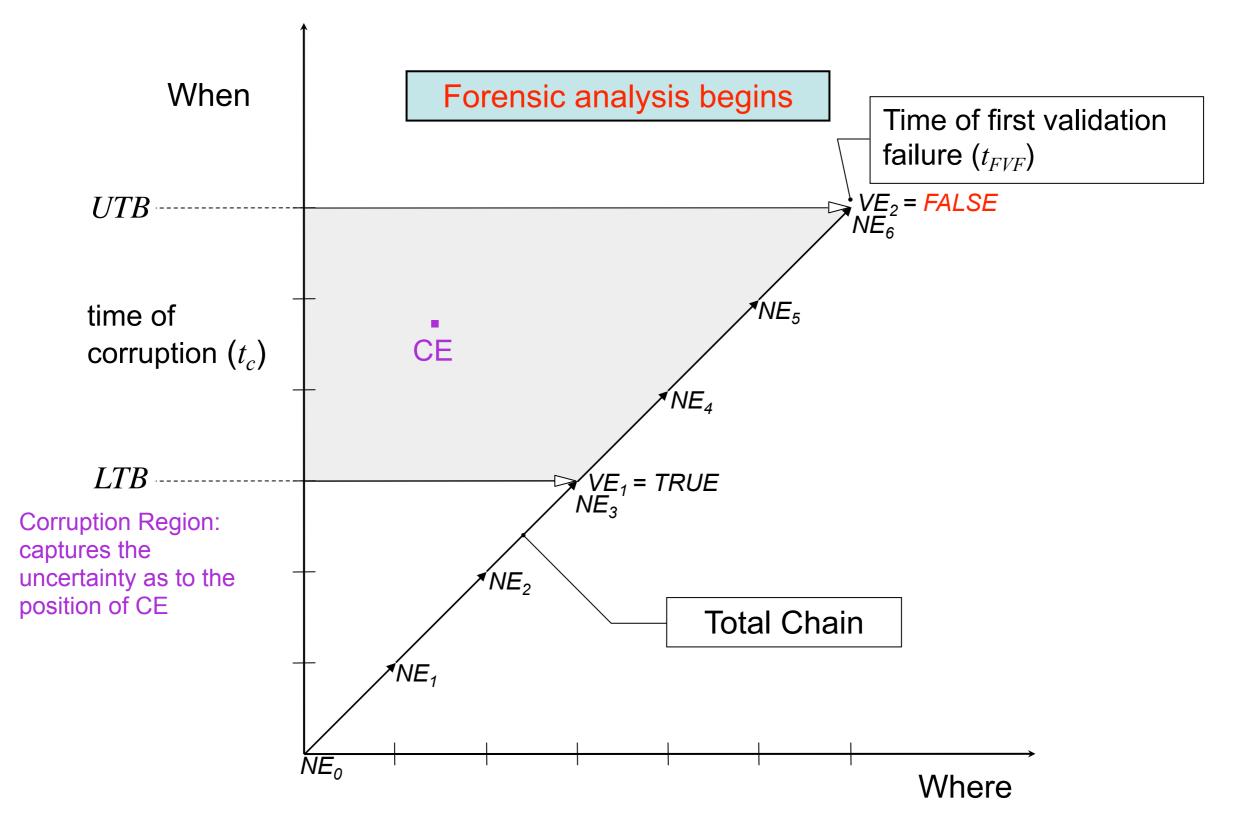


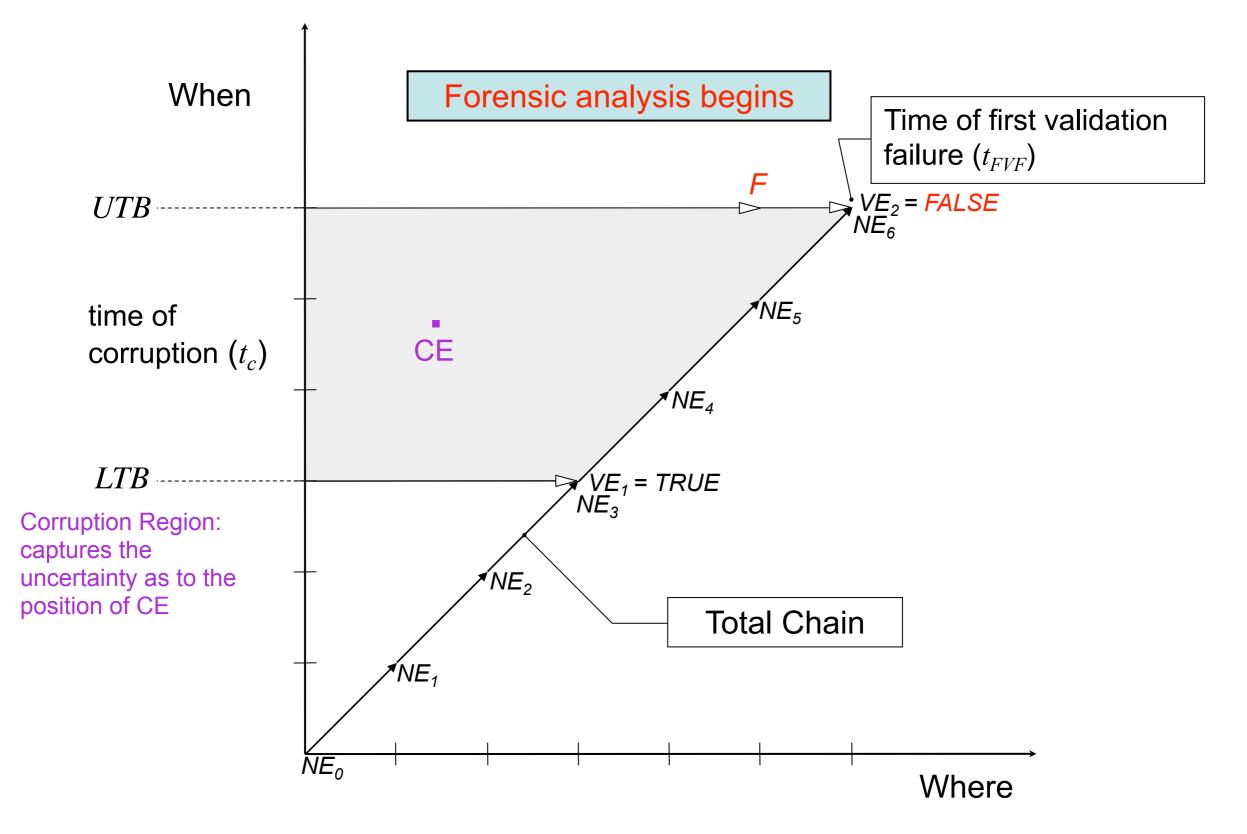


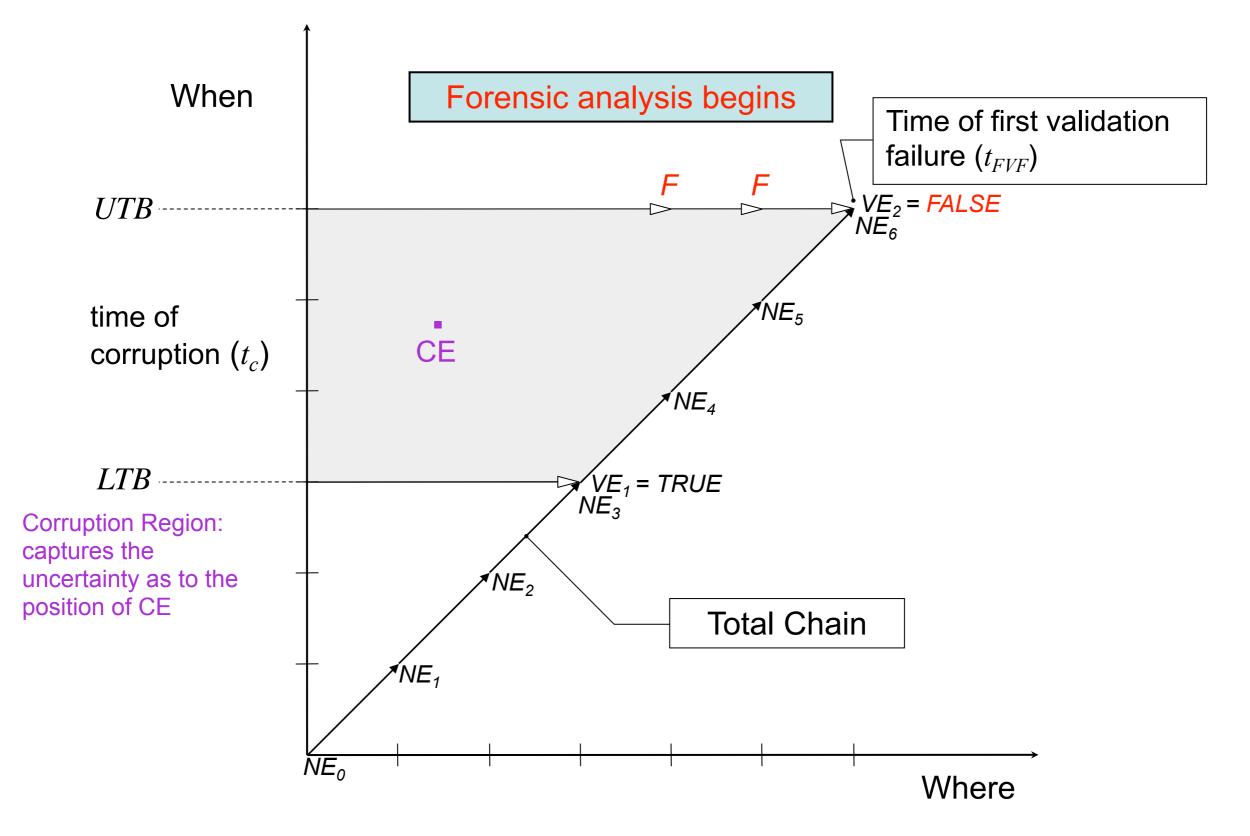


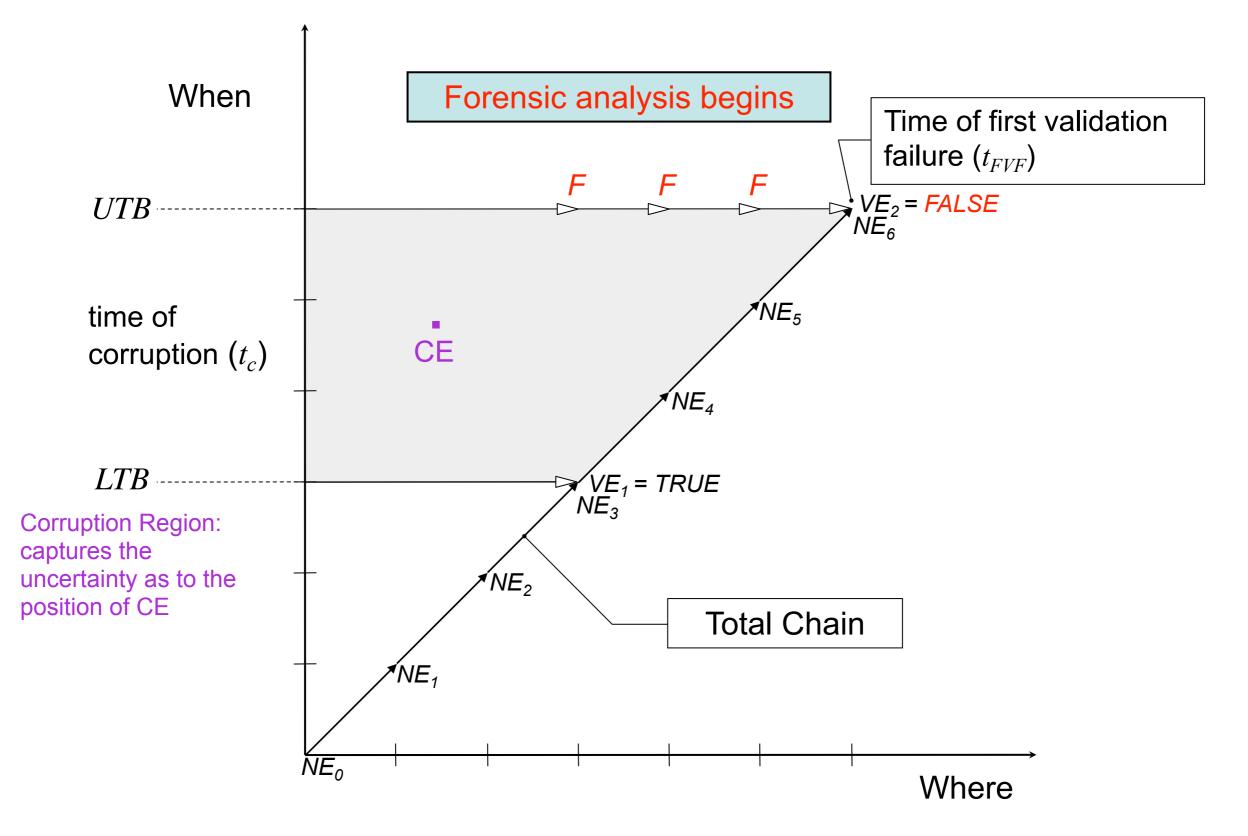


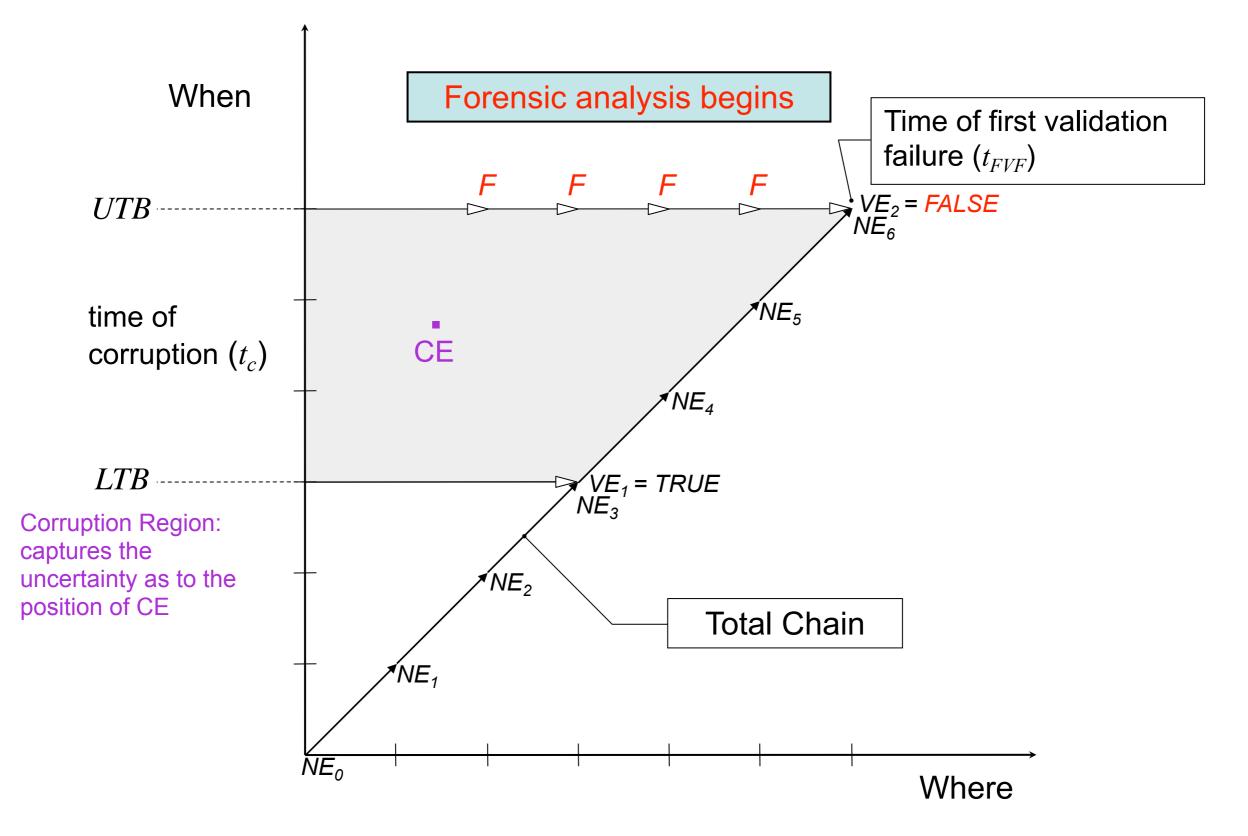


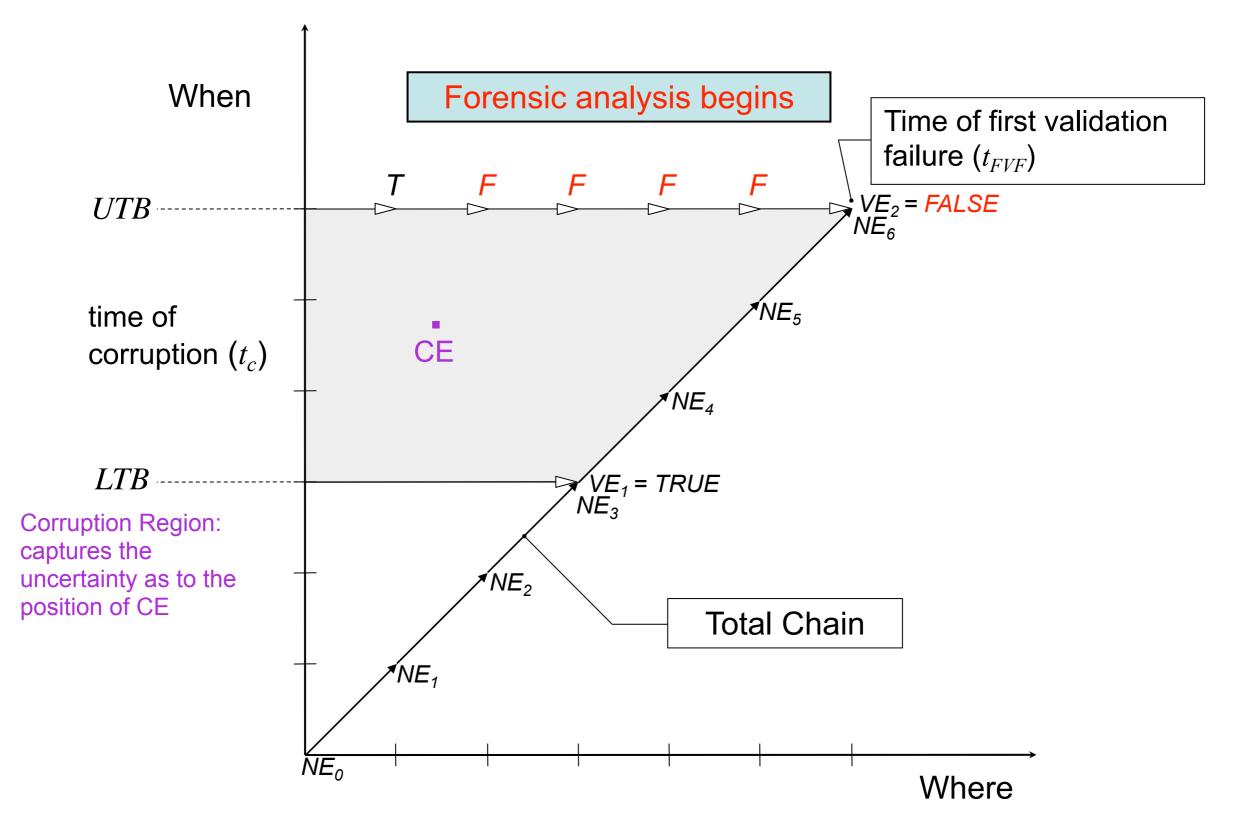


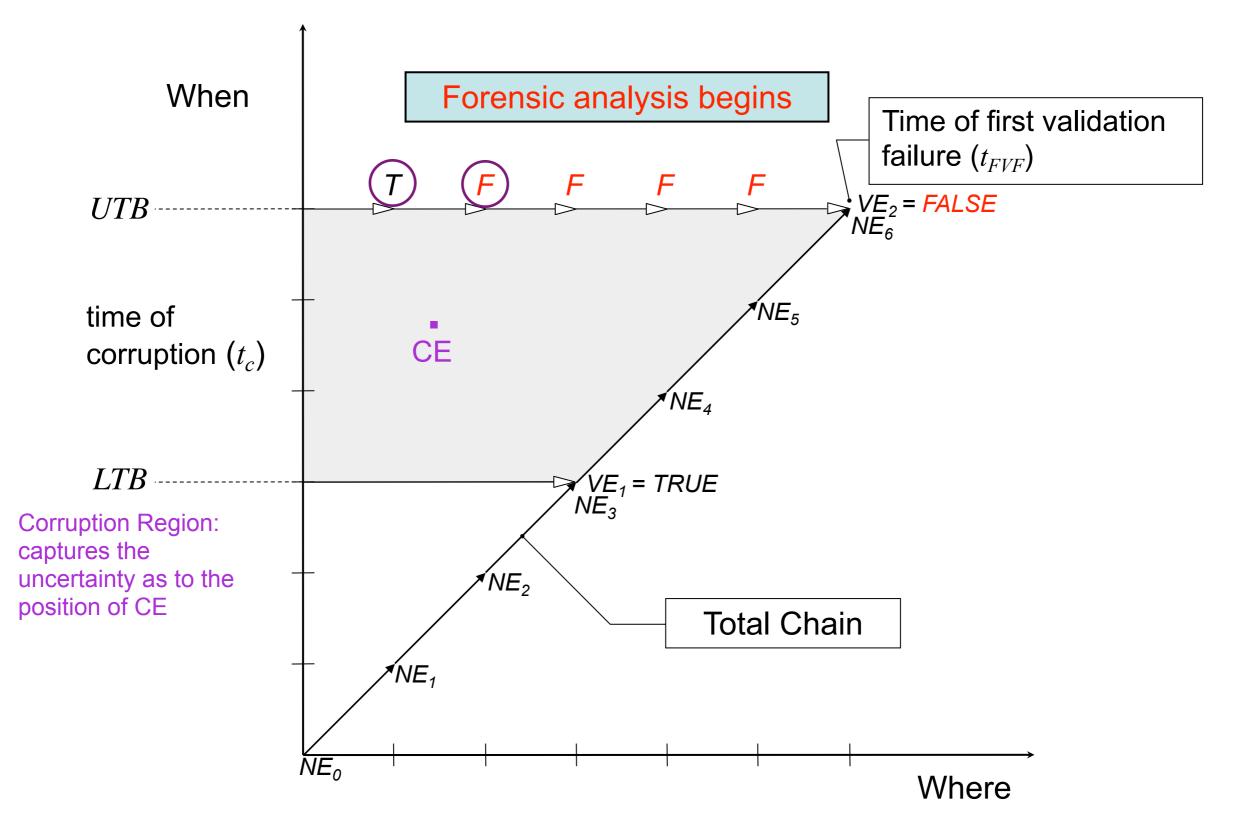


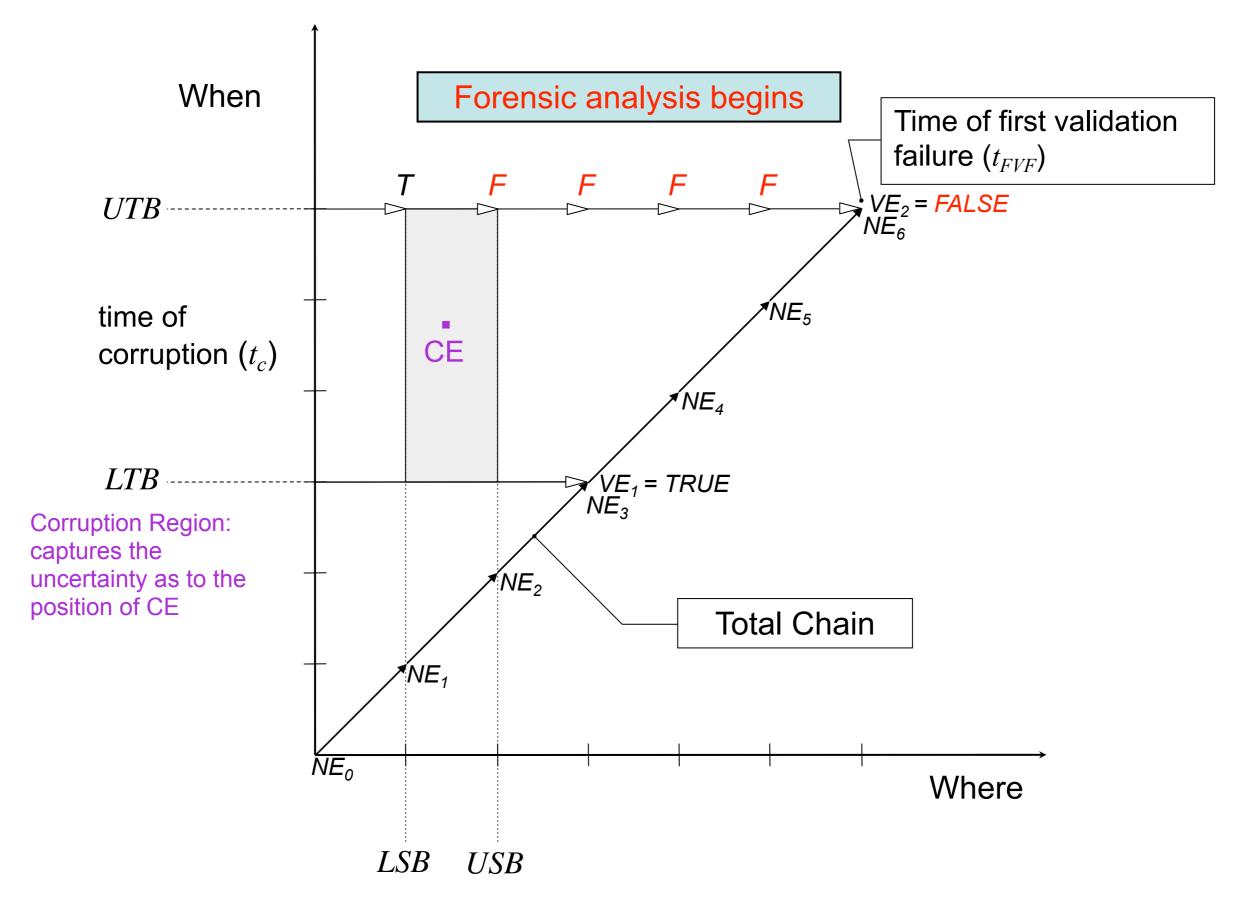


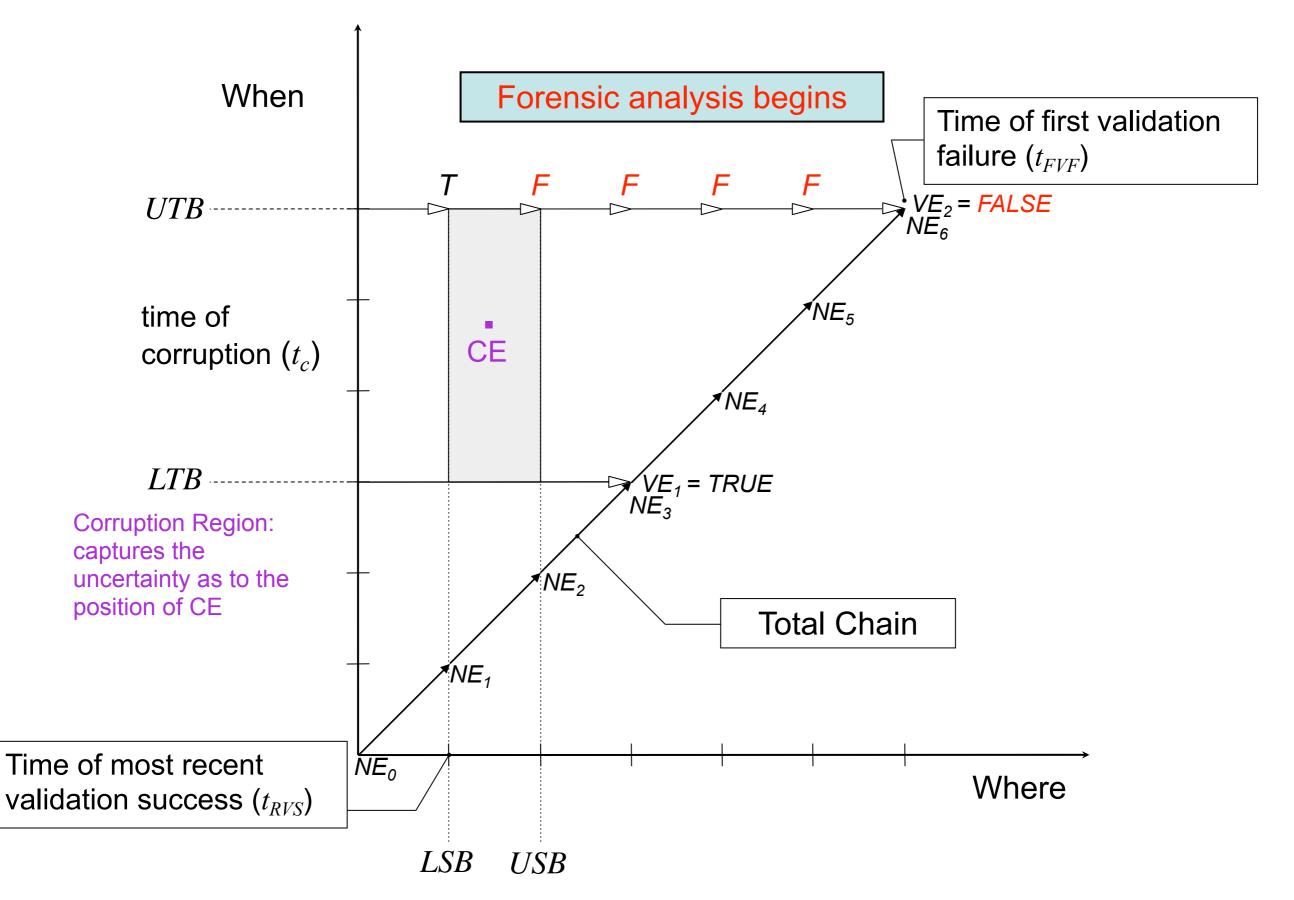


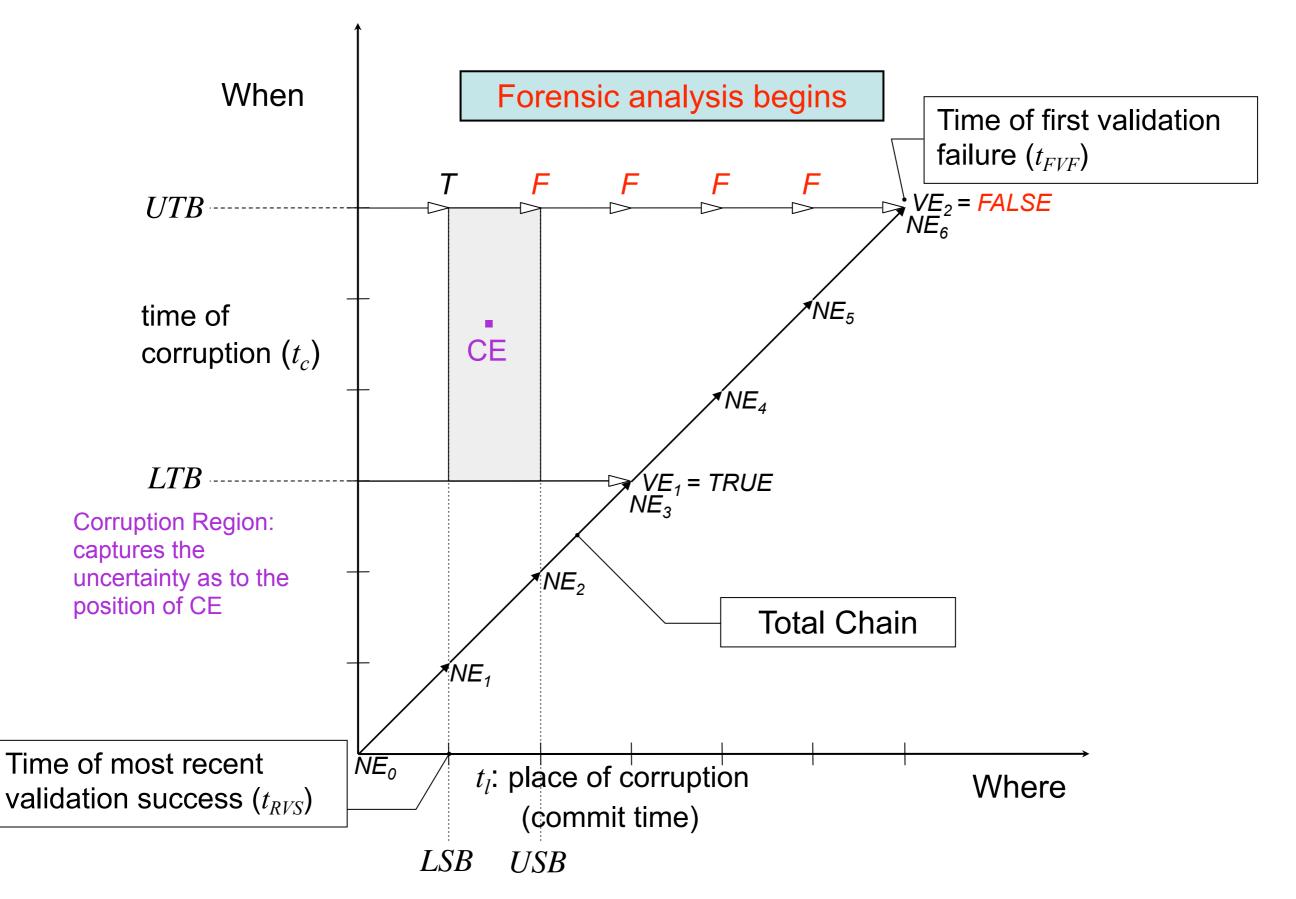


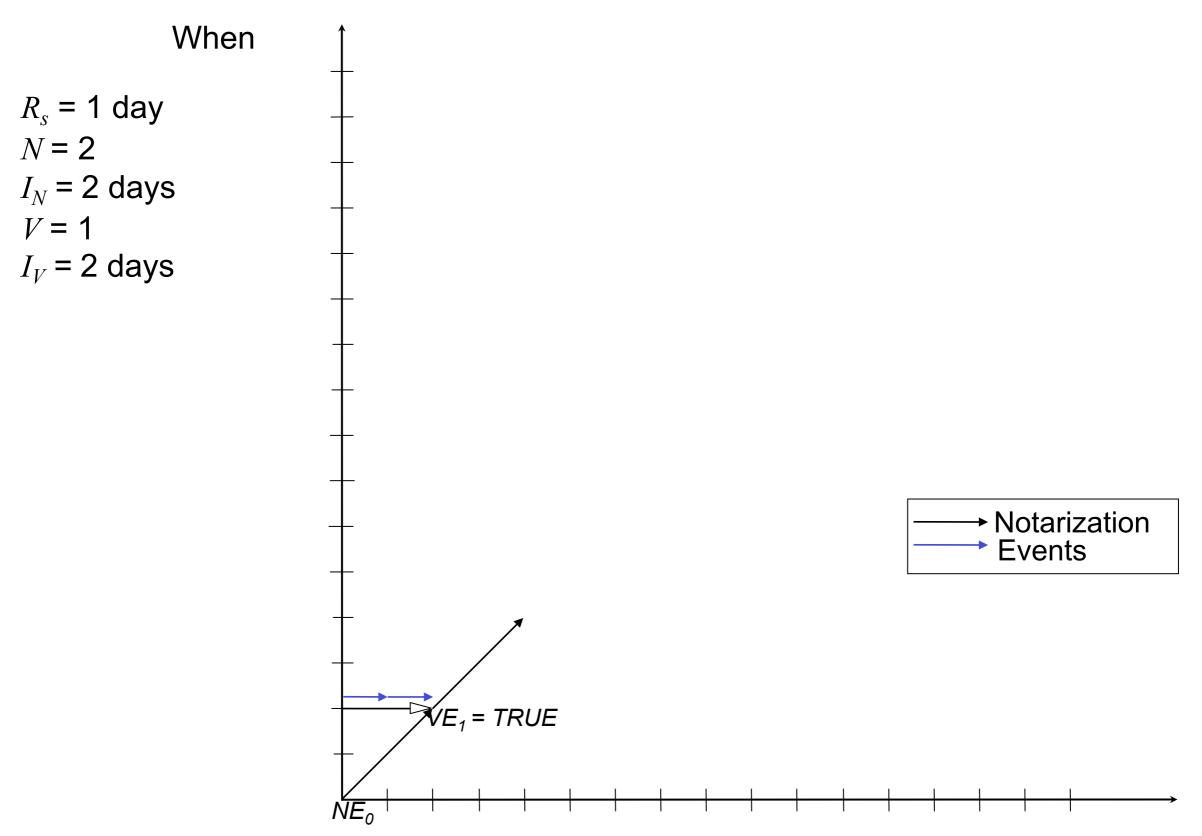


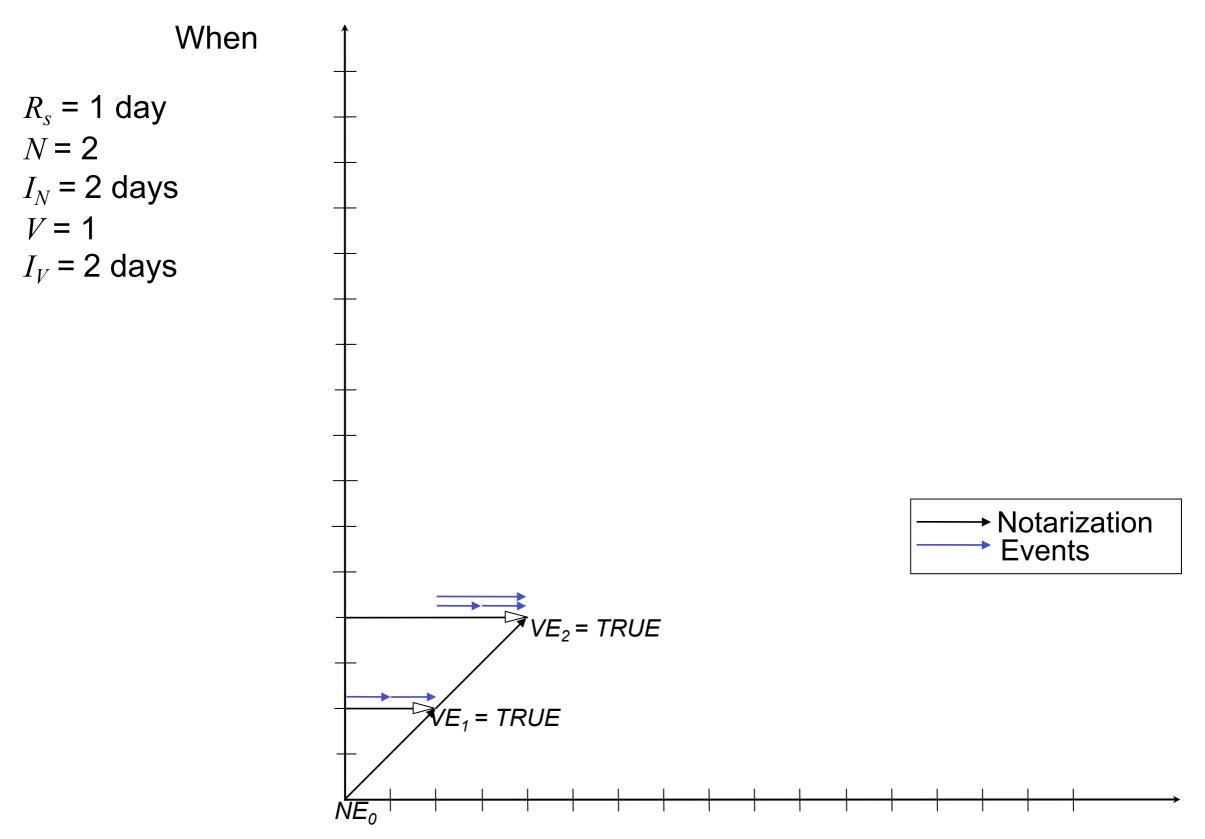


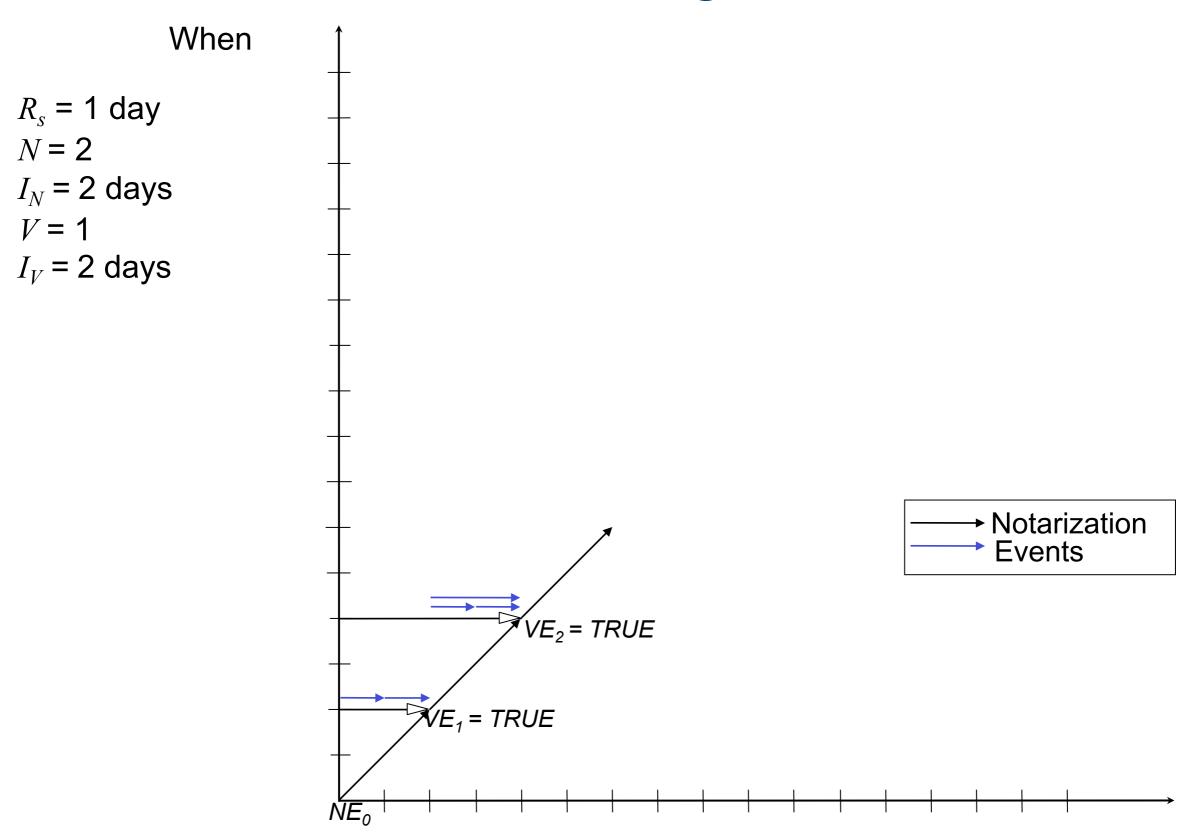


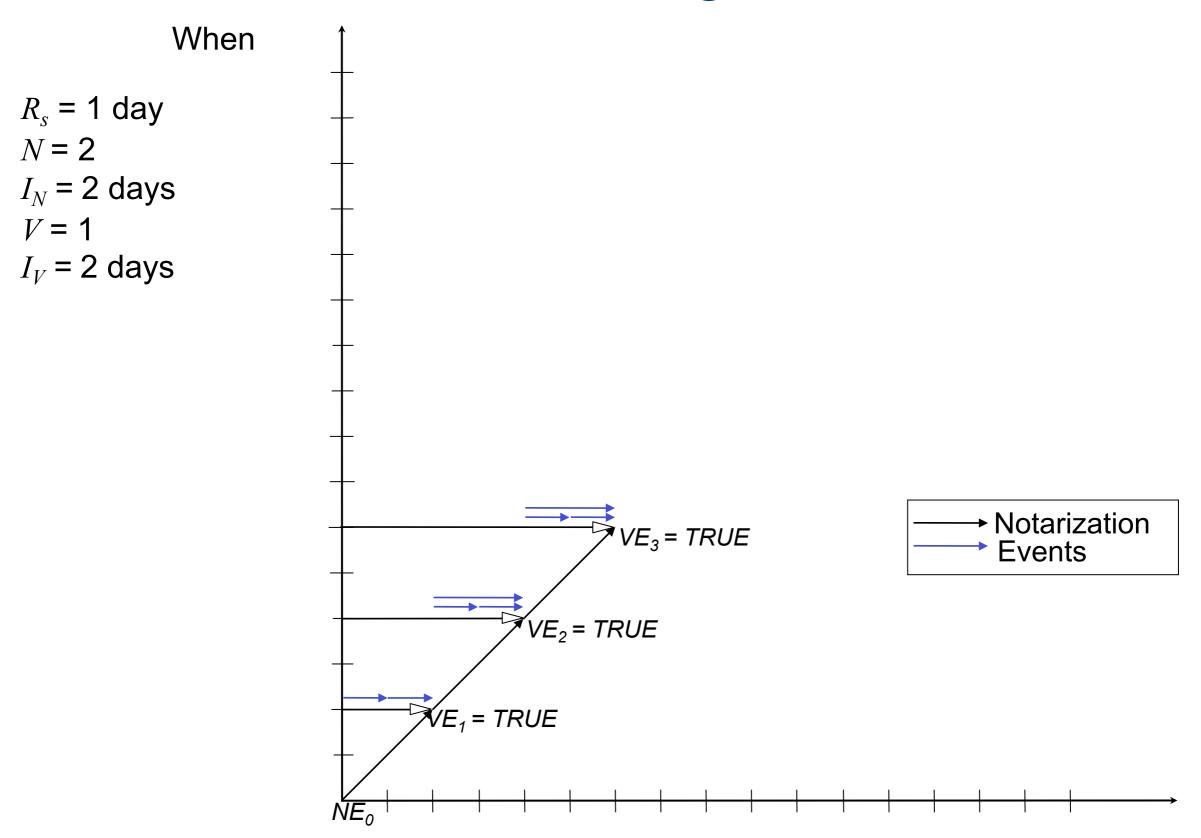


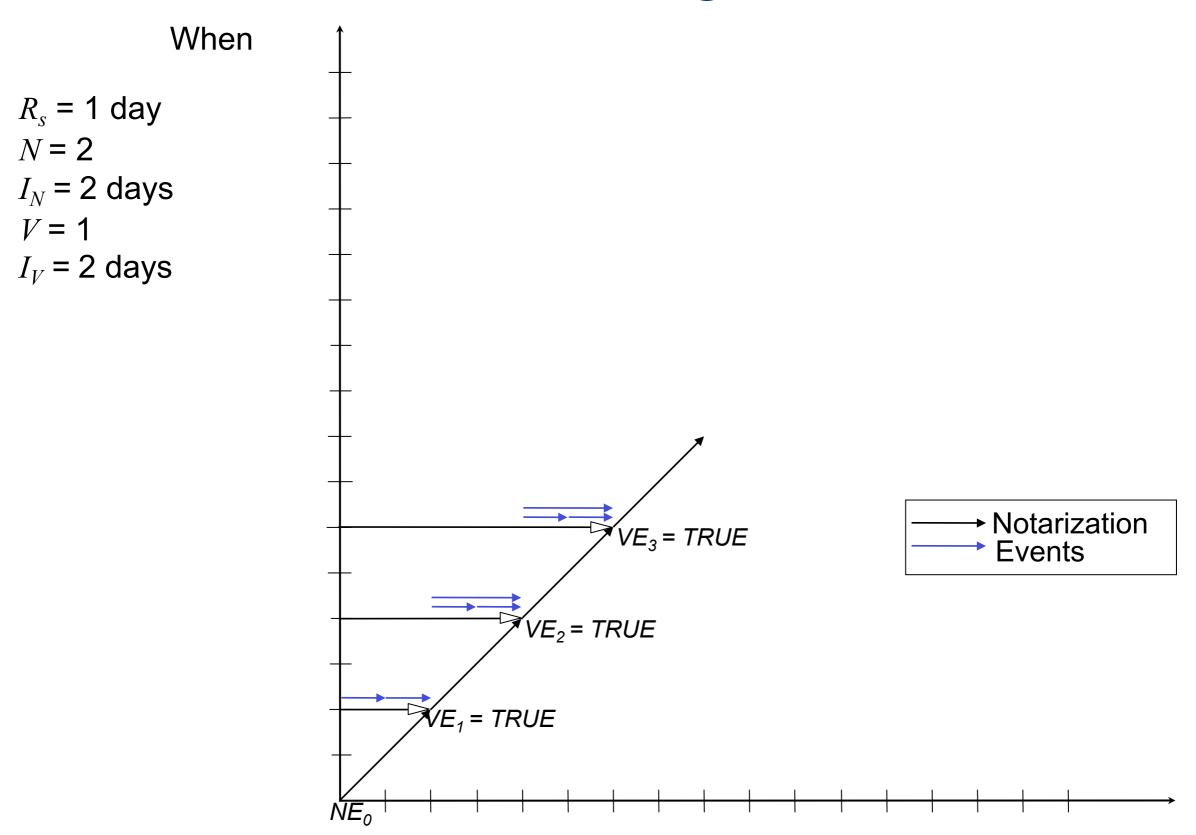


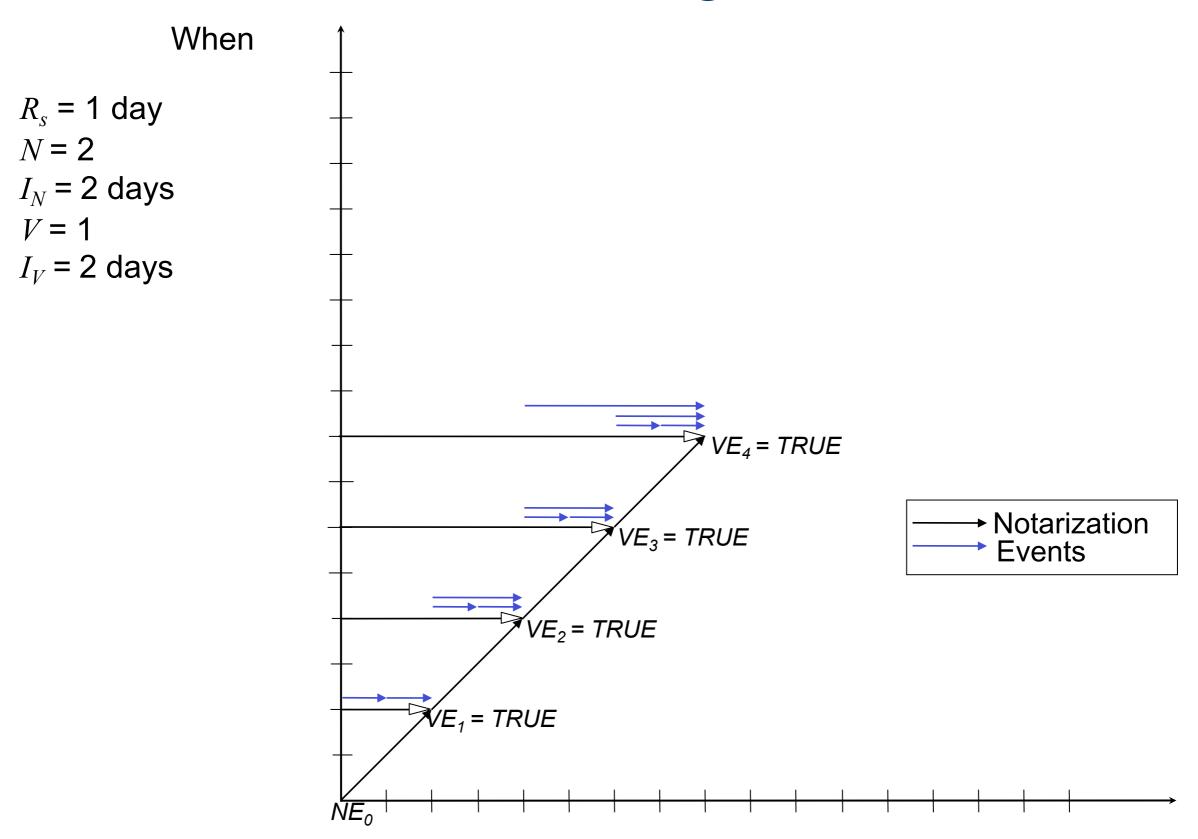


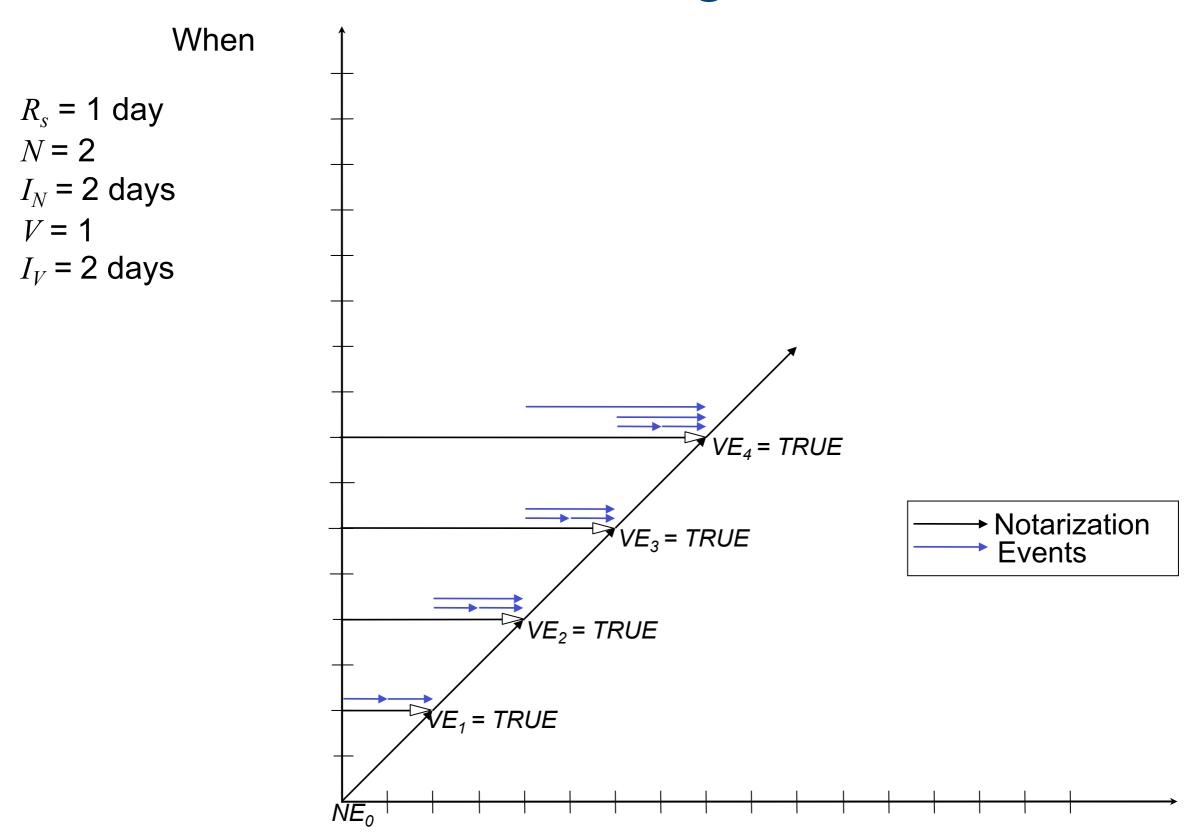


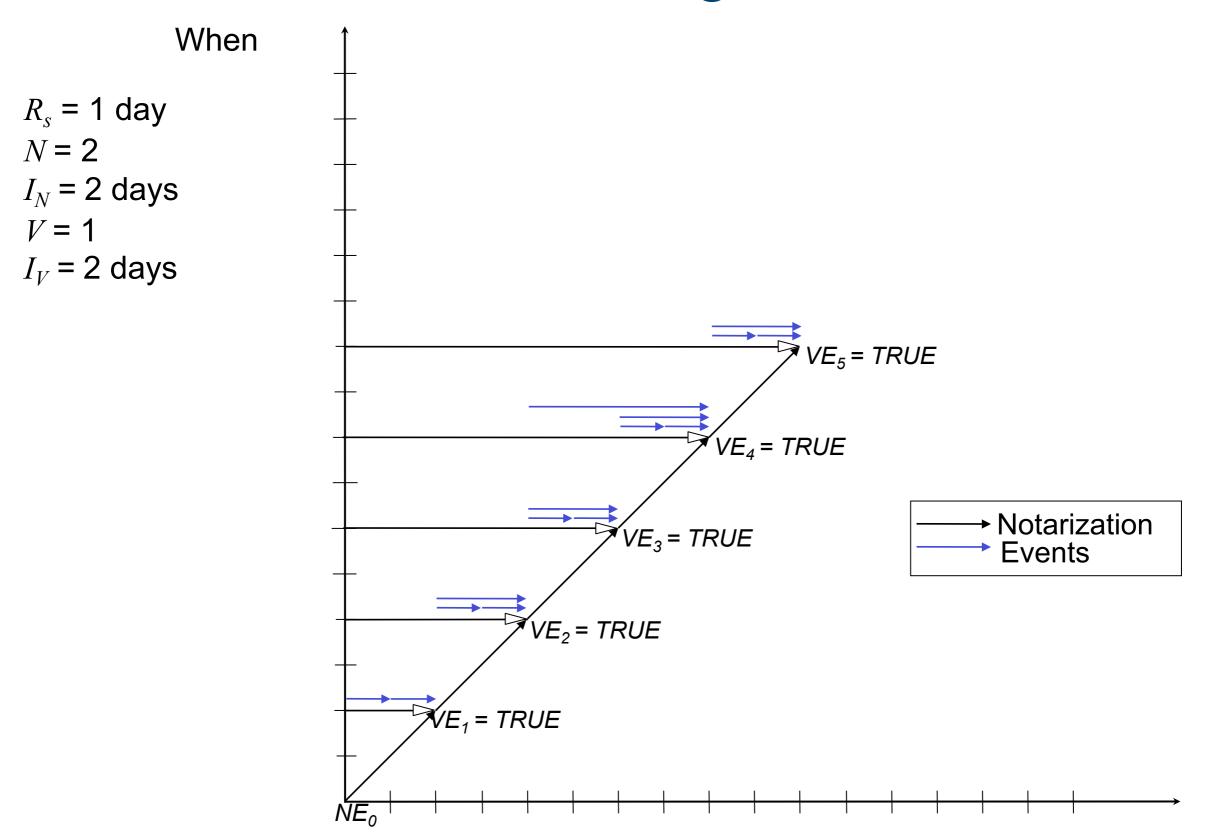


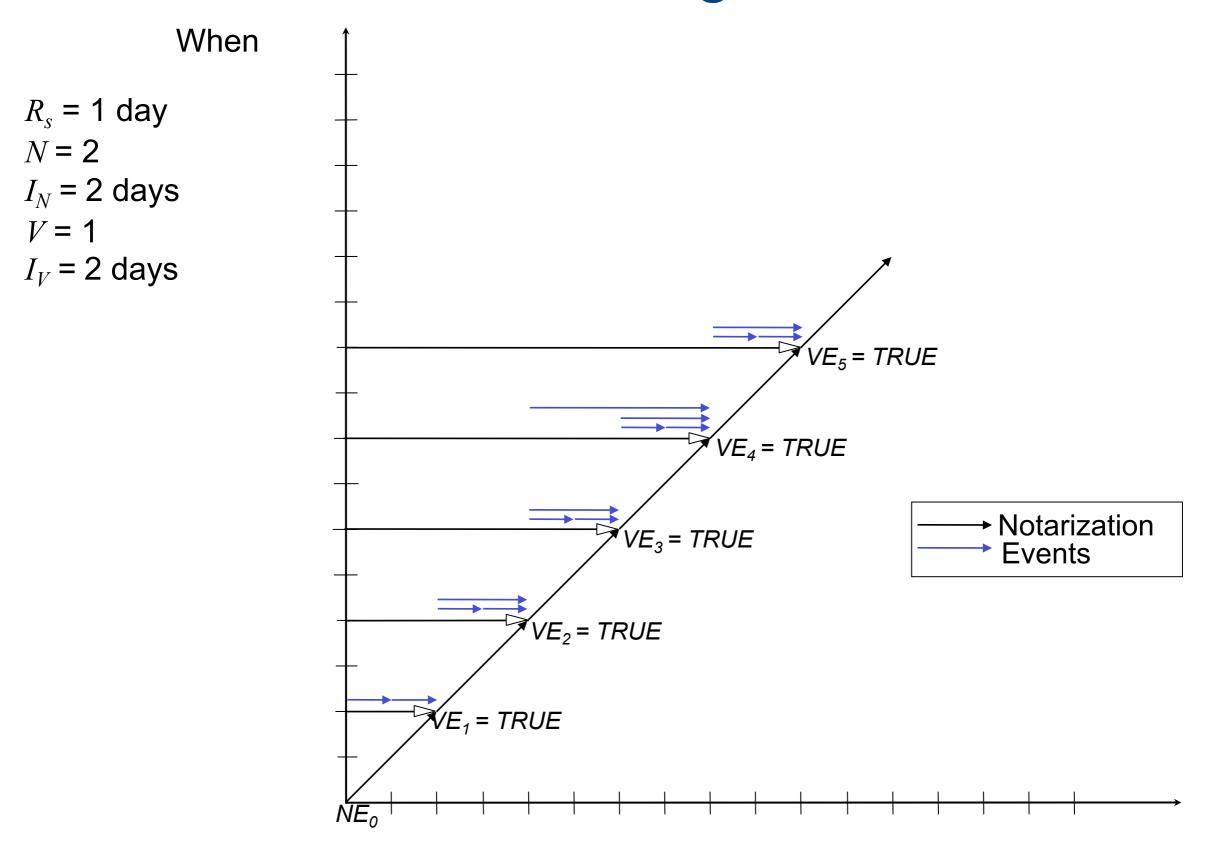


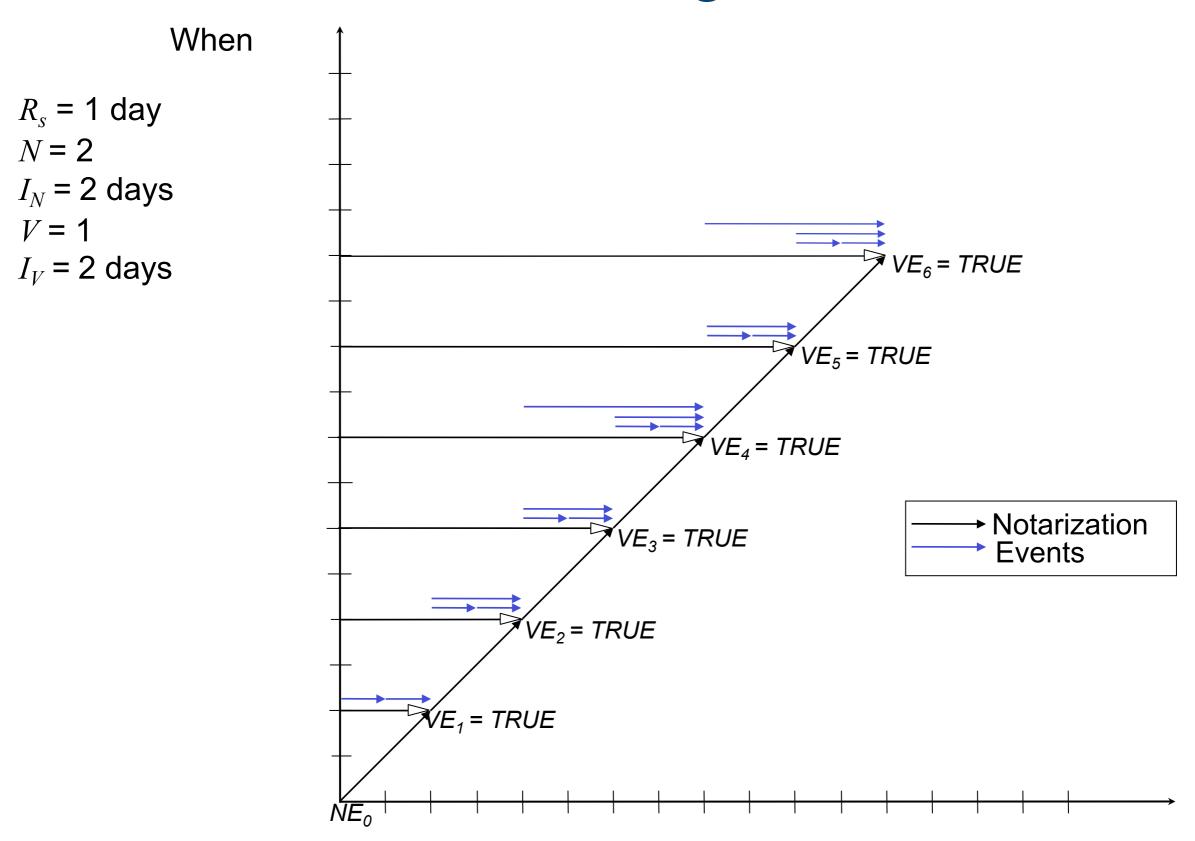


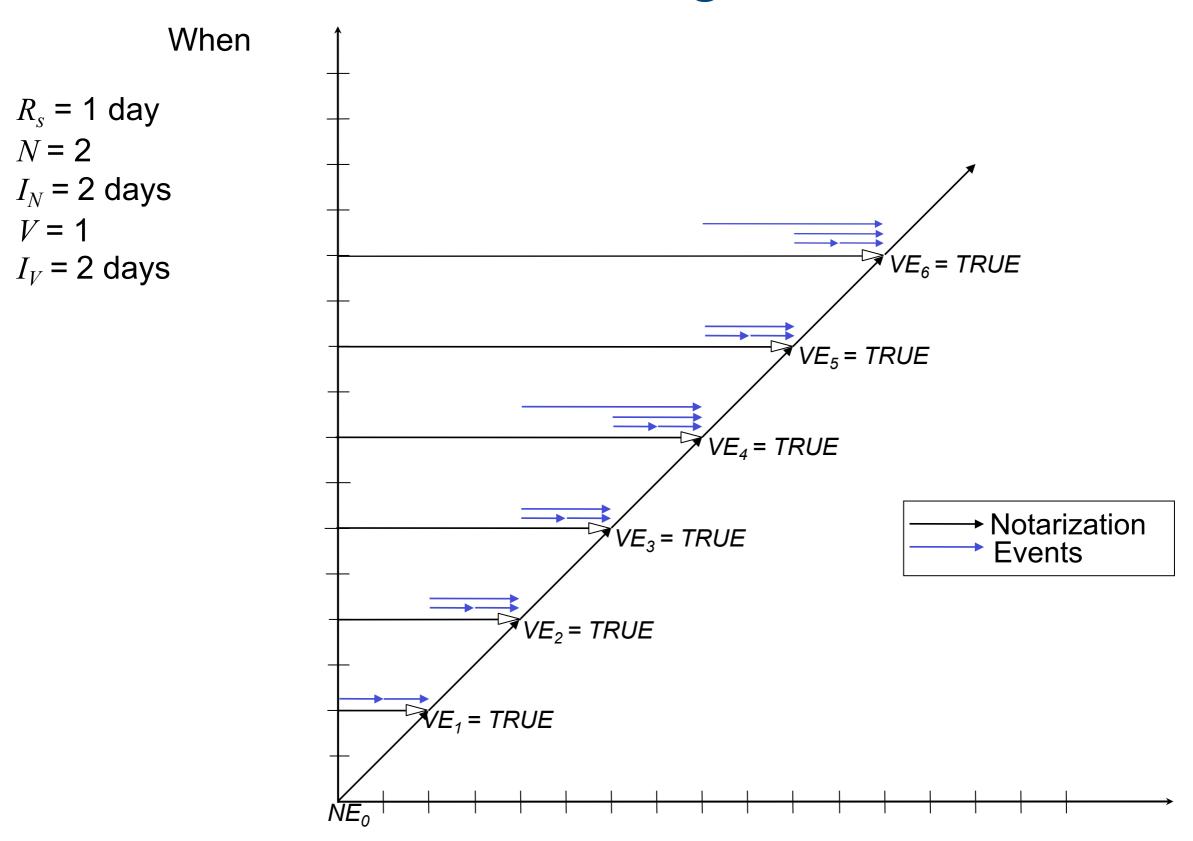


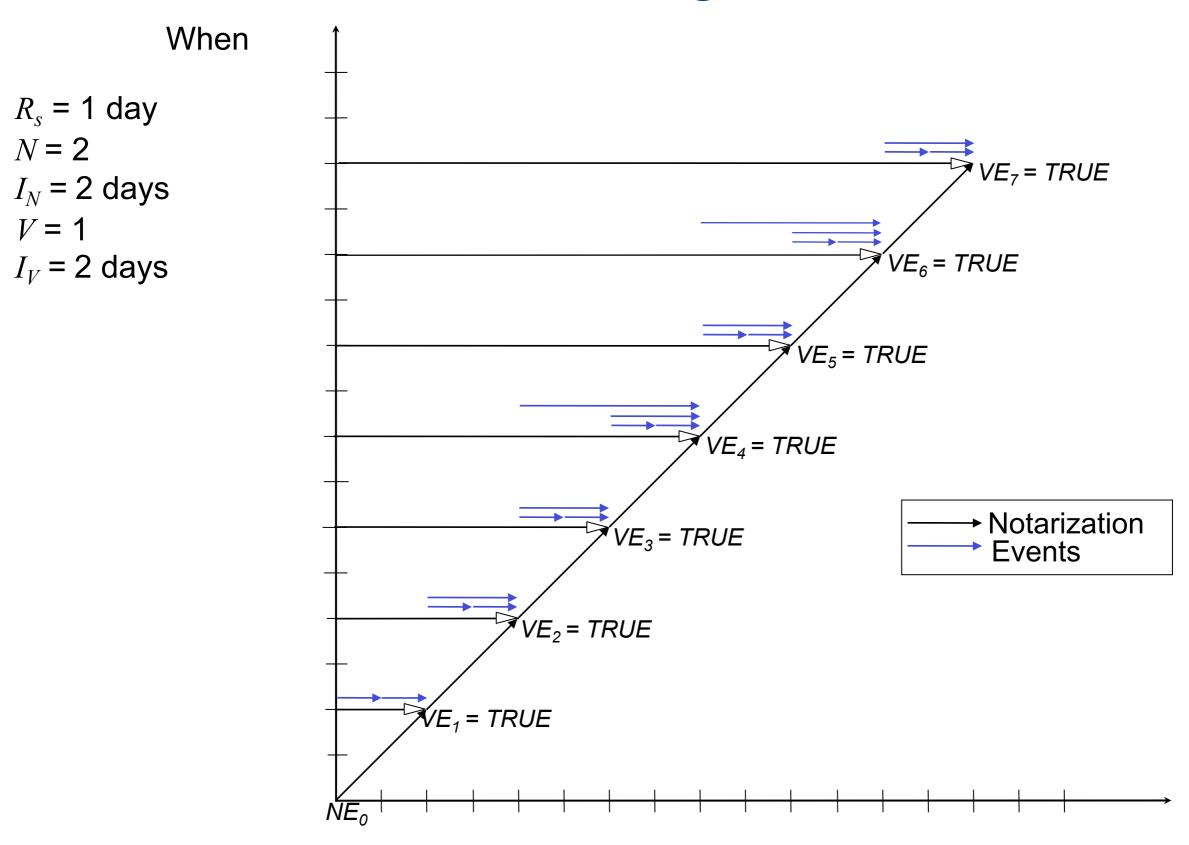


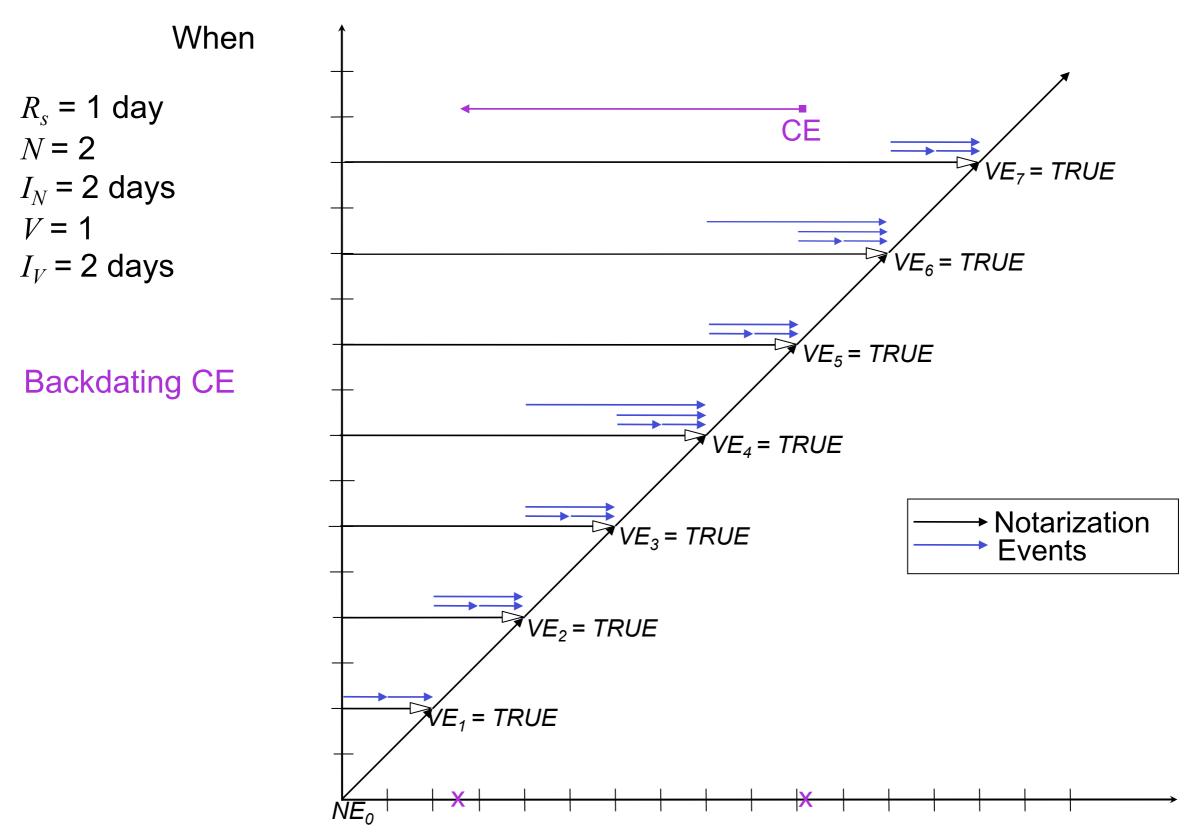


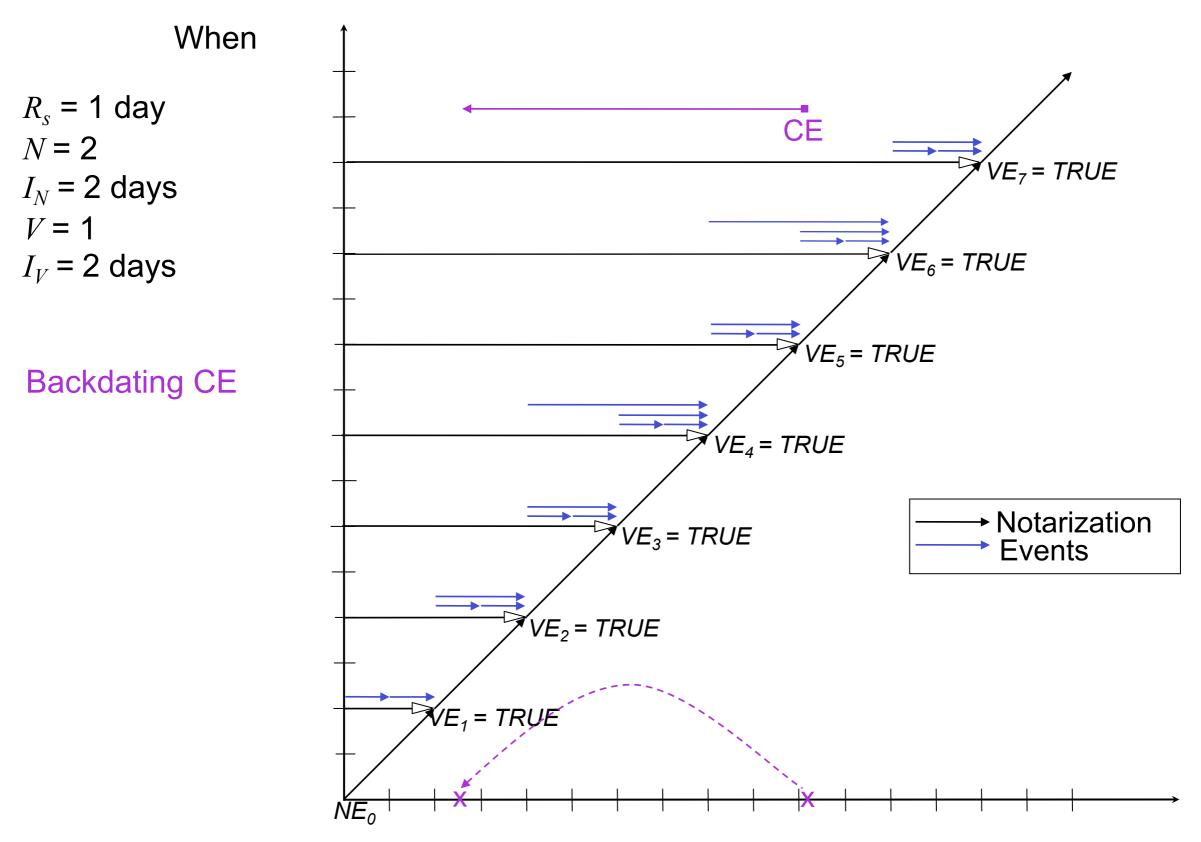


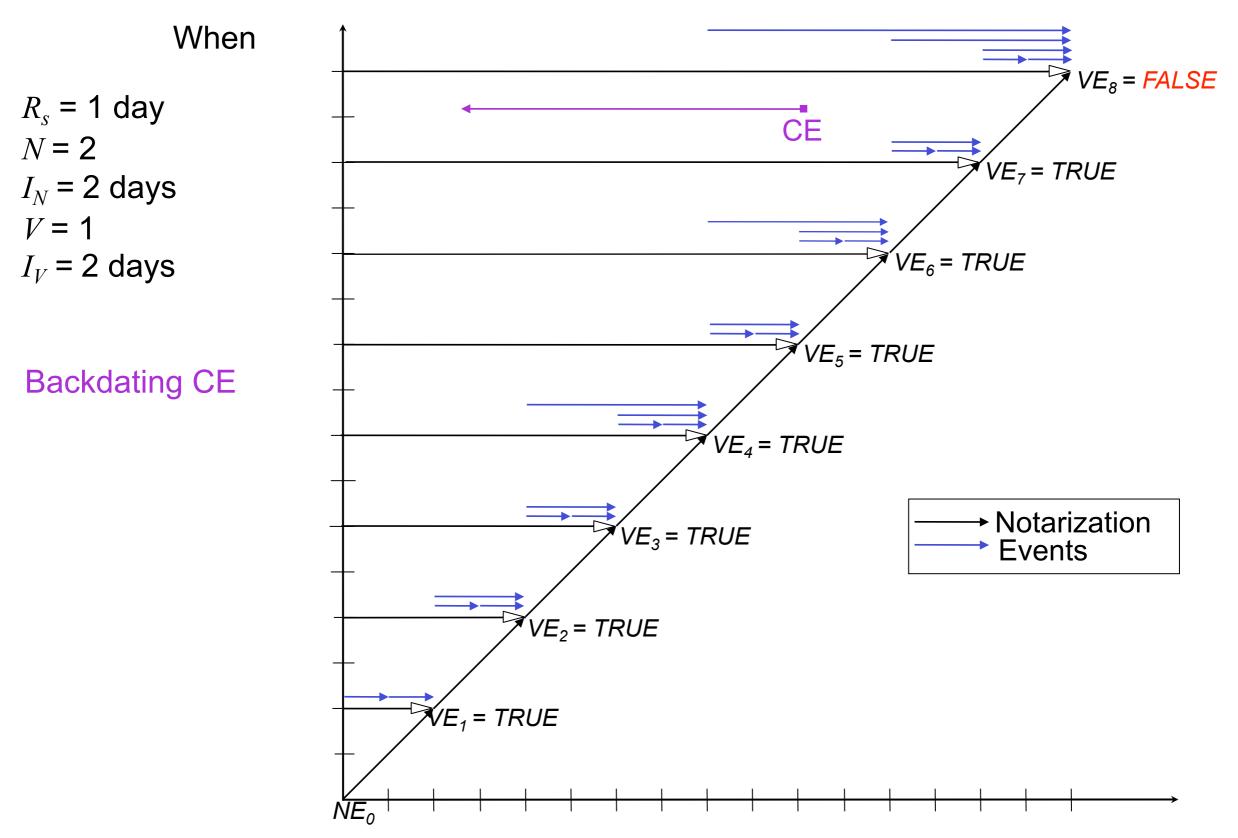


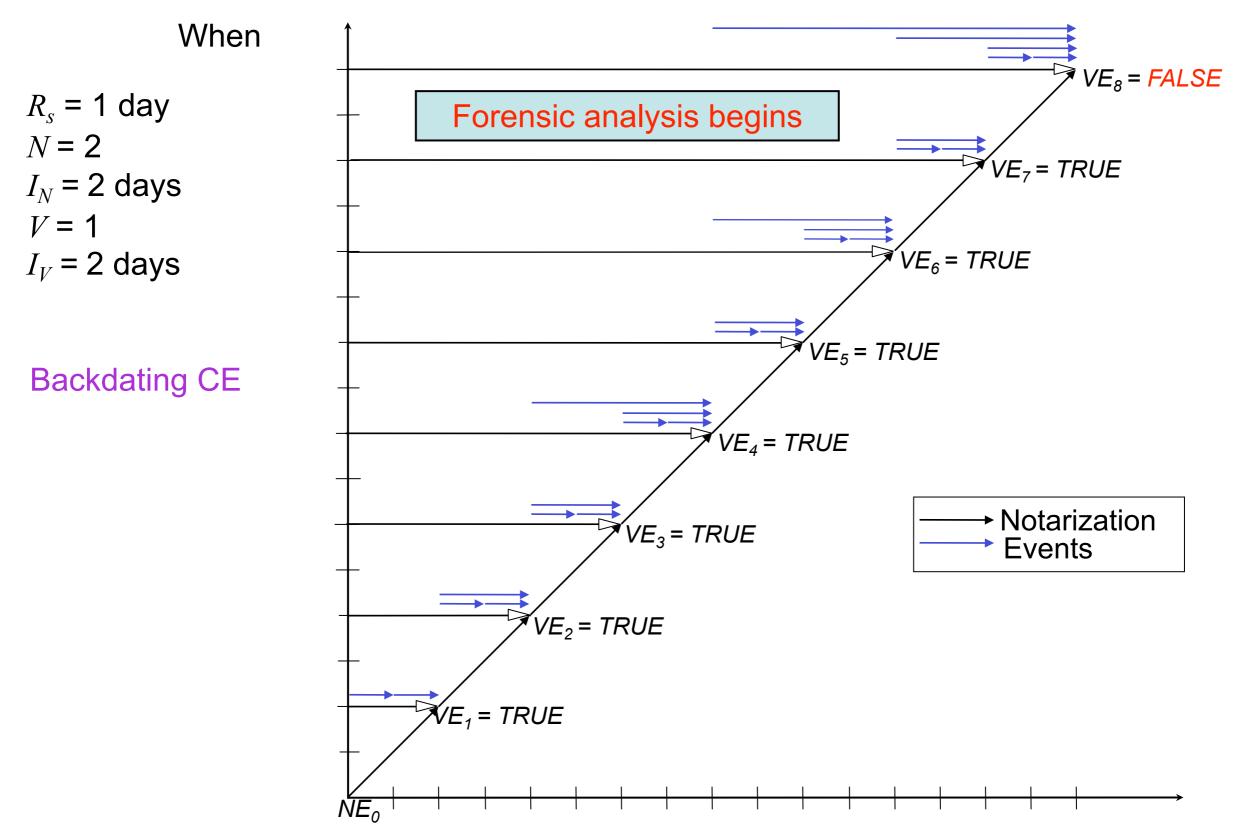


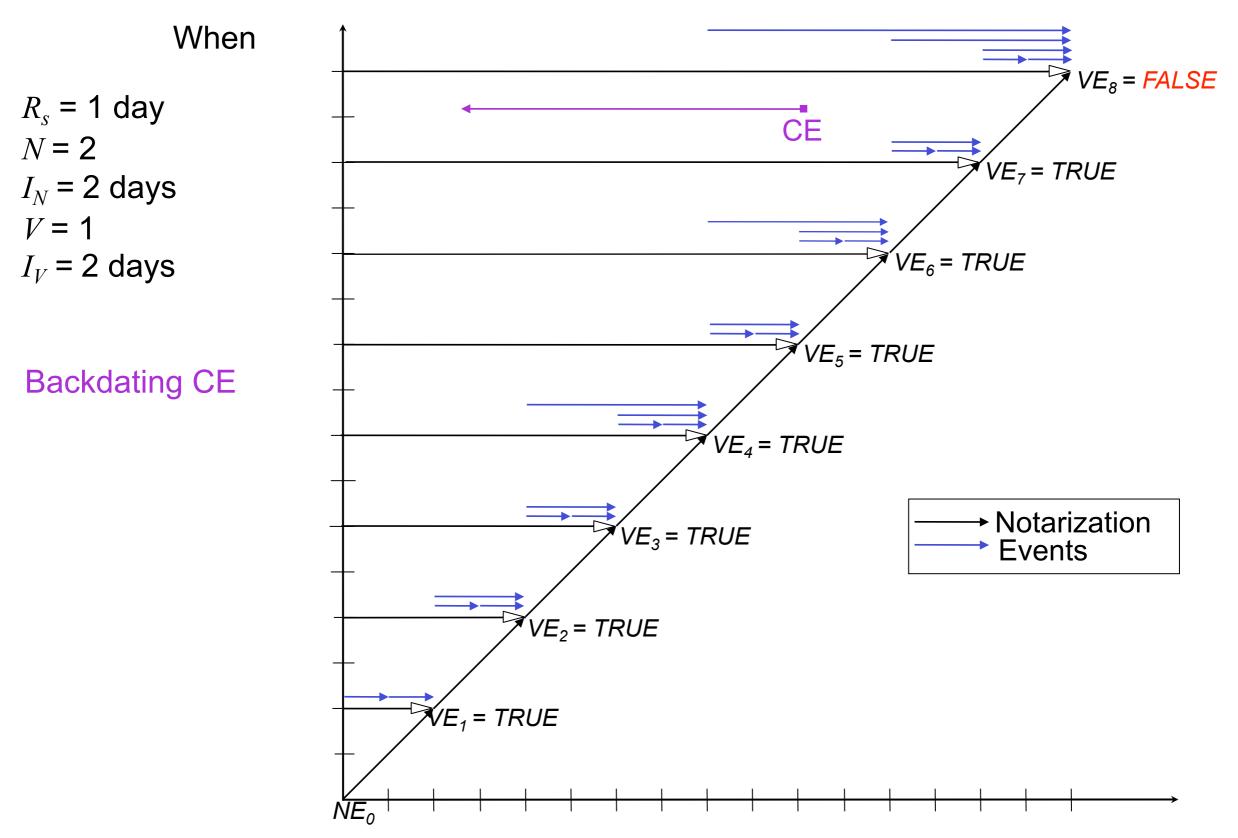


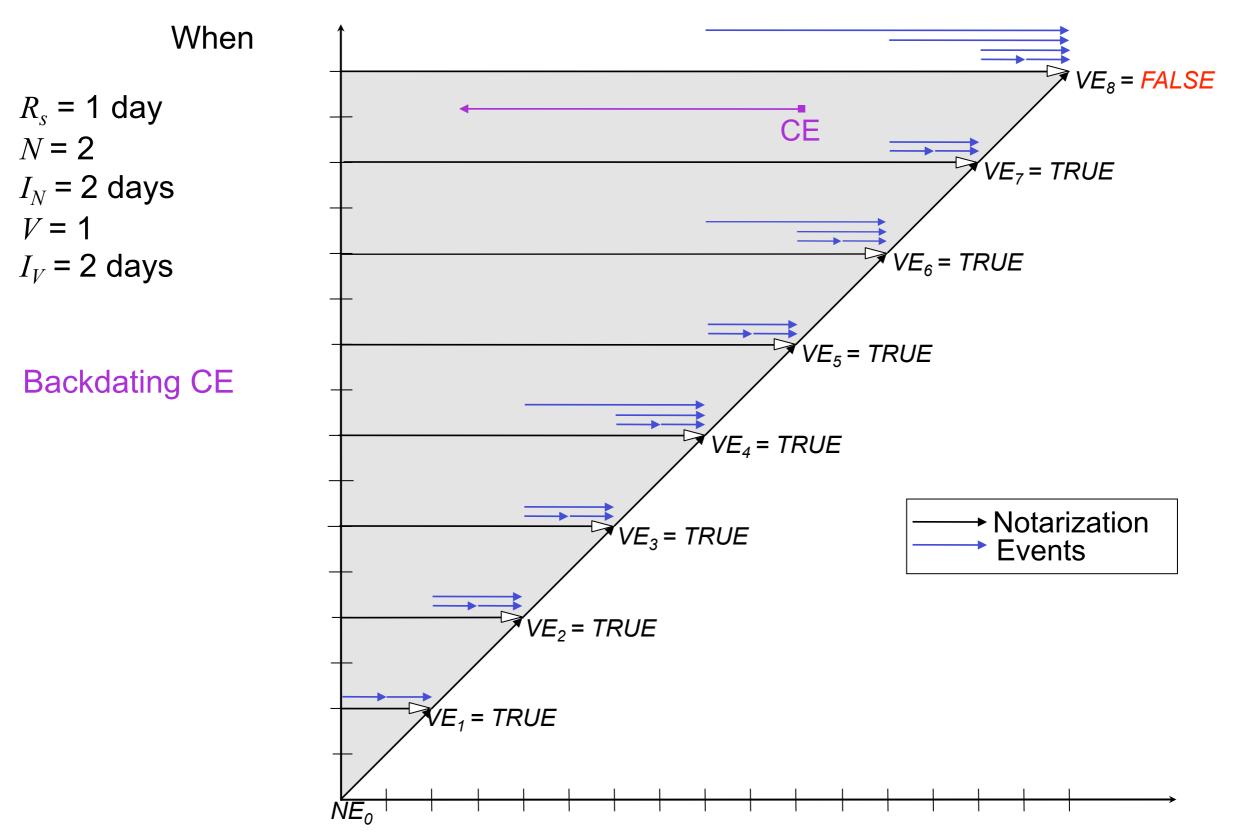


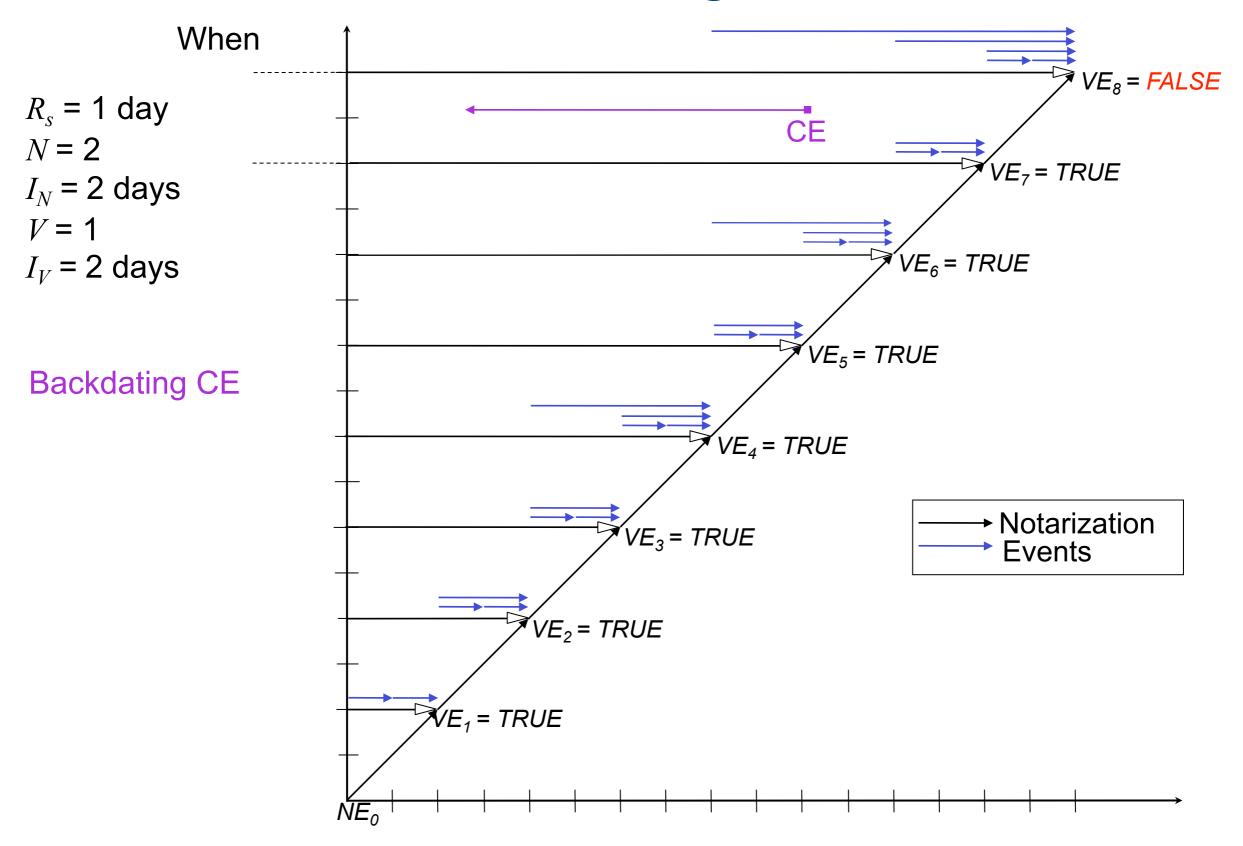


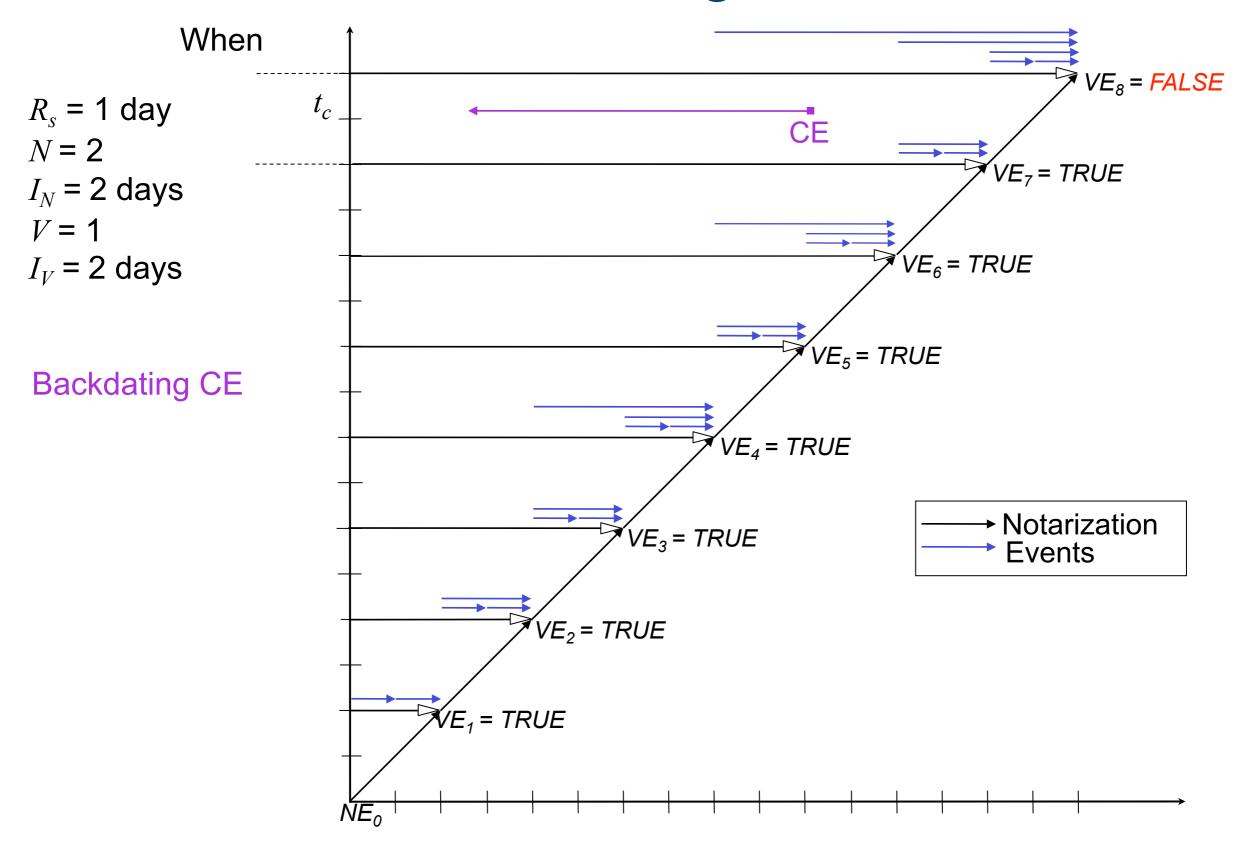


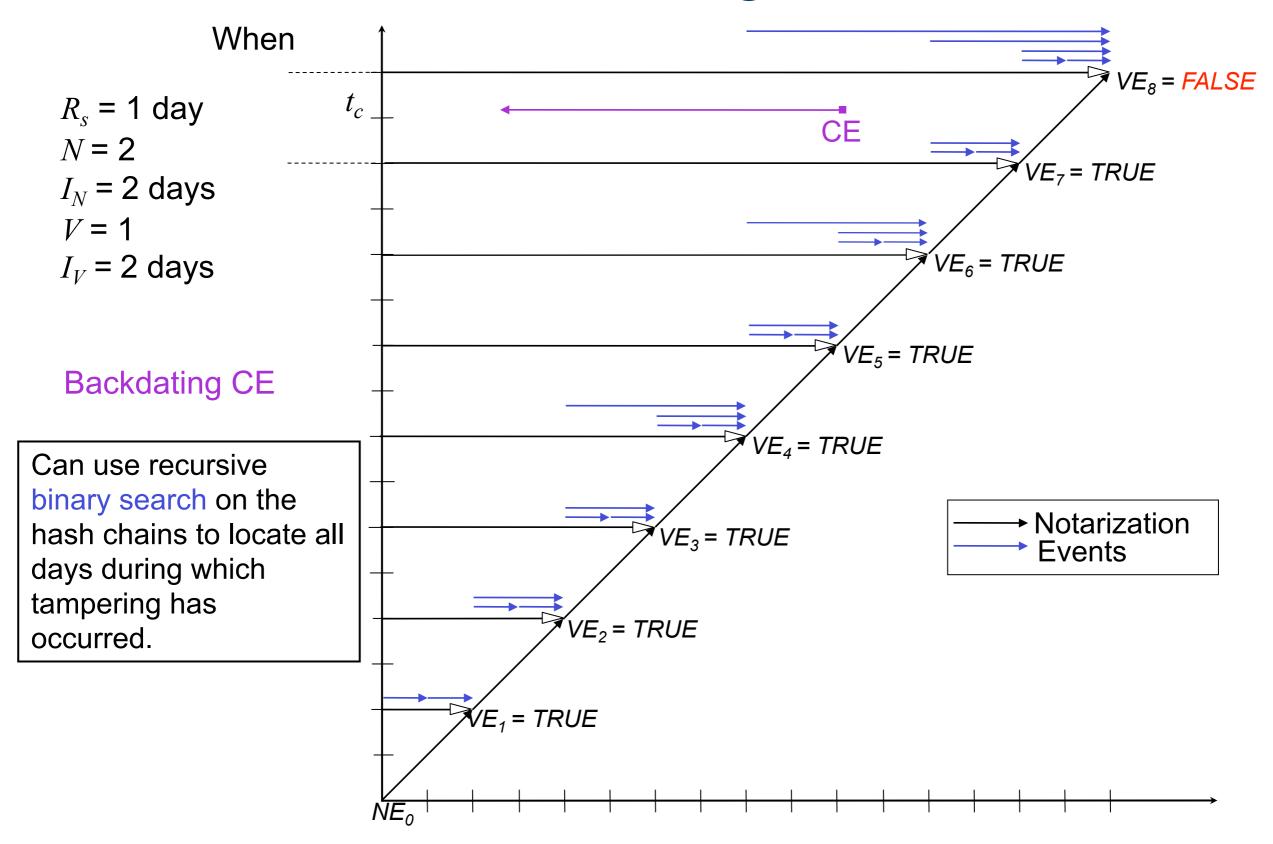


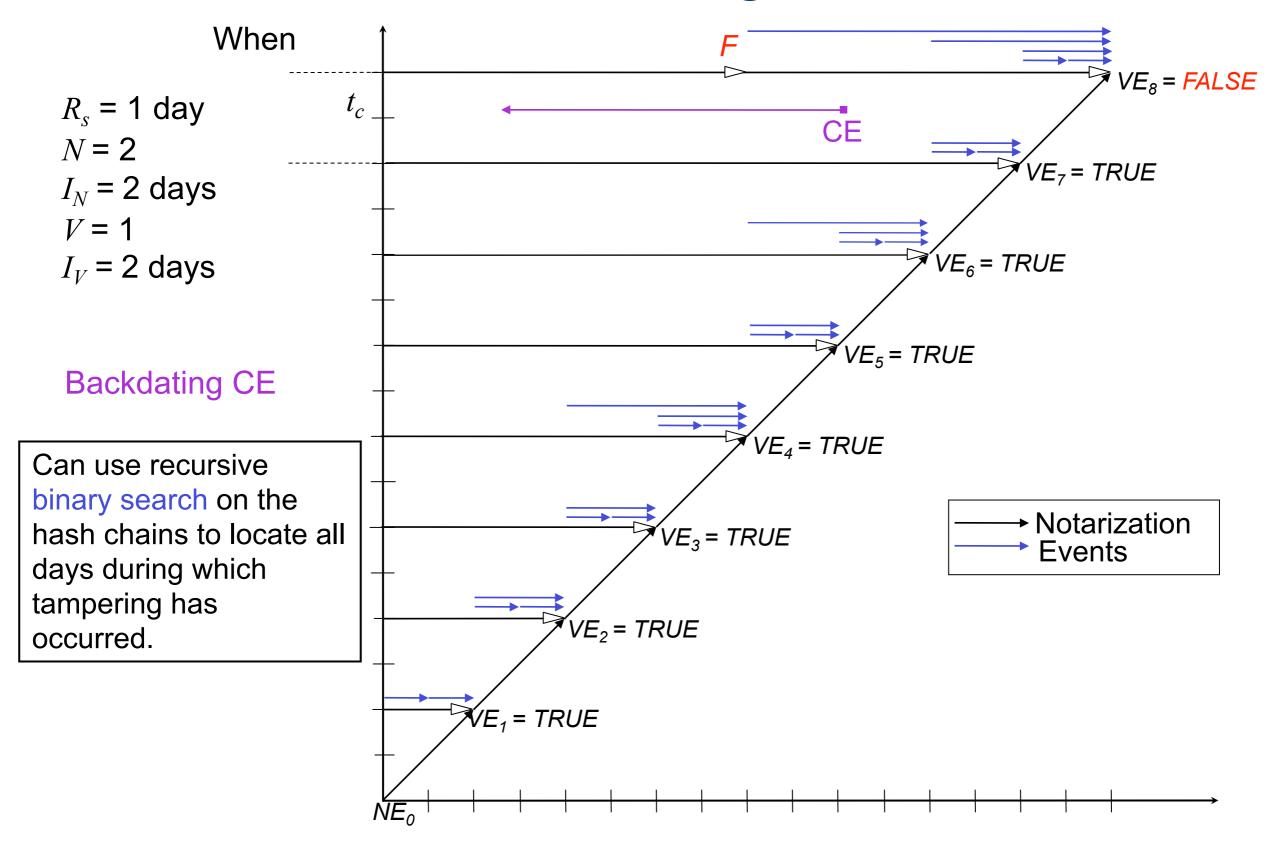


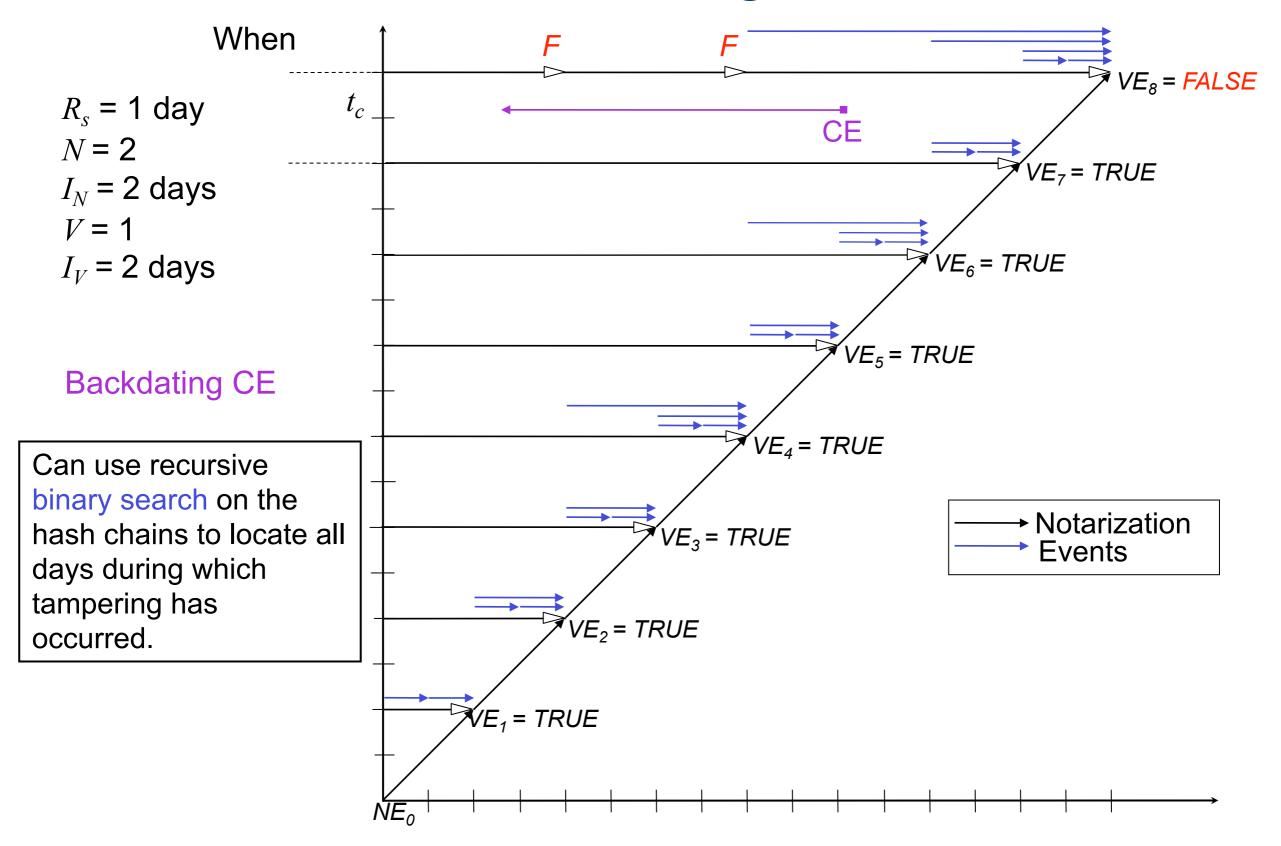


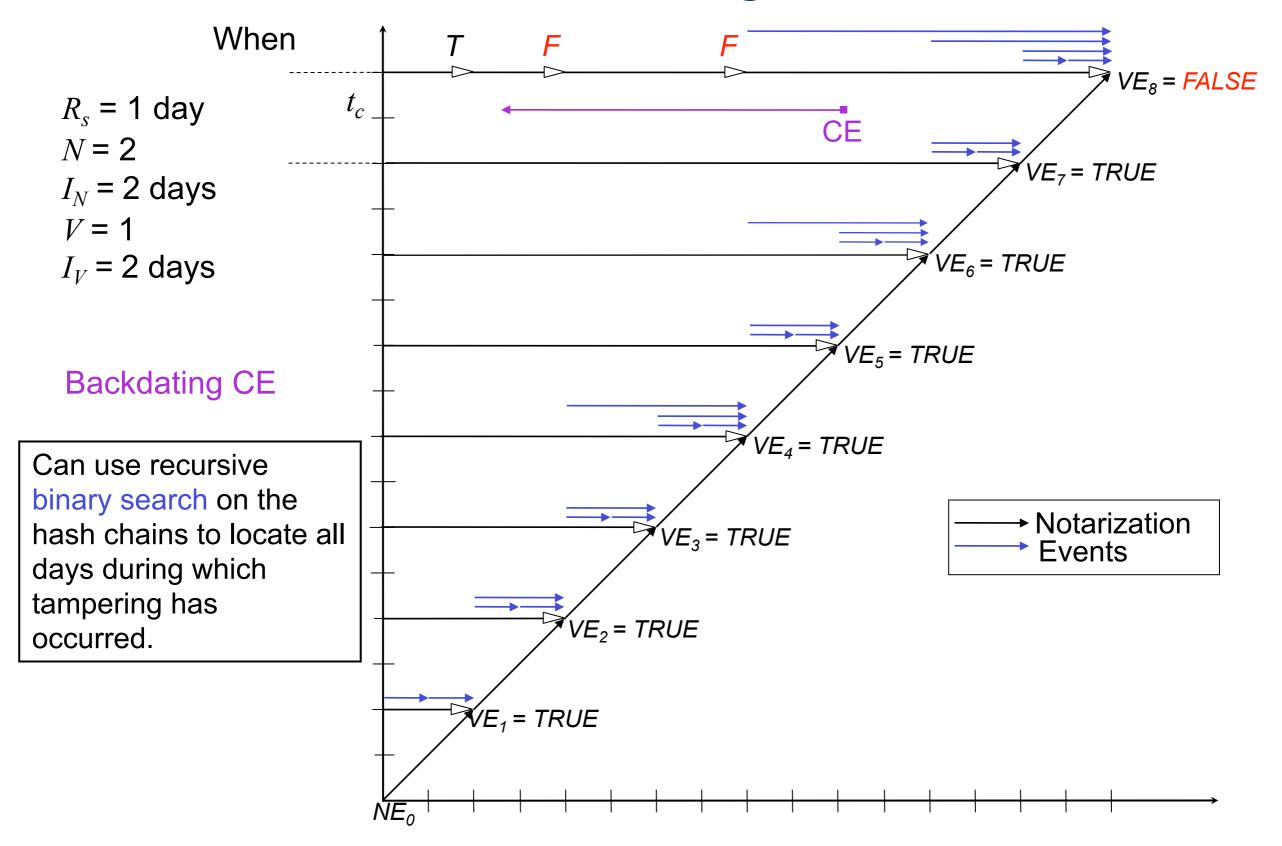


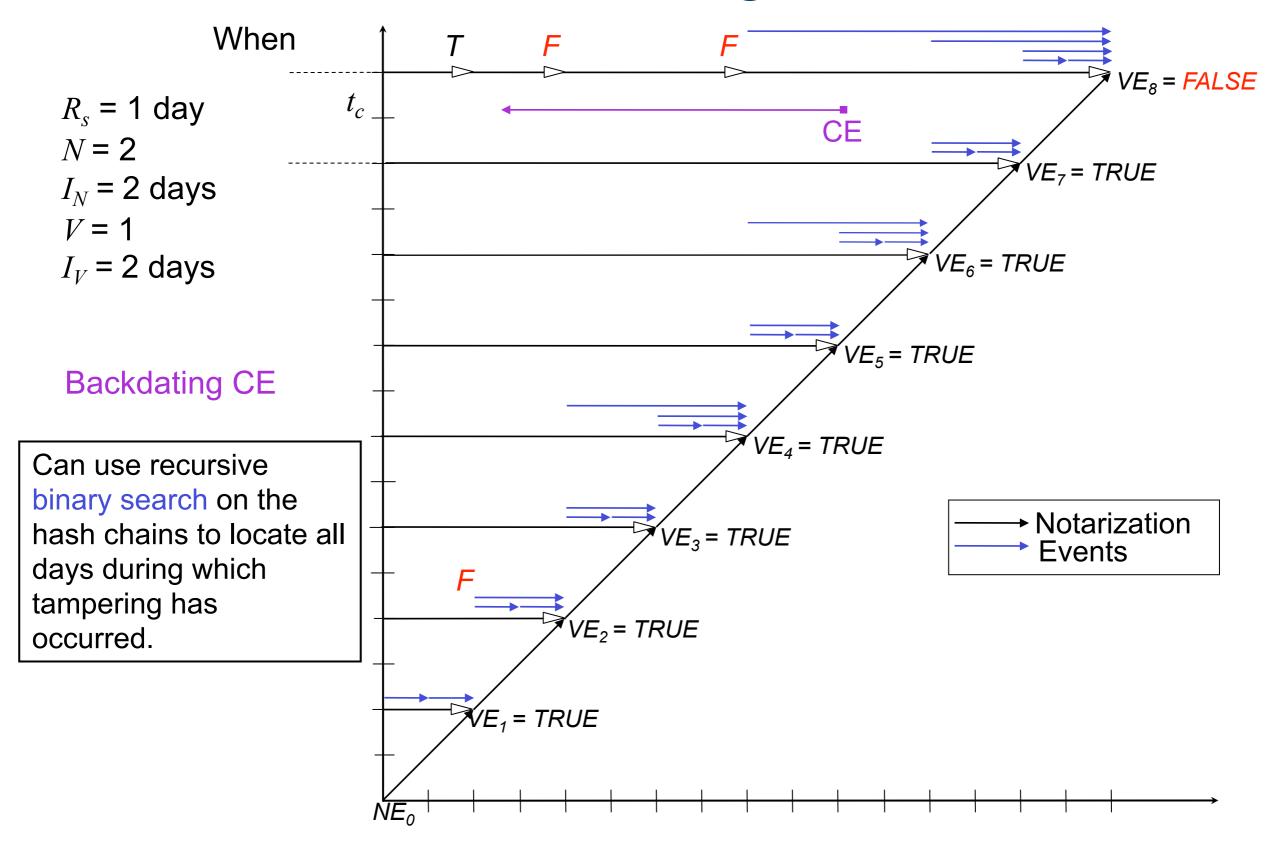


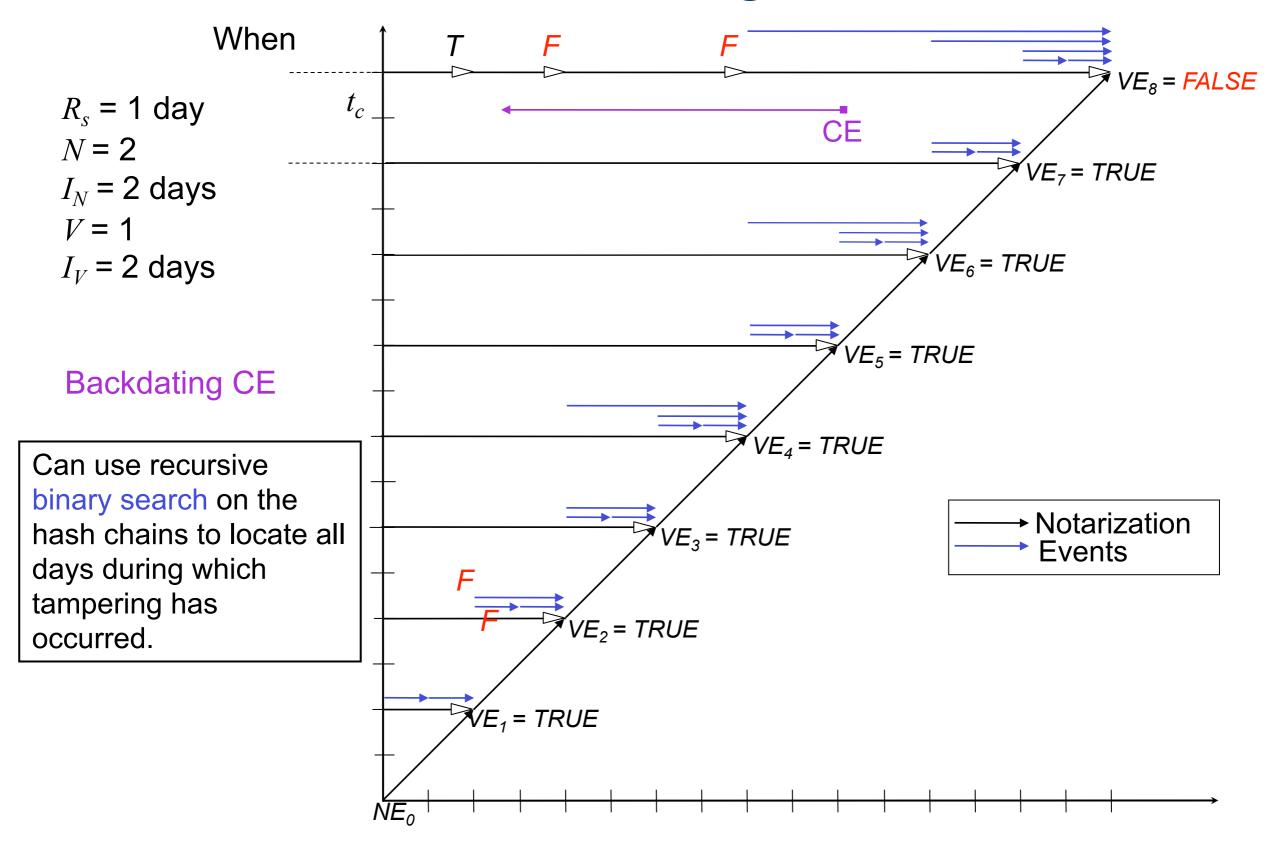


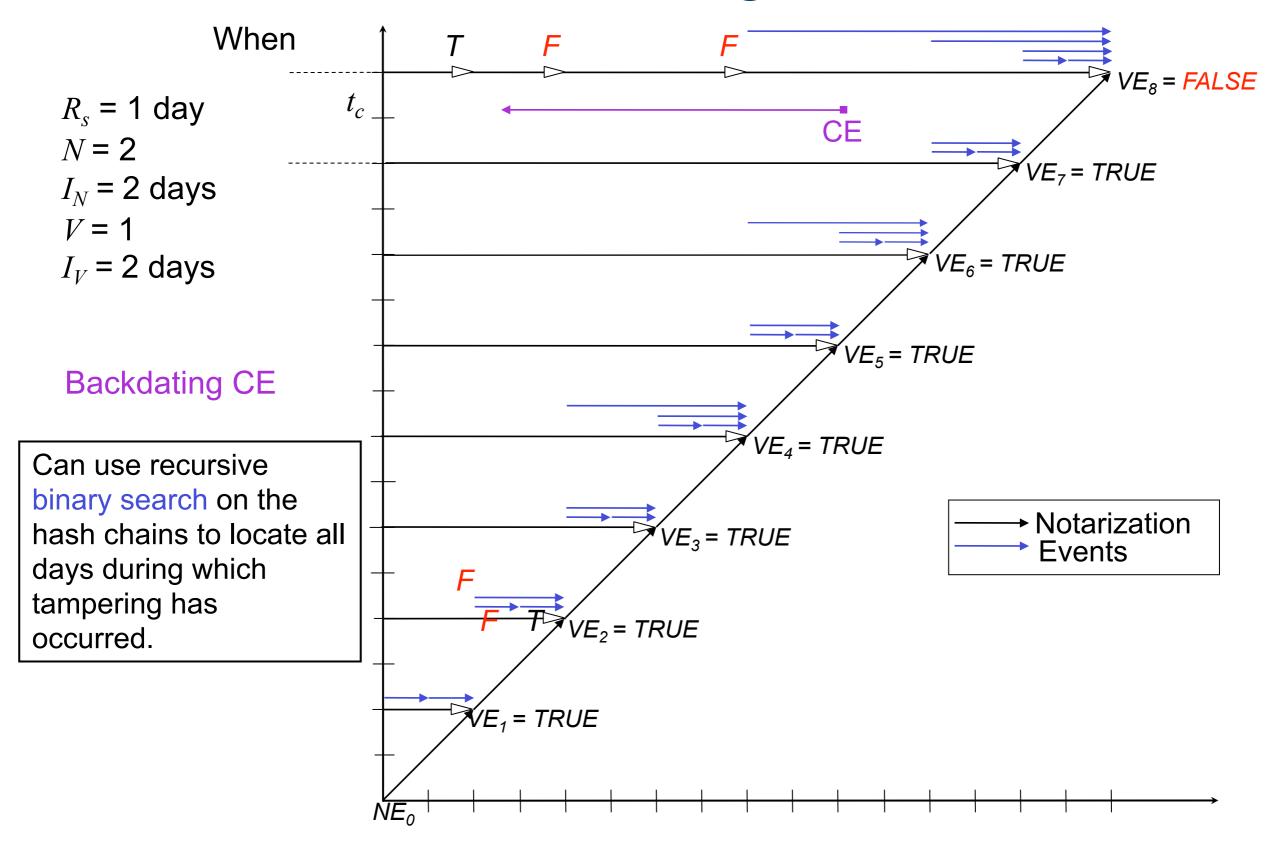




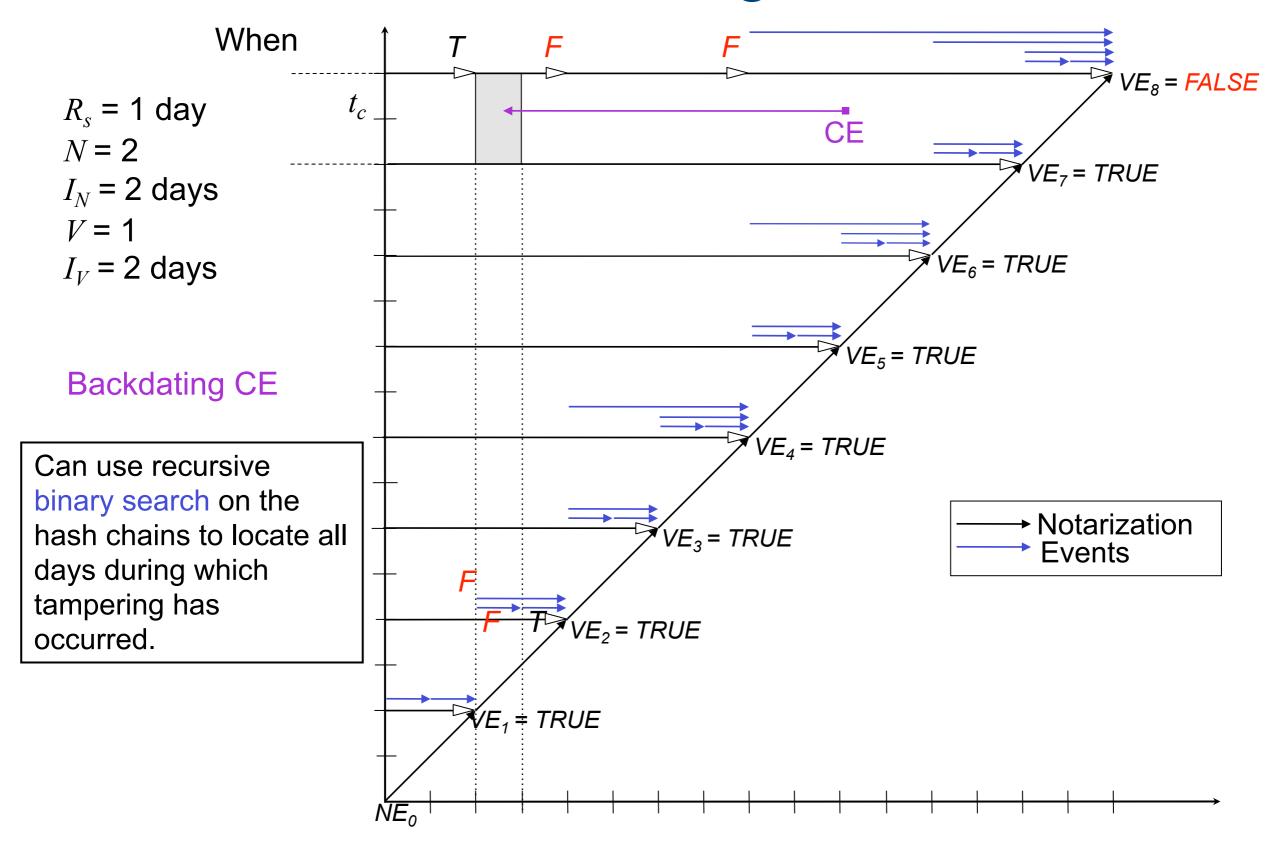


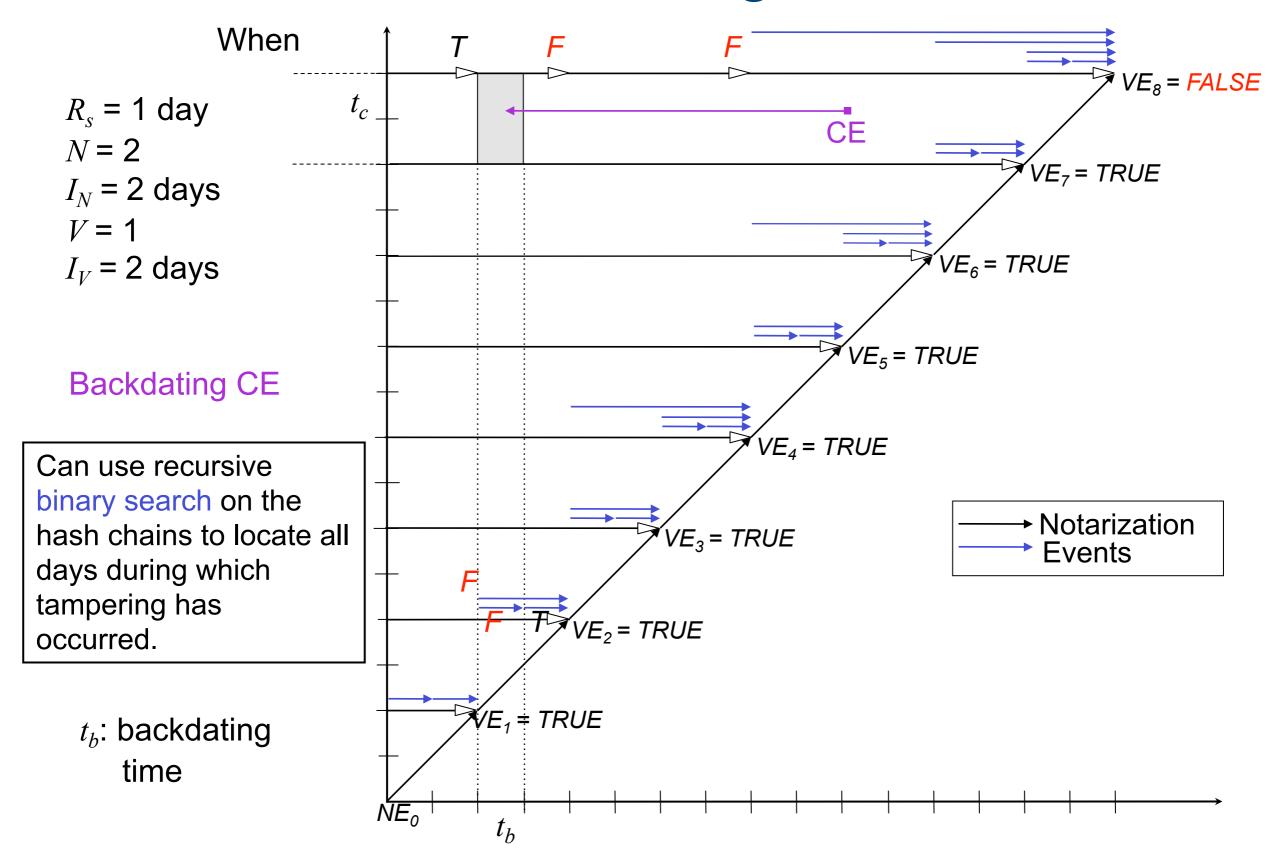


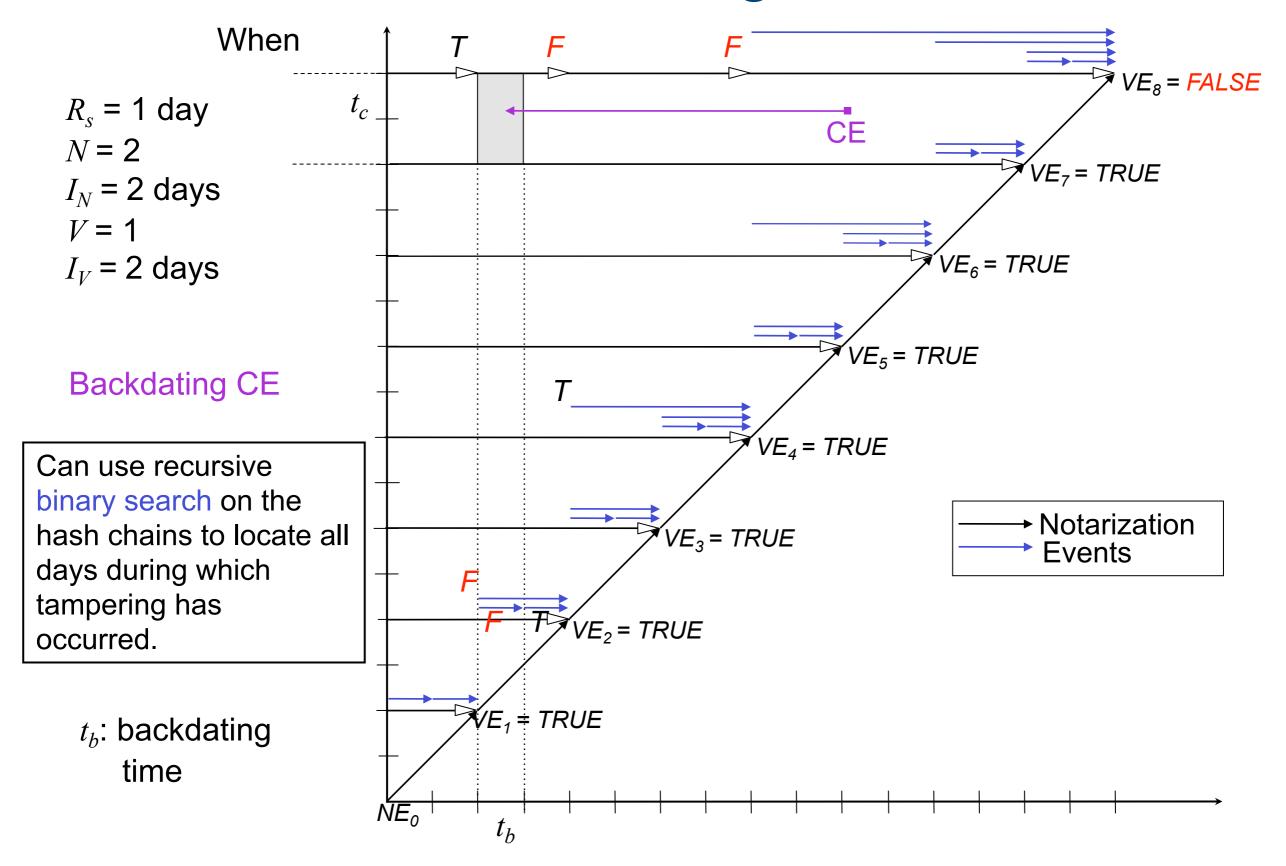


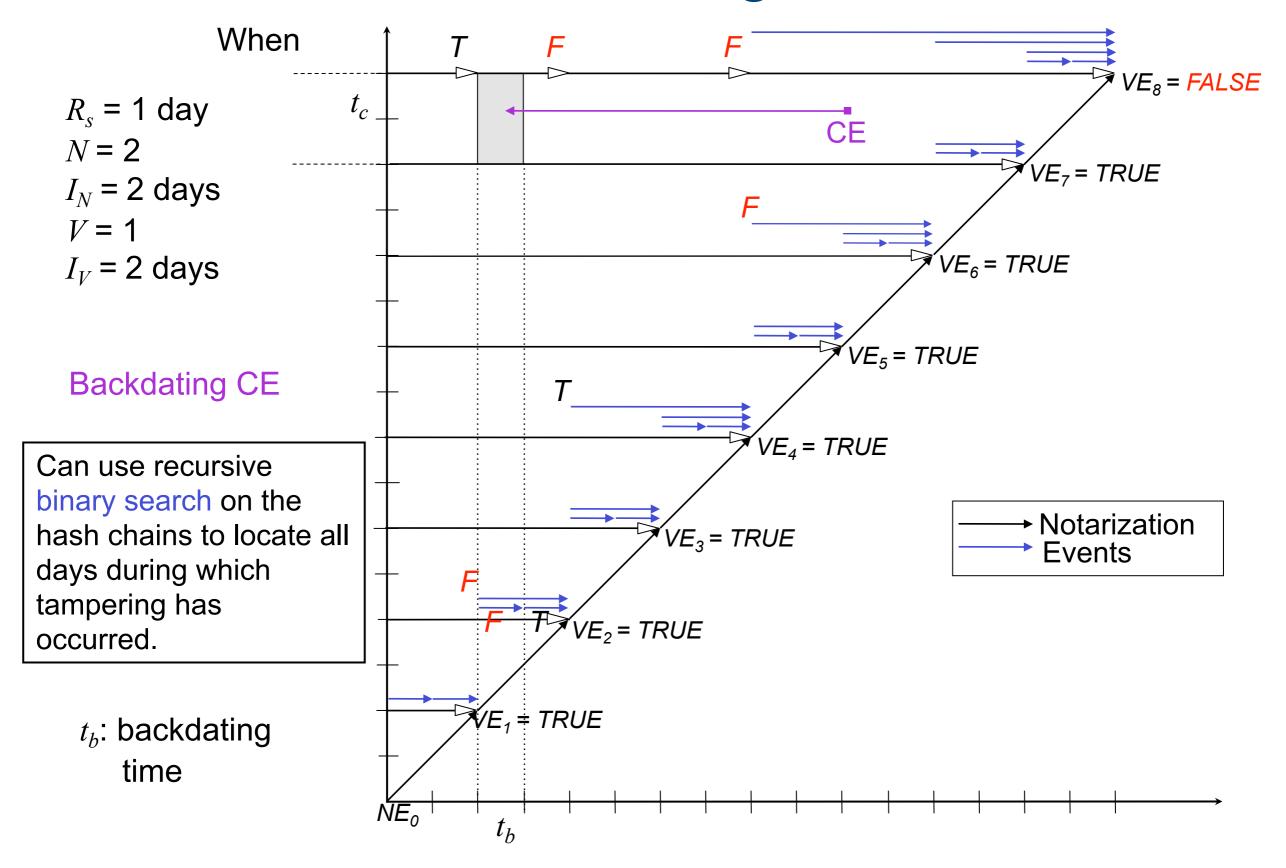


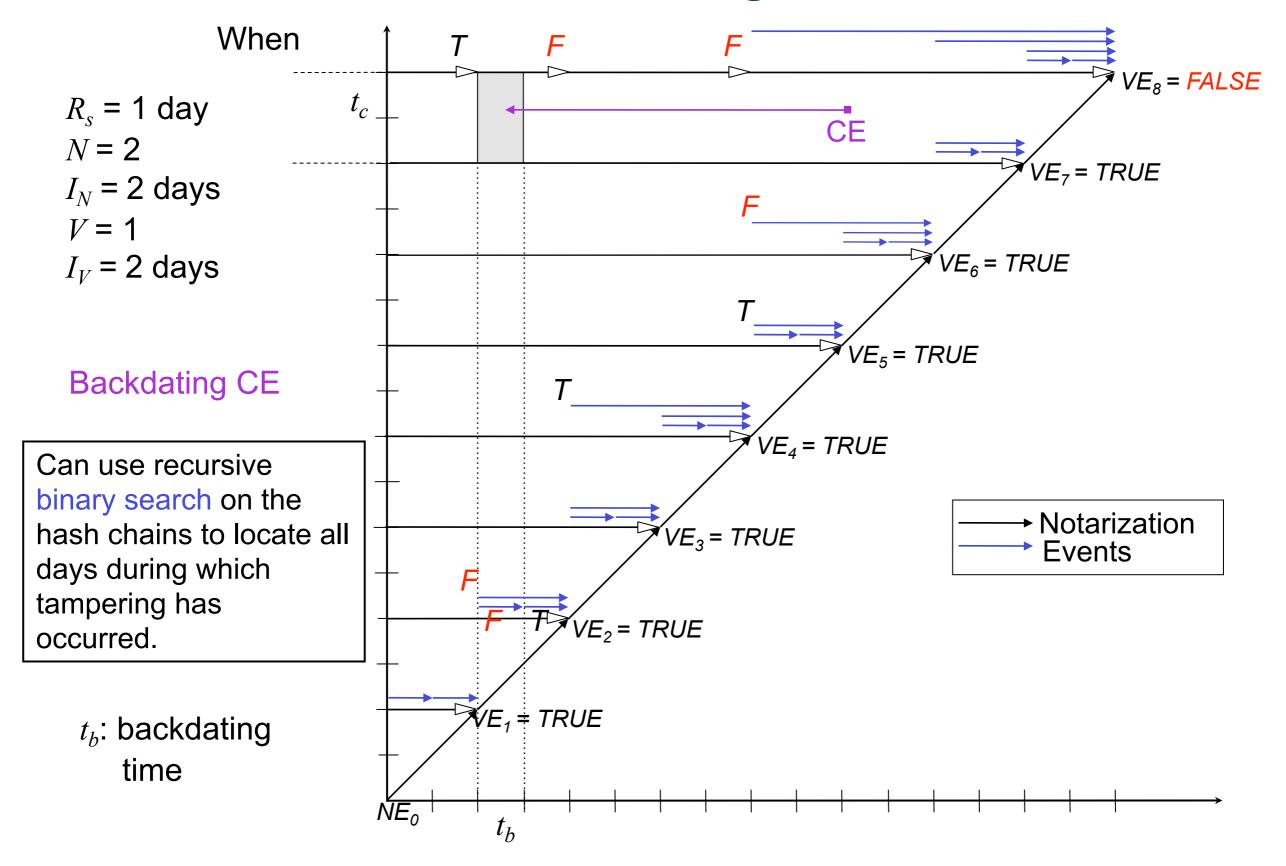
#### The a3D Algorithm When $\overrightarrow{VE}_8 = FALSE$ $t_c$ $R_s = 1 \text{ day}$ CE *N* = 2 $VE_7 = TRUE$ $I_N$ = 2 days V = 1 $VE_6 = TRUE$ $I_V$ = 2 days VE₅ = TRUE **Backdating CE** VE₄ = TRUE Can use recursive binary search on the Notarization hash chains to locate all $VE_3 = TRUE$ **Events** days during which tampering has $VE_2 = TRUE$ occurred. $VE_1 = TRUE$ NE<sub>0</sub>

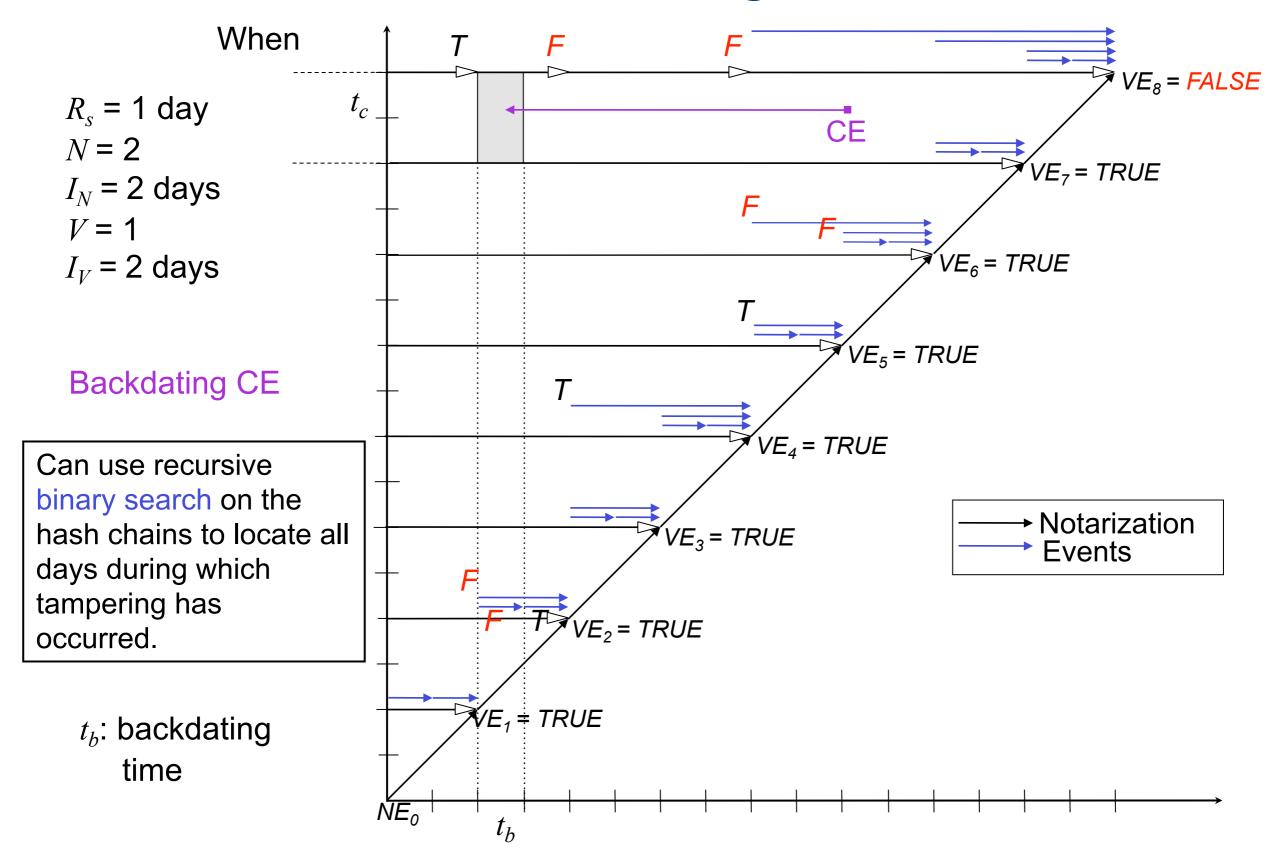


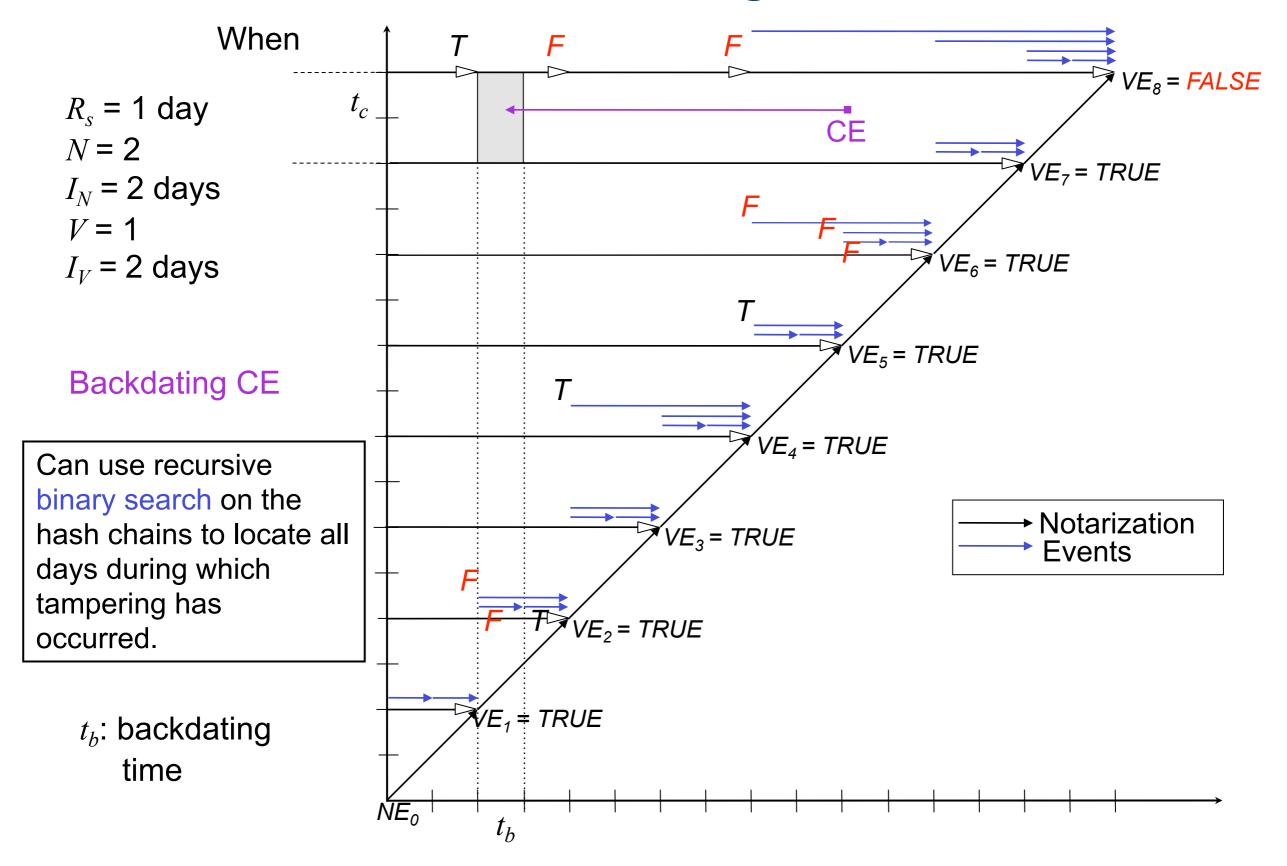


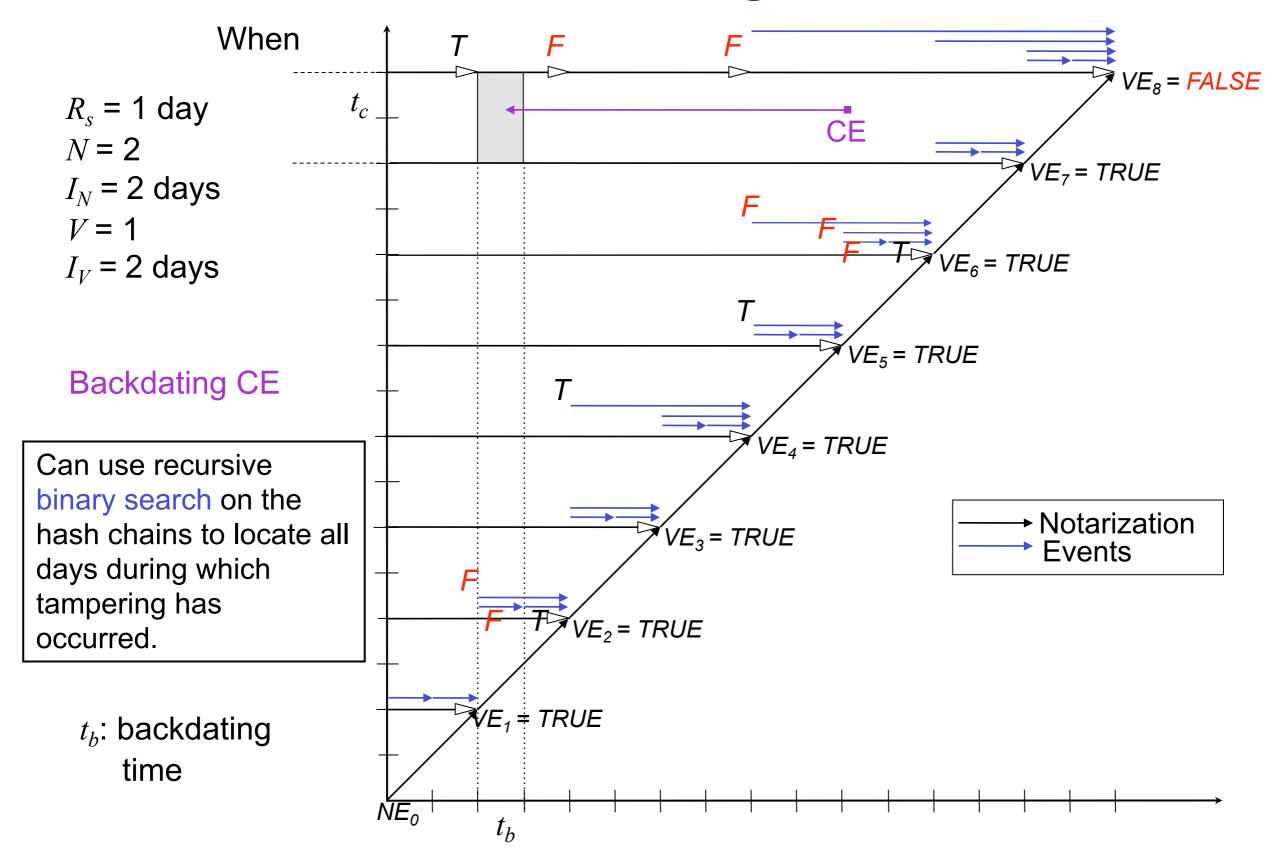


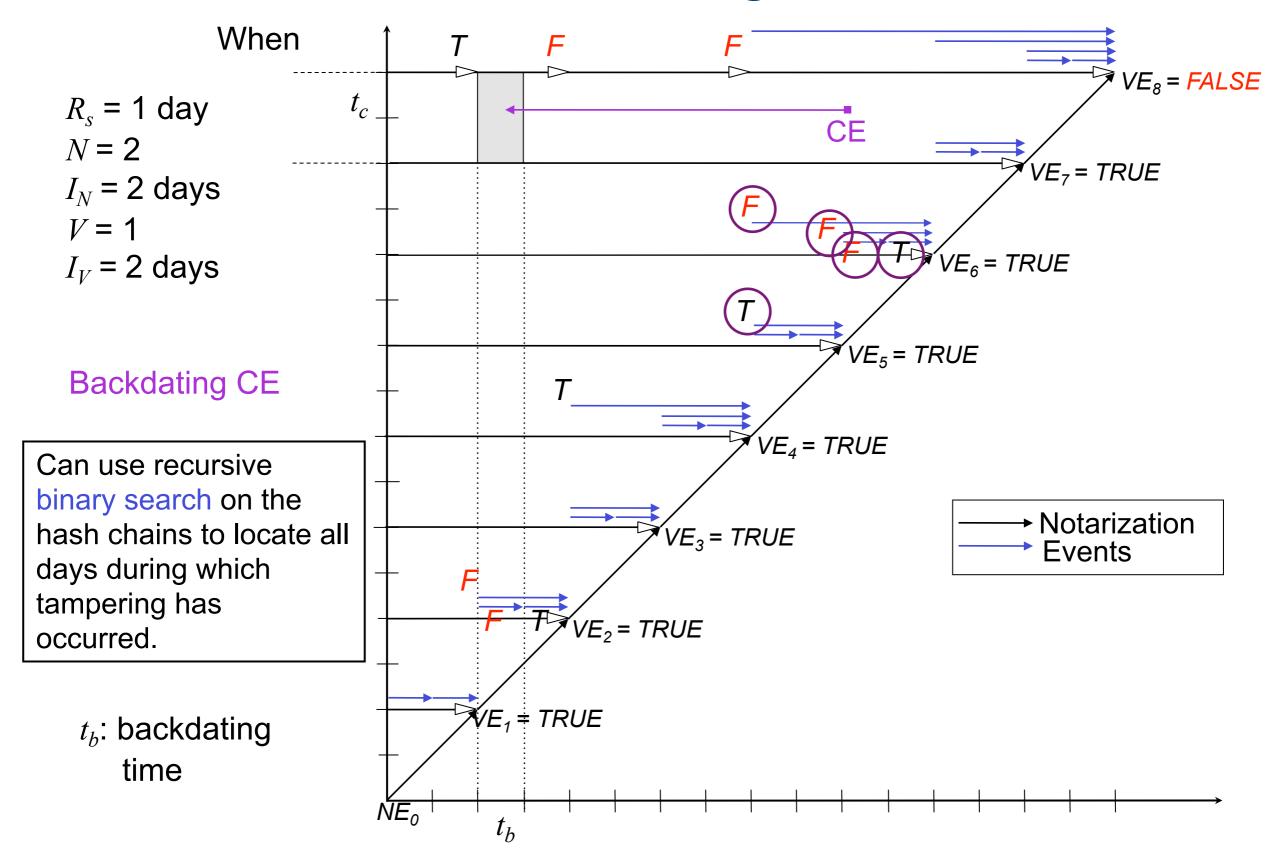


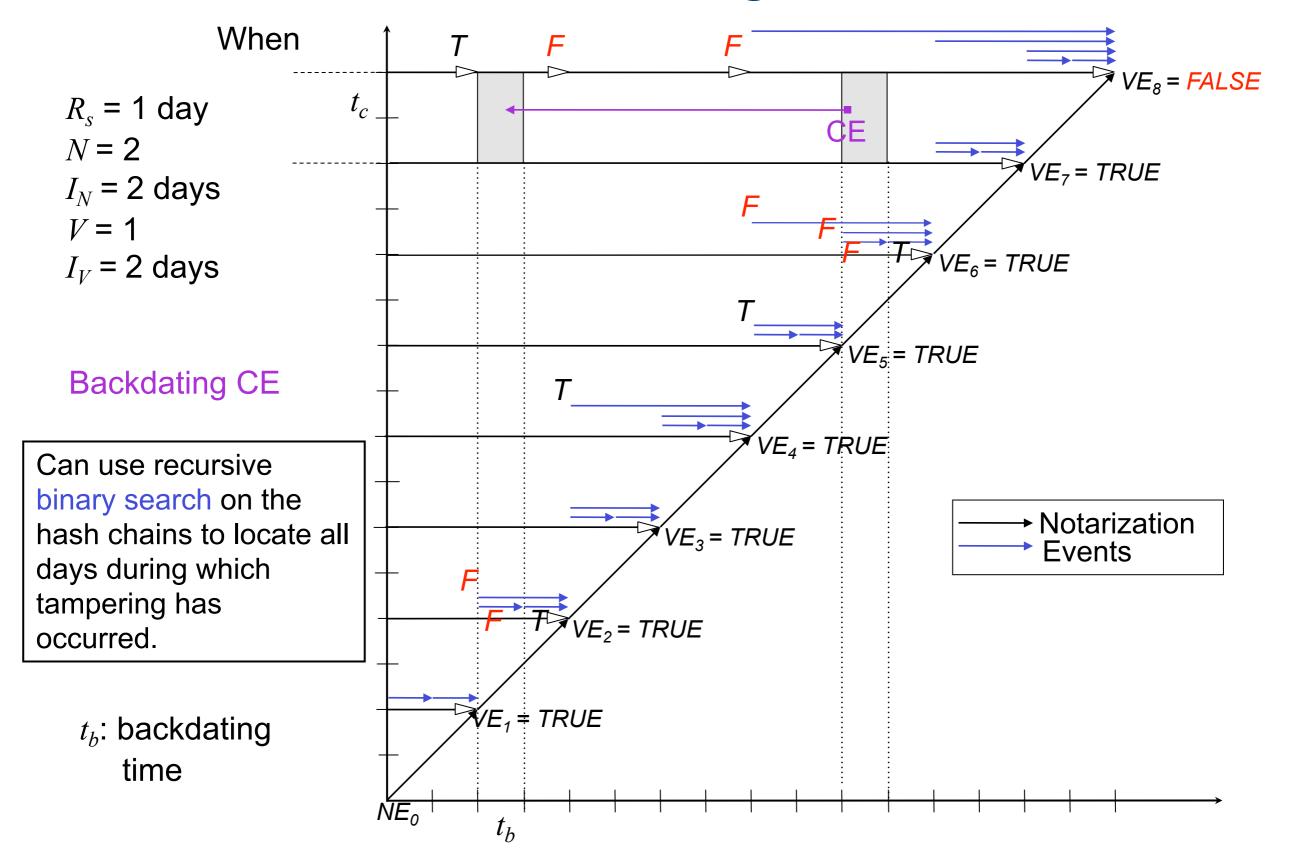


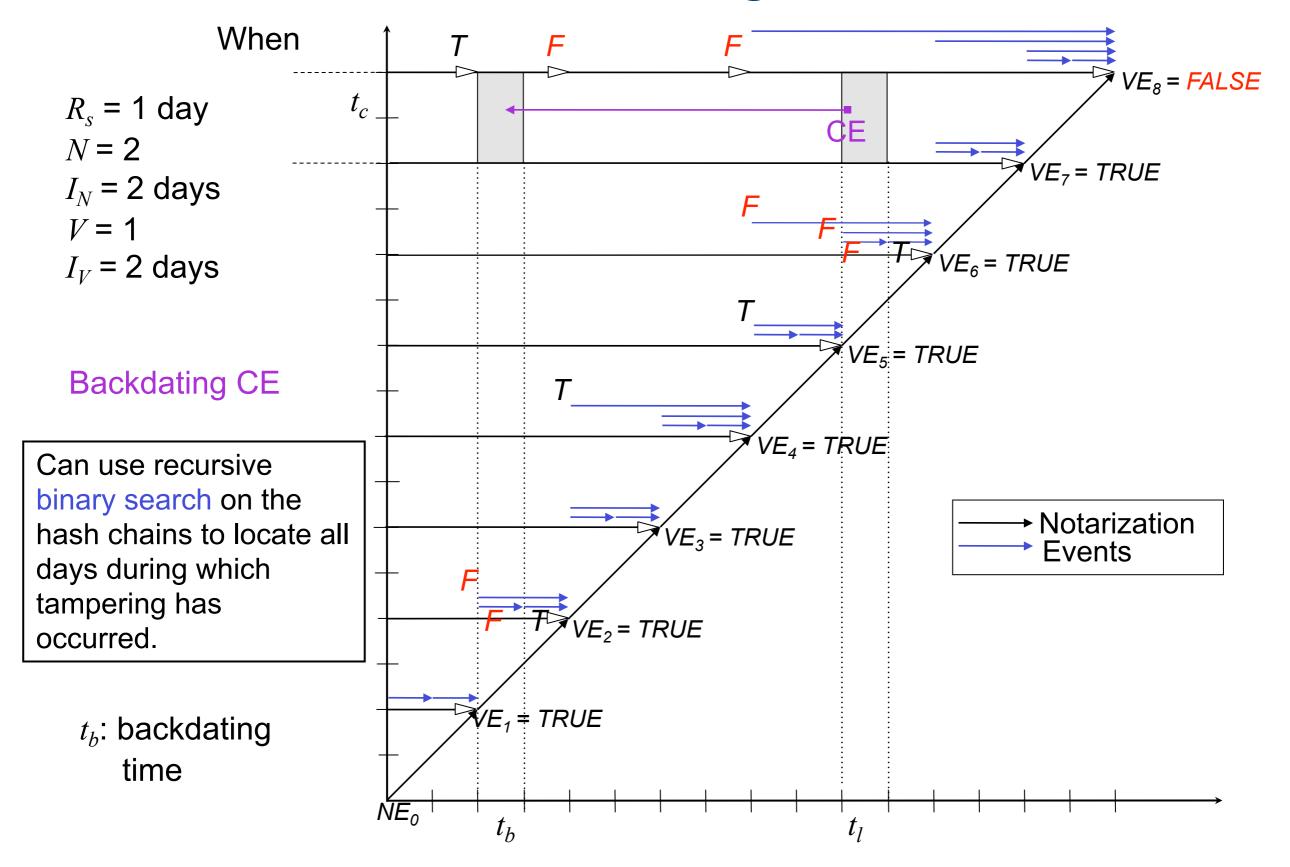


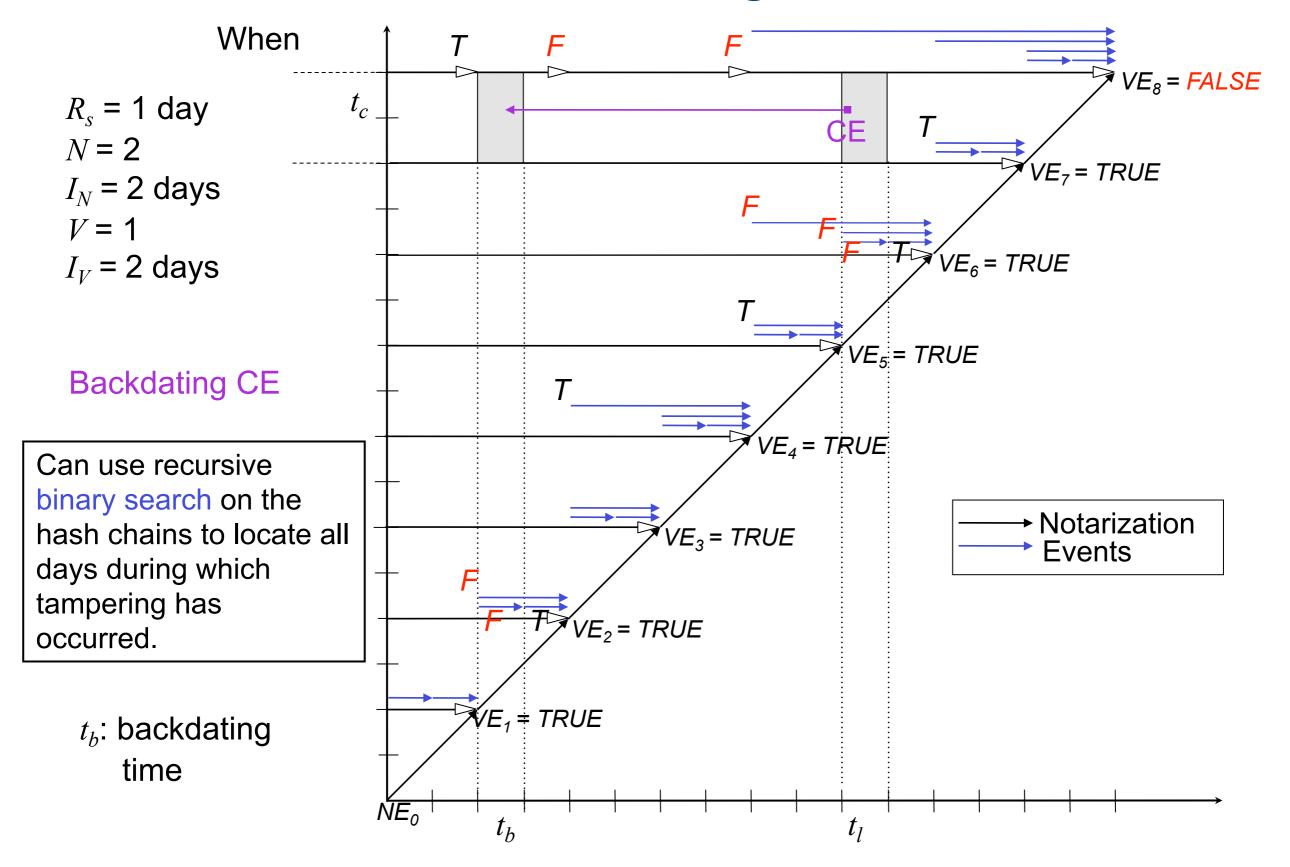


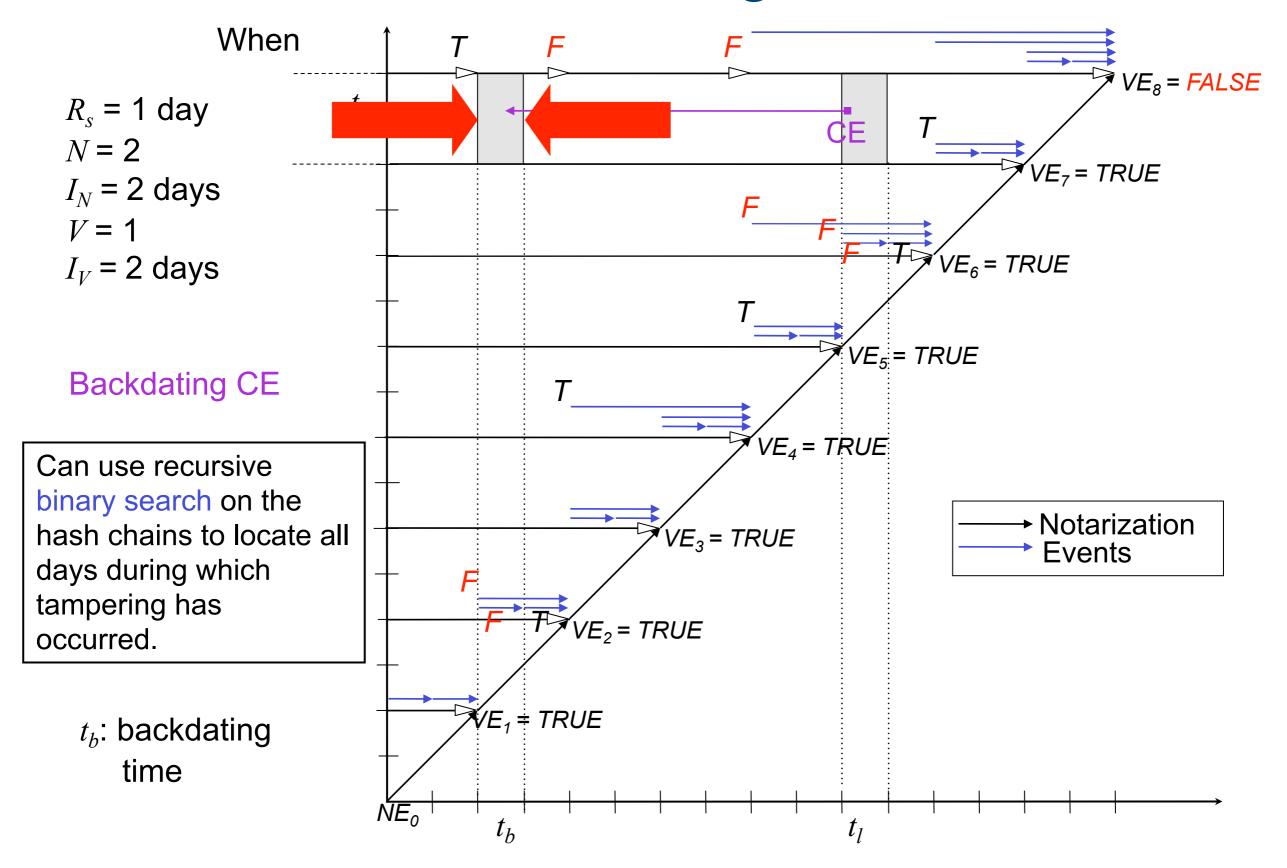


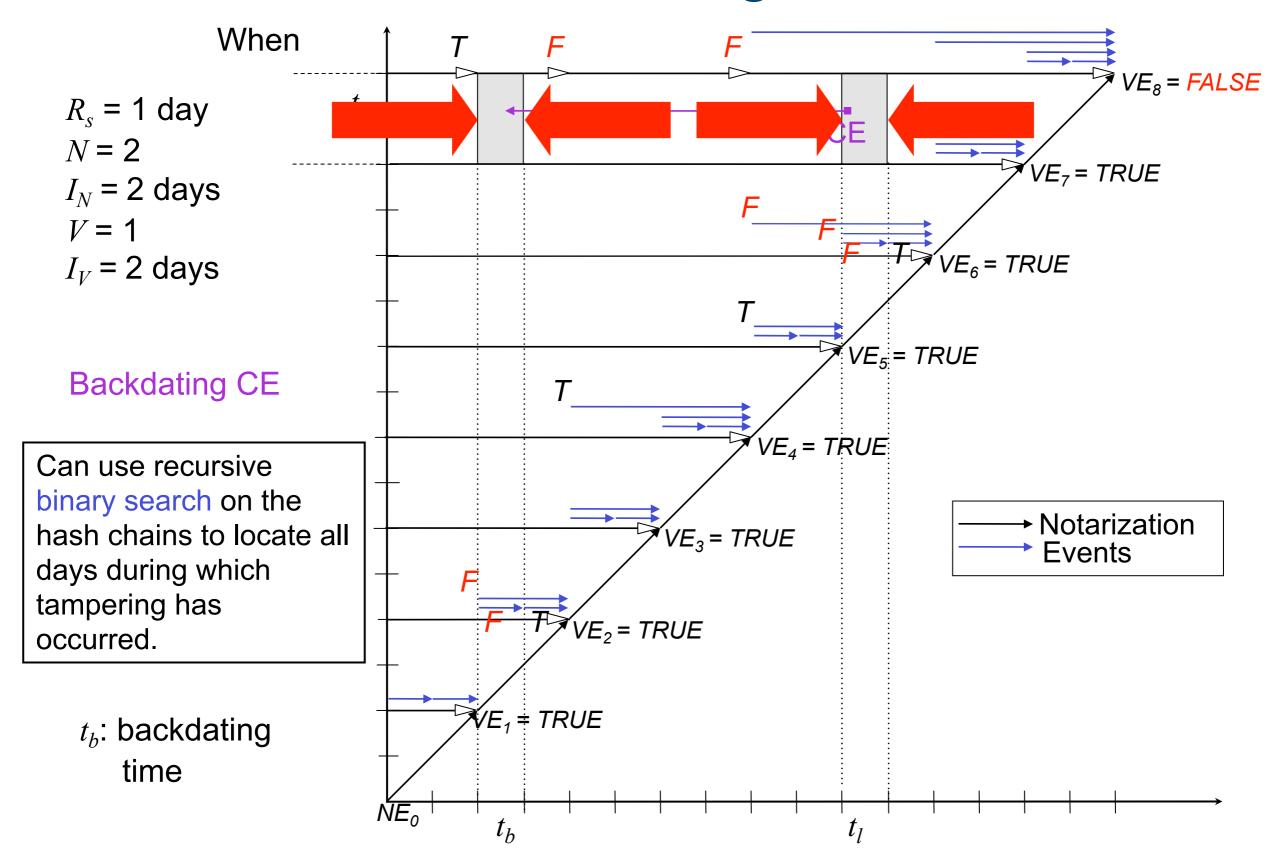






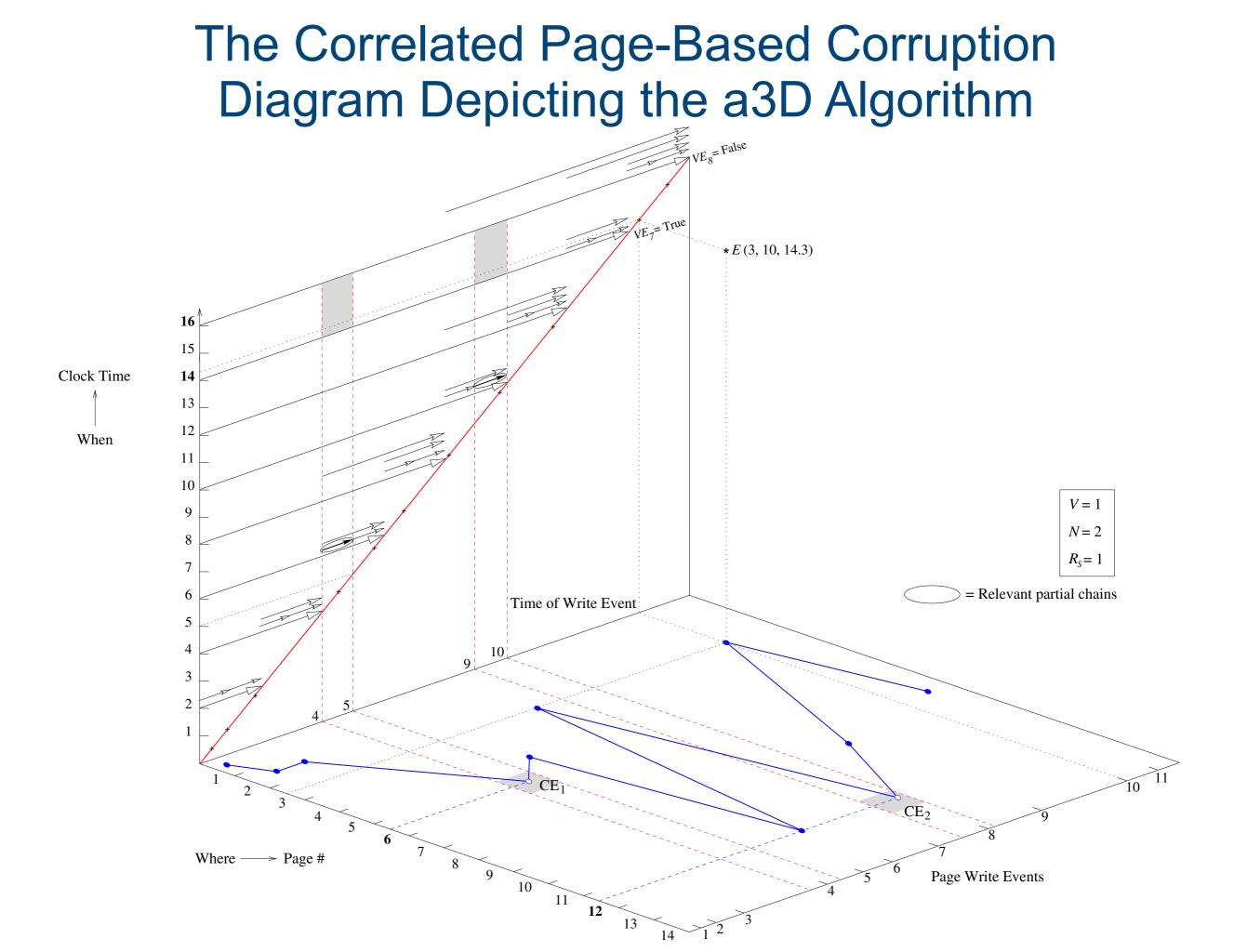


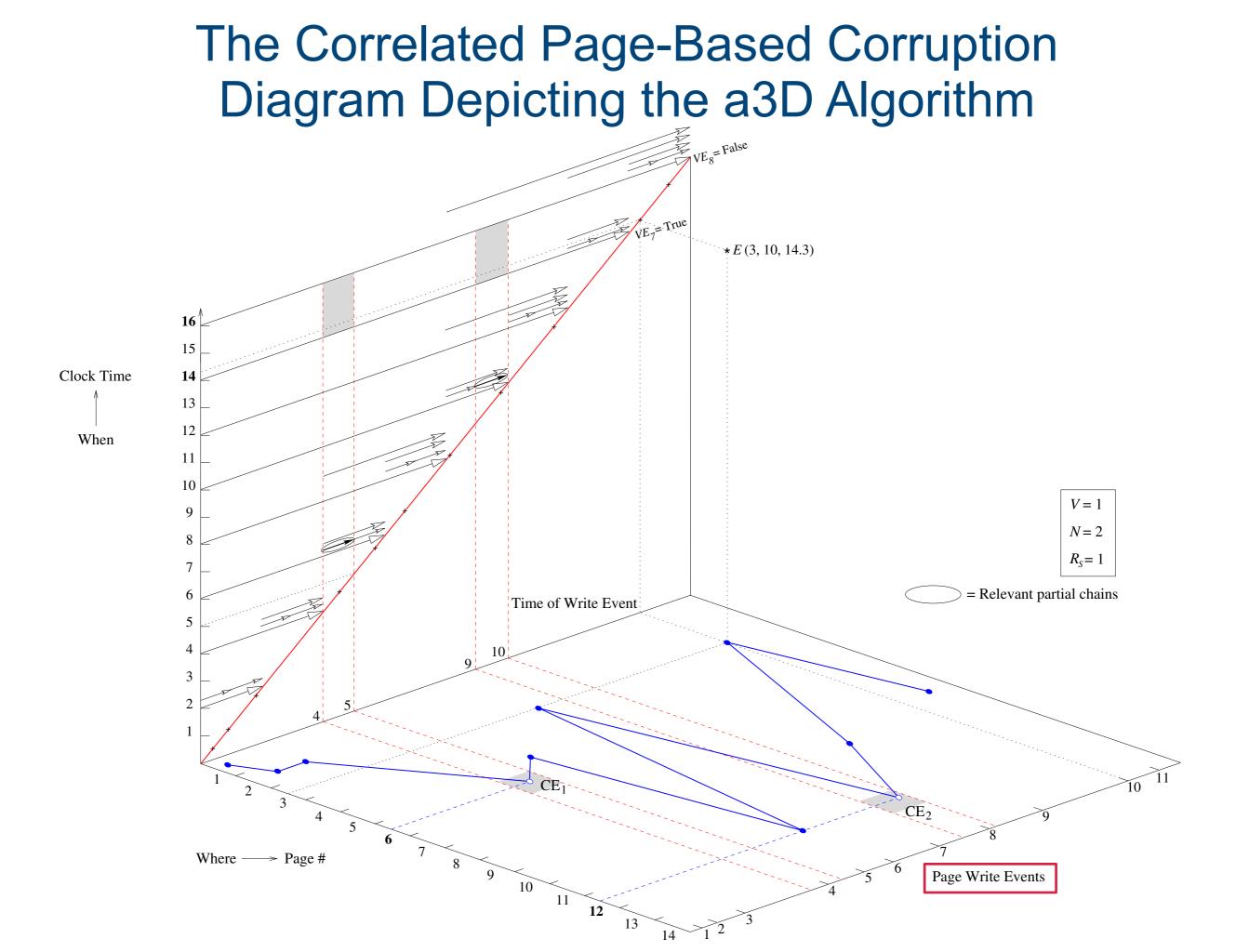


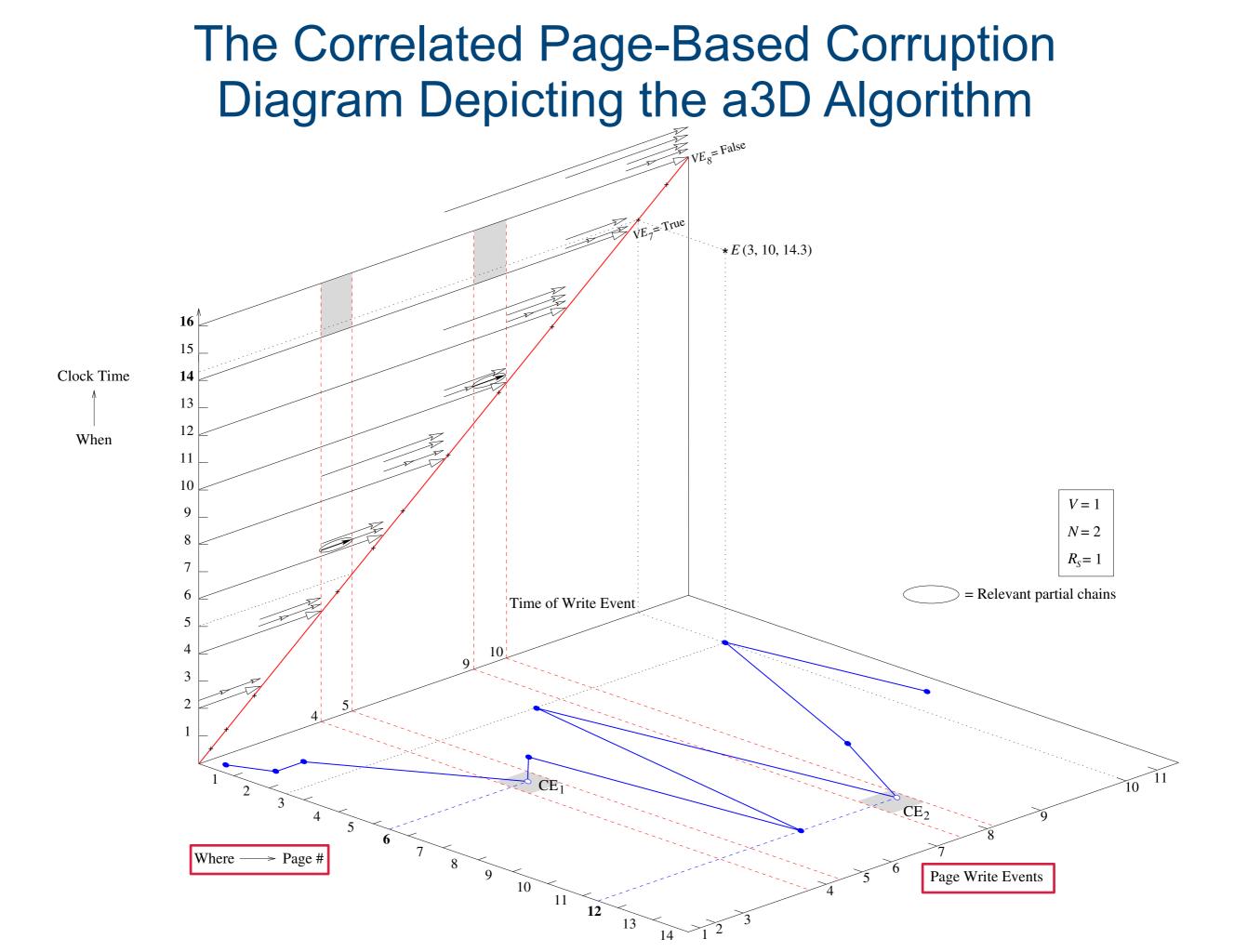


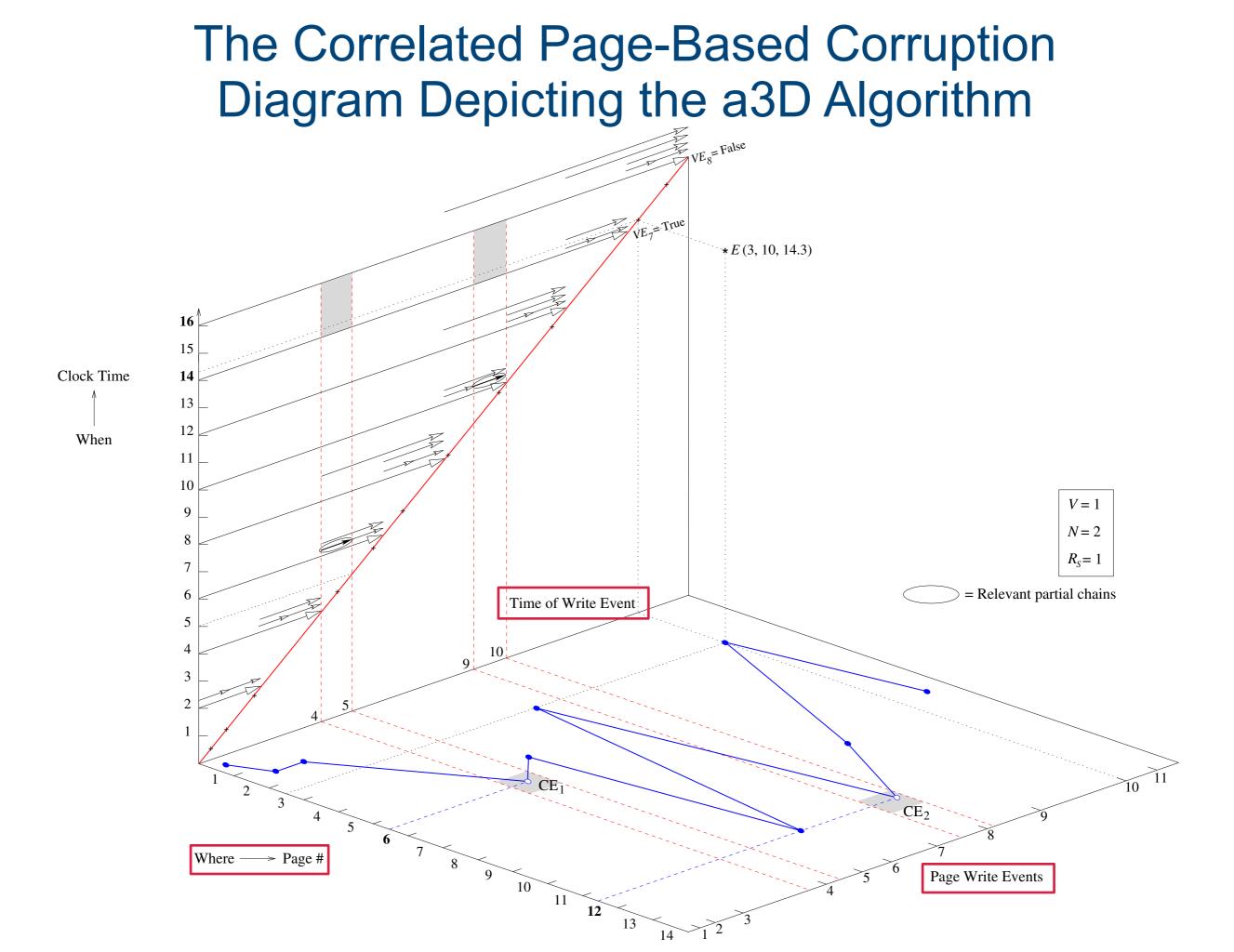
# **Comparison of Forensic Algorithms**

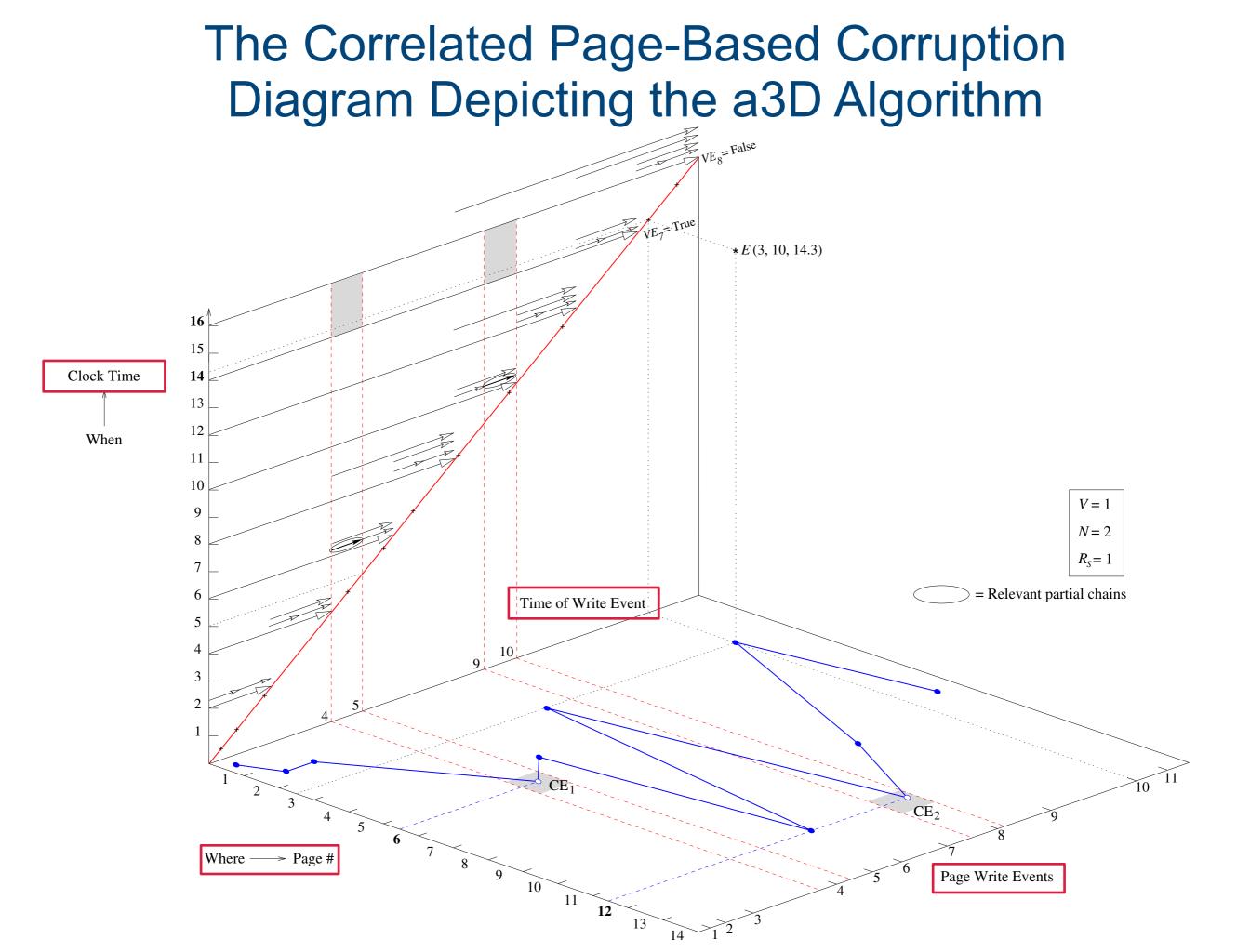
	Partitioning		
	Commit-Time-Based	Page-Based	Attribute-Based
Tables Affected	Any number	Any number	One or several of those containing the designated attribute
$R_s$	Time interval	Time interval	Time interval
$R_d$	N/A	N/A	Number of subsets of domain values
Segment	One of the contiguous <i>periods</i> induced by $R_s$ , starting from a particular anchor. Contiguous periods form a chronologically ordered partition.	One of the contiguous <i>periods</i> induced by $R_s$ , starting from a particular anchor. Contiguous periods form a chronologically ordered partition.	One of the contiguous <i>periods</i> induced by $R_s$ , starting from a particular anchor. Contiguous periods form a chronologically ordered partition.
Granule	Encompasses all tuples with <i>commit times</i> within the associated segment (one granule has tuples from many transactions committing in that segment).	Encompasses all tuples whose physical location is in a page mentioned within the associated segment.	Encompasses all tuples with <i>commit times</i> within the associated segment (one granule has tuples from many transactions committing in that segment).
Hashing order	Transactions hashed in order of increasing commit time.	Granules hashed in chronological order of "page write" event of the page. Granules not hashed in order of page number.	Transactions hashed in order of increasing commit time.
Segment	When the last tuple in the	When the last	When the last tuple in the
Completion Event	granule associated with that segment commits	page write event in the segment occurs.	granule associated with that segment commits
Notarization Factor ( N)	Specified by DBA	Specified by DBA	Specified by DBA
Validation Factor (V)	Specified by DBA	Specified by DBA	Specified by DBA
$I_N$	$N \times R_s$	$N \times R_s$	$N \times R_s$
$I_V$	$V \times I_N$	$V \times I_N$	$V \times I_N$
Notarization	Occurs as soon as N granules are hashed.	Occurs as soon as N granules are hashed.	Occurs as soon as N granules are hashed.
Validation	Occurs as soon as V notarizations have occurred.	Occurs as soon as V notarizations have occurred.	Occurs as soon as V notarizations have occurred.

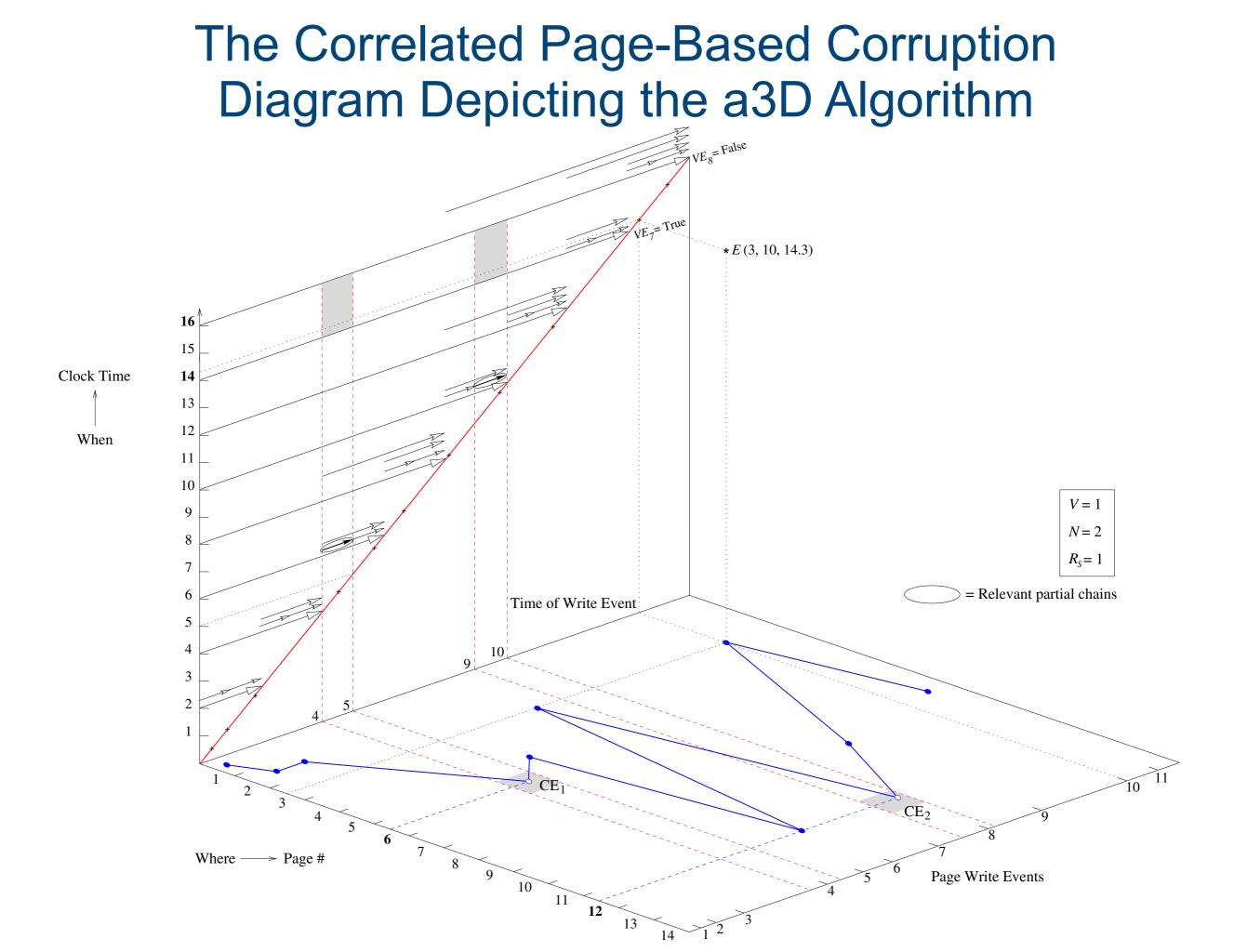


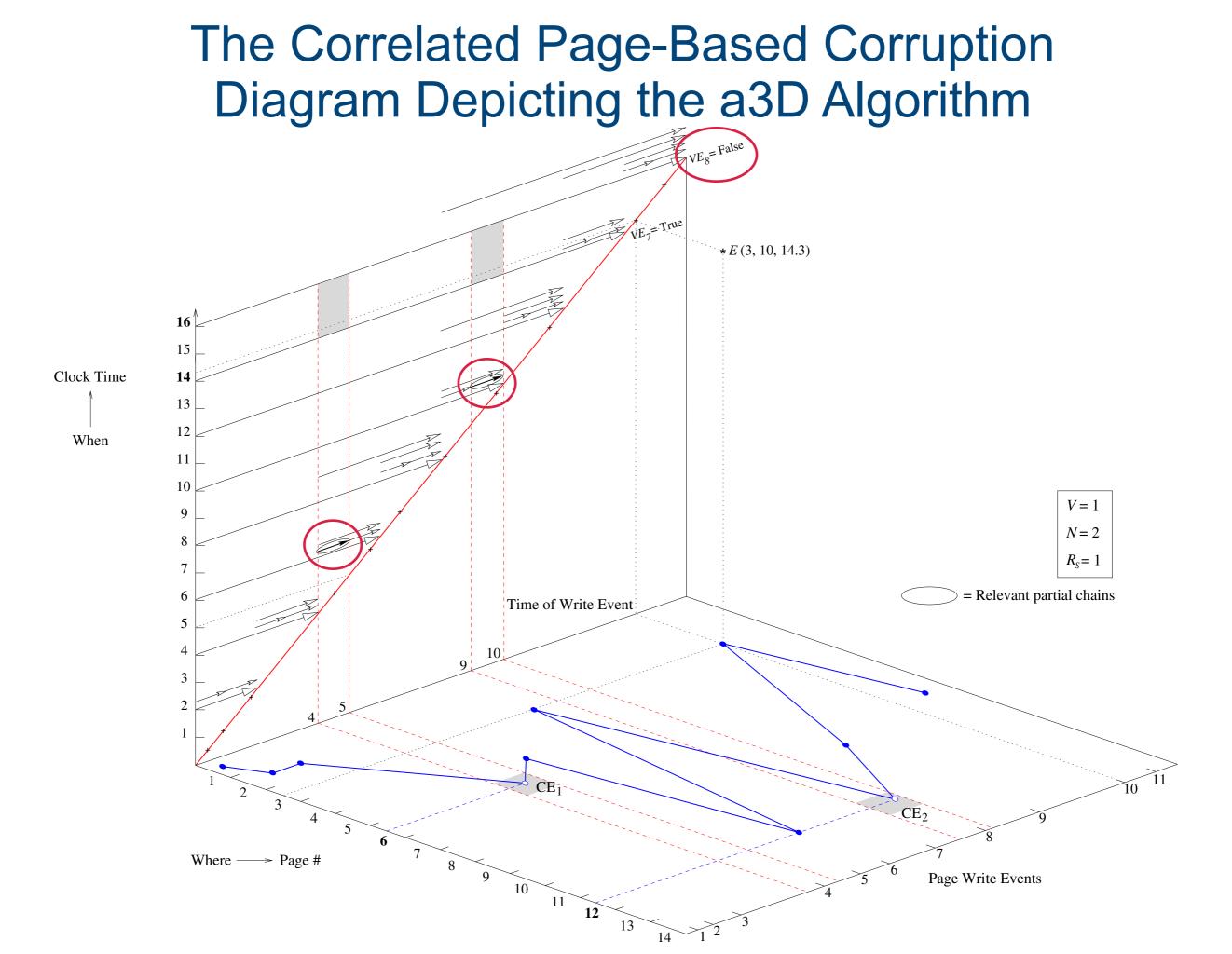


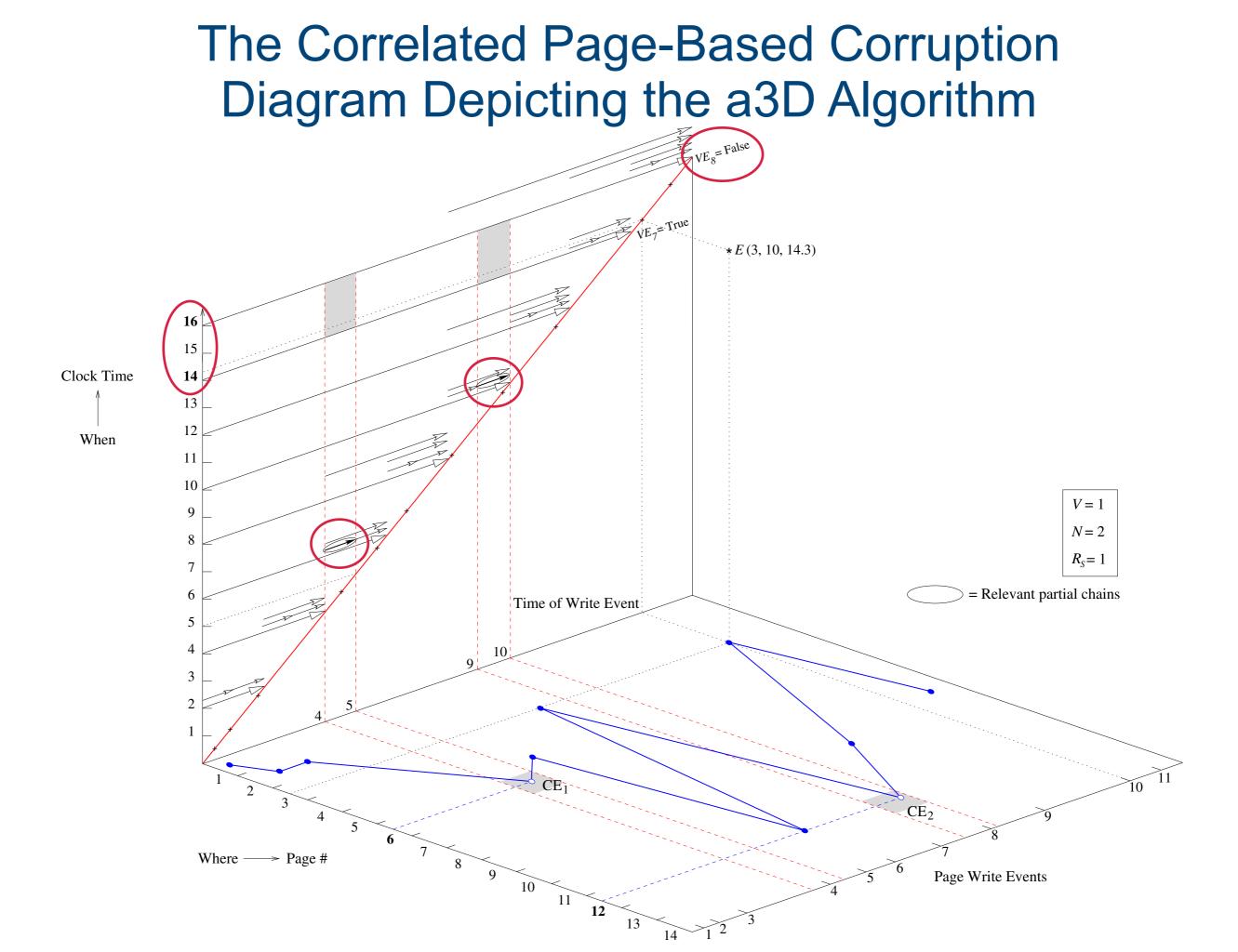


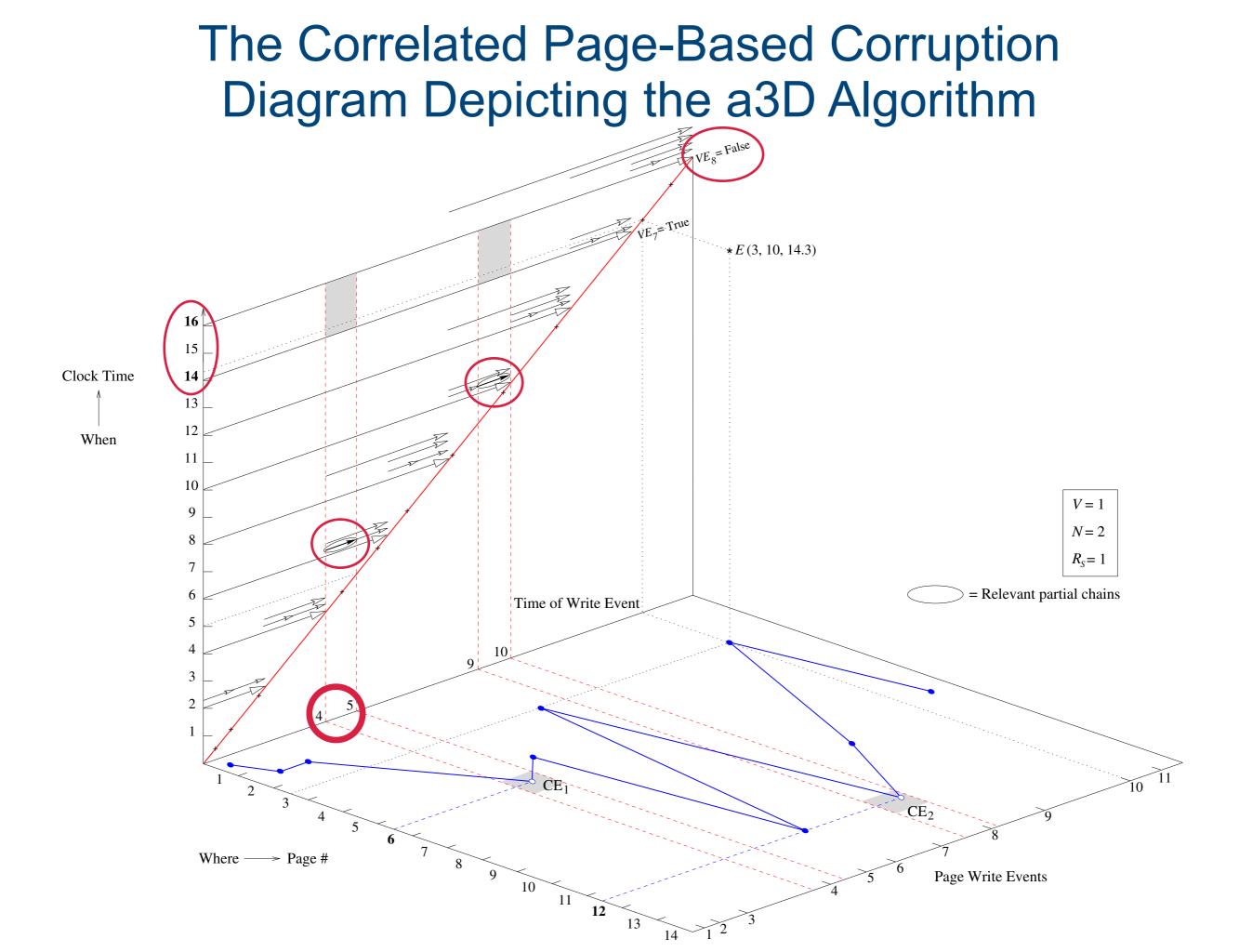


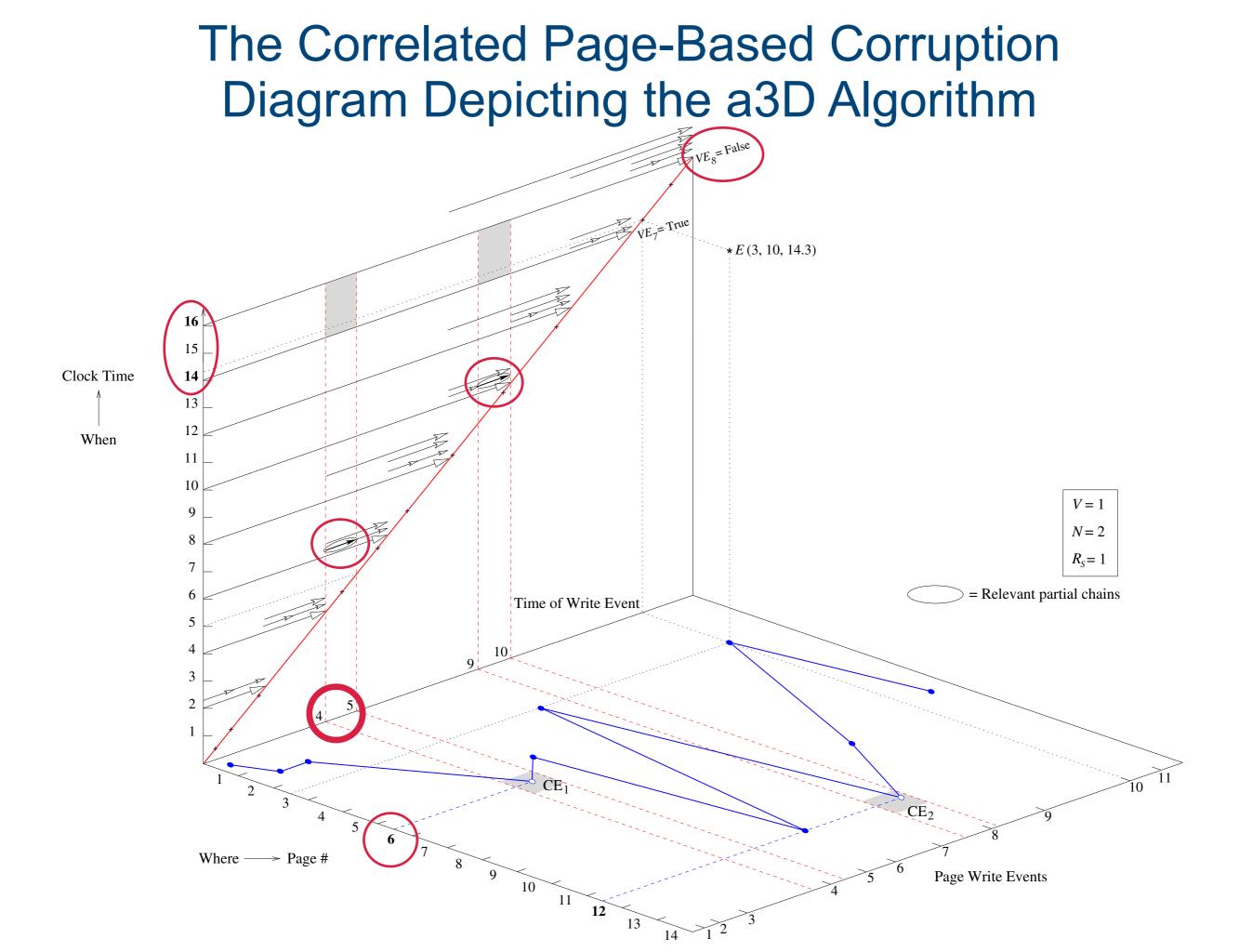






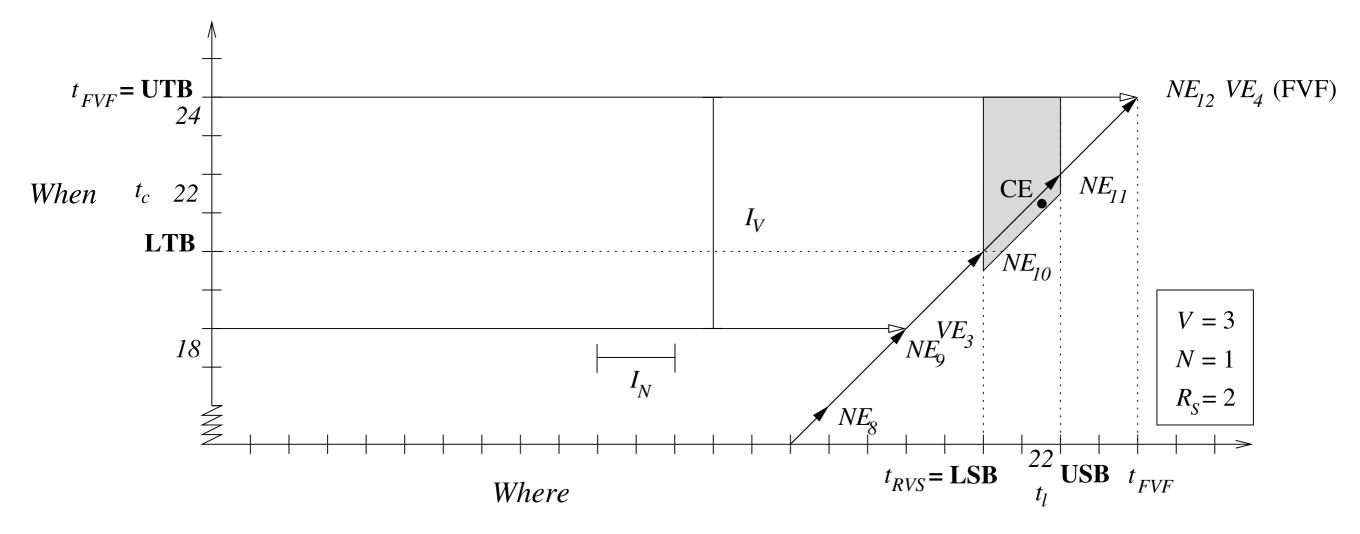


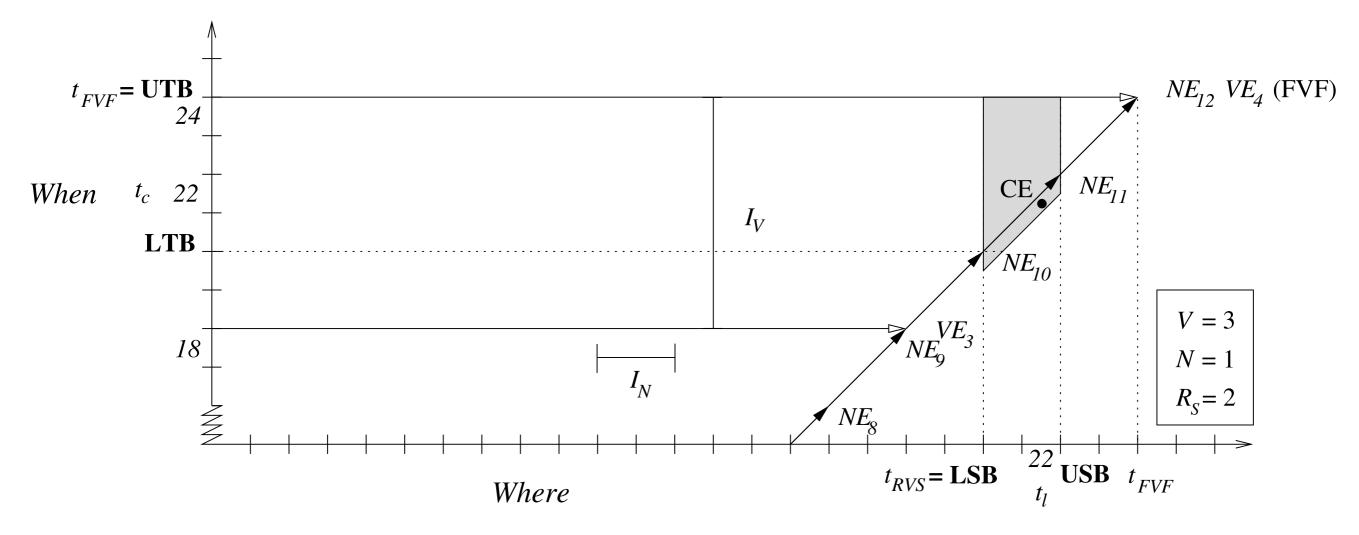




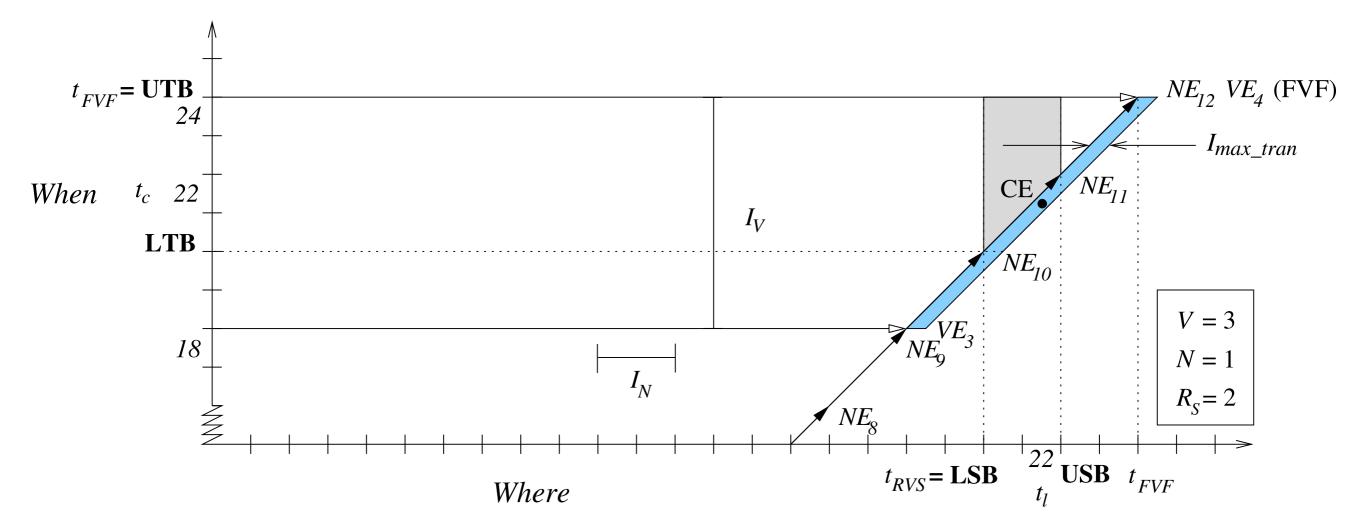
# Outline

- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations





- In general  $t_c > t_l$
- Exception: corruption affects currently executing transaction



- In general  $t_c > t_l$
- Exception: corruption affects currently executing transaction
- Introduce "envelope" of width  $I_{max tran}$

Different solution: Exploit the Regret Interval

- Different solution: Exploit the Regret Interval
- The regret interval  $(I_R)$  is the minimum time interval before any adversary can reverse the change they made.

- Different solution: Exploit the Regret Interval
- The regret interval ( $I_R$ ) is the minimum time interval before any adversary can reverse the change they made.
- *I<sub>R</sub>* is intrinsic to the semantics and social use of application.
  We have no control over it.

- Different solution: Exploit the Regret Interval
- The regret interval  $(I_R)$  is the minimum time interval before any adversary can reverse the change they made.
- *I<sub>R</sub>* is intrinsic to the semantics and social use of application.
  We have no control over it.
- We use an estimate  $I_R^* \leq I_R$

- Different solution: Exploit the Regret Interval
- The regret interval  $(I_R)$  is the minimum time interval before any adversary can reverse the change they made.
- *I<sub>R</sub>* is intrinsic to the semantics and social use of application.
  We have no control over it.
- We use an estimate  $I_R^* \leq I_R$
- No introactive corruptions:  $0 \le I_N \le I_V \le I_R^* \le I_R$



# Shredding

- Transaction time semantics require that data are never physically deleted.
  - Performance overhead
  - Privacy and liability threat

# Shredding

- Transaction time semantics require that data are never physically deleted.
  - Performance overhead
  - Privacy and liability threat
- Retention period: a sliding time frame  $I_{RP}$ 
  - Determined by regulations & company policy
  - Record physically deleted after exiting  $now I_{RP}$

# Shredding

- Transaction time semantics require that data are never physically deleted.
  - Performance overhead
  - Privacy and liability threat
- Retention period: a sliding time frame  $I_{RP}$ 
  - Determined by regulations & company policy
  - Record physically deleted after exiting  $now I_{RP}$
- Shredding ensures information restriction.
  - Breaks semantics of information accountability
  - Reconcile shredding with tamper detection and forensic analysis?

• Litigation holds can be issued on the data for a duration of time as specified by a court.

- Litigation holds can be issued on the data for a duration of time as specified by a court.
- Override retention period regulations

- Litigation holds can be issued on the data for a duration of time as specified by a court.
- Override retention period regulations
- Litigation holds "restore" info accountability.

- Litigation holds can be issued on the data for a duration of time as specified by a court.
- Override retention period regulations
- Litigation holds "restore" info accountability.
- The capability to impose litigation holds prevents indiscriminate shredding and ensures accountability.

## Outline

- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations

- There are three GUIs:
  - Chief Security Office (CSO)
  - Database Administrator (DBA)
  - Crime Scene Investigator (CSI)

- There are three GUIs:
  - Chief Security Office (CSO)
  - Database Administrator (DBA)
  - Crime Scene Investigator (CSI)
- Configure the security policies by
  - selecting a database to be monitored
  - setting the security parameters, e.g.,  $R_s$ , N, V,  $I_N$

- There are three GUIs:
  - Chief Security Office (CSO)
  - Database Administrator (DBA)
  - Crime Scene Investigator (CSI)
- Configure the security policies by
  - selecting a database to be monitored
  - setting the security parameters, e.g.,  $R_s$ , N, V,  $I_N$
- Calculate the forensic cost for normal processing and forensic analysis

- There are three GUIs:
  - Chief Security Office (CSO)
  - Database Administrator (DBA)
  - Crime Scene Investigator (CSI)
- Configure the security policies by
  - selecting a database to be monitored
  - setting the security parameters, e.g.,  $R_s$ , N, V,  $I_N$
- Calculate the forensic cost for normal processing and forensic analysis
- Create corruption diagrams

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 👻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 🔽 hrs 1 🖵 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

4	DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File	
acmedb(	/home/tau/software/audit/auditdb) 👻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings	Detected Tampering
	Detection Resolution Unit: 0 🔻 days 0 🖝 hrs 1 🖝 mins.
	Forensic Algorithm: Monochromatic 🛛 🖛
	Number of Resolution Units Between Notarizations 1
	Time between notarizations: 0 days 0 hrs 1 mins.
	Number of Notarizations Between Validations 1
	Time between validations: 0 days 0 hrs 1 mins.
	Cost Per Unit: \$ 0.01
	Predicted Cost:
	Tampering Detection:
	Per Day \$: 0.02 Per Year \$: 7.30
	Forensic Analysis (Worst Case):
	One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
	Start On: 7 🗸 / 22 🗸 / 2012 🗸 at 23 <table-cell> : 52 🗸</table-cell>
	Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 💌 hrs 1 💌 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🚽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 🔽 hrs 1 🔽 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 0 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

DBA: Nirav Merchant Email:nirav@email.arizona.ee	du 💶 🗙
File	
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass	rts@cs.arizona.edu)
Settings Detected Tampering	
Detection Resolution Unit: 0 🔽 days 0 🖵 hrs 1 🗨	mins.
Forensic Algorithm: Monochromatic 🔽	
Number of Resolution Units Between Notarizations 1	-
Time between notarizations: 0 days 0 hrs 1	mins.
Number of Notarizations Between Validations 1	•
Time between validations: 0 days 0 hrs 1	mins.
Cost Per Unit: \$ 0.01	
Predicted Cost:	
Tampering Detection:	
Per Day \$: 0.02 Per Year \$: 7.30	
Forensic Analysis (Worst Case):	
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.0	0
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼:	52 🔽
Save These Settings	

<u>ب</u>	DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File	
acmedb(/	home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings	Detected Tampering
	Detection Resolution Unit: 0 🔽 days 0 🔽 hrs 1 🔽 mins.
	Forensic Algorithm: Monochromatic 💌
	Number of Resolution Units Between Notarizations 1
	Time between notarizations: 0 days 0 hrs 1 mins.
	Number of Notarizations Between Validations 1
	Time between validations: 0 days 0 hrs 1 mins.
	Cost Per Unit: \$ 0.01
	Predicted Cost:
	Tampering Detection:
	Per Day \$: 0.02 Per Year \$: 7.30
	Forensic Analysis (Worst Case):
	One Corruption \$: 0.00 0 🔽 Corruptions \$: 0.00
	Start On: 7 🗸 / 22 🗸 / 2012 🗸 at 23 <table-cell> : 52 🗸</table-cell>
	Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 👻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 💌 hrs 1 💌 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 🗕 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 👻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 🔻 hrs 1 💌 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 👻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🖛 days 0 🖛 hrs 1 🖛 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔻 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🖵 days 0 🖵 hrs 1 🖵 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

📓 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🖛 days 0 🖛 hrs 1 🖛 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: O days O hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

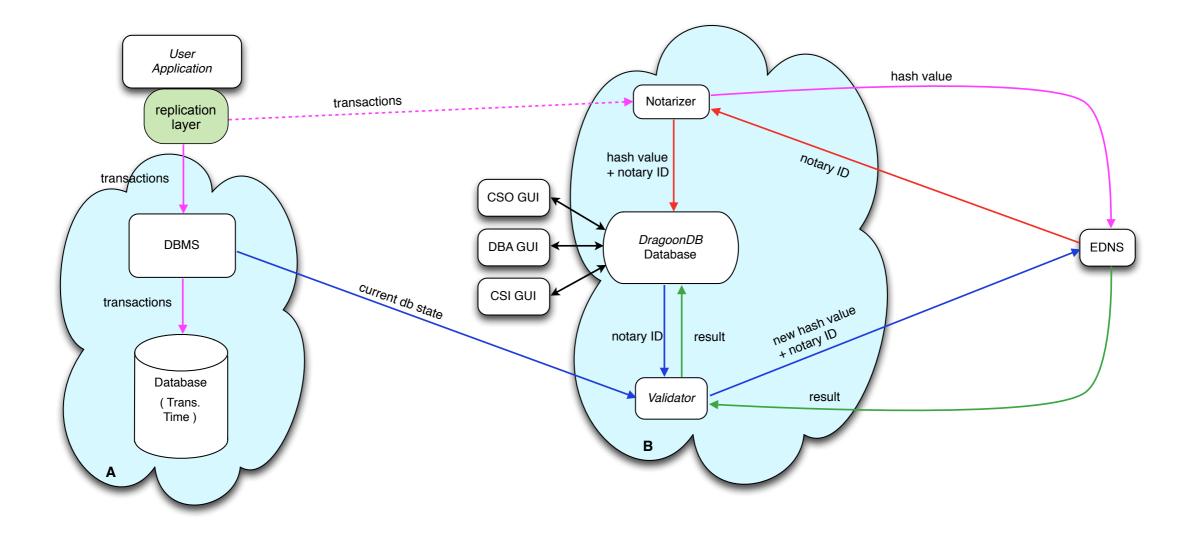
🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🖛 days 0 🖛 hrs 1 🖛 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

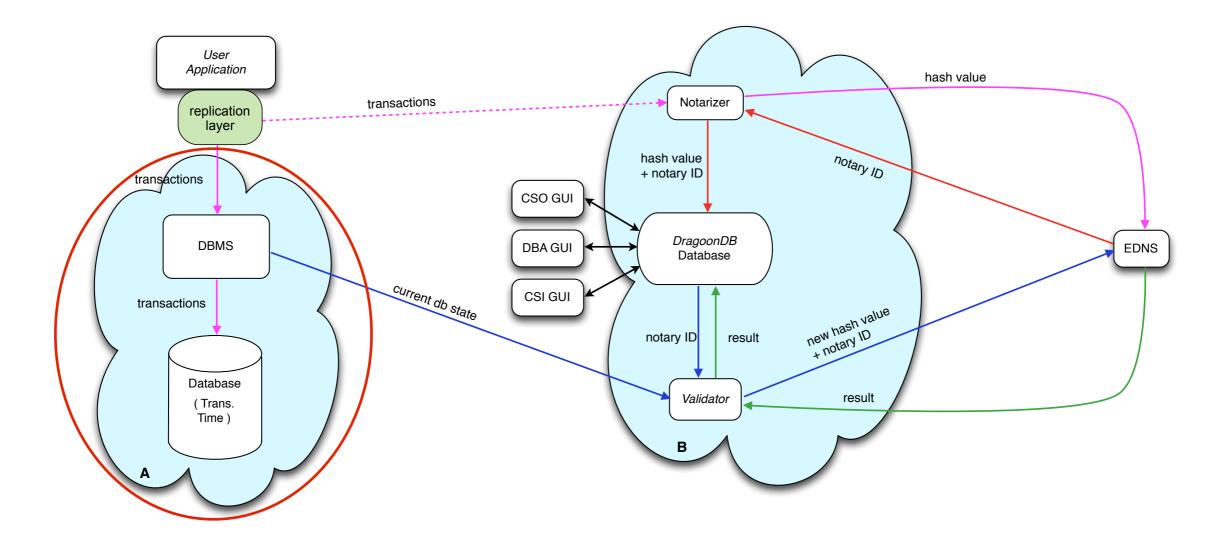
#### **DBA: Database-Specific Settings**

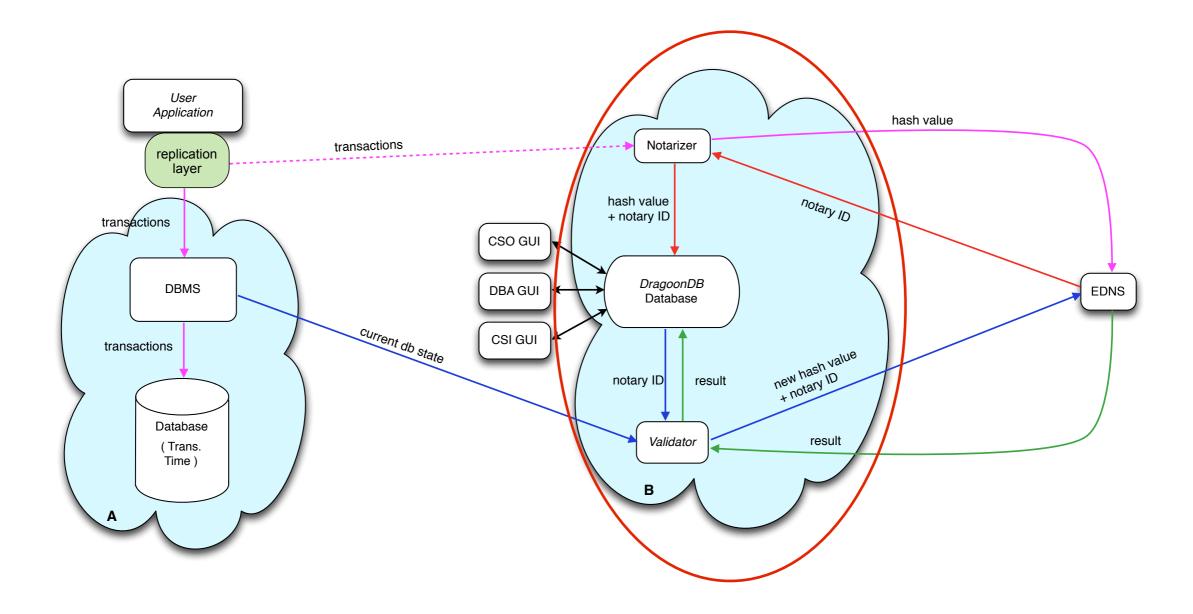
📓 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 🔽 hrs 1 🖵 mins.
Forensic Algorithm: Monochromatic 💌
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼ / 22 ▼ / 2012 ▼ at 23 ▼ : 52 ▼
Save These Settings

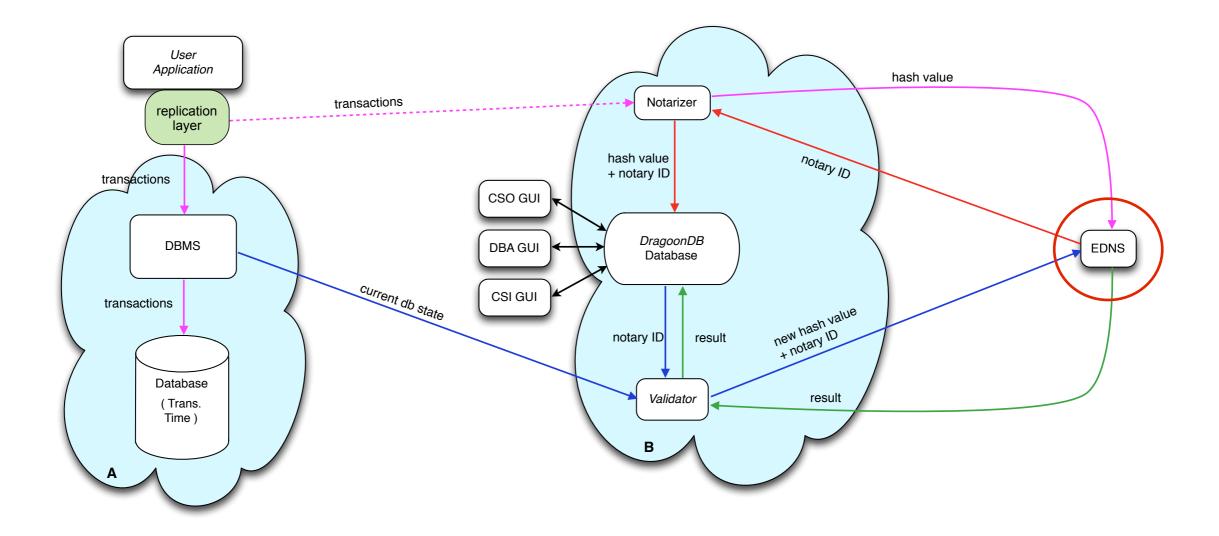
#### **DBA: Database-Specific Settings**

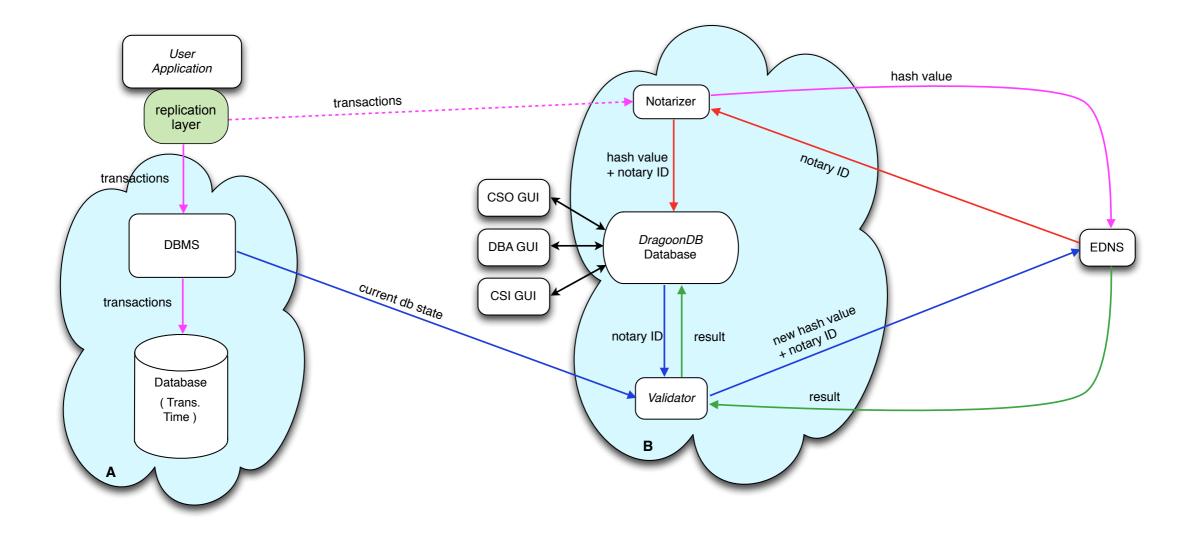
🙆 DBA: Nirav Merchant Email:nirav@email.arizona.edu 📃 🗆 🗙
File
acmedb(/home/tau/software/audit/auditdb) 🔽 CSI: Rick Snodgrass(rts@cs.arizona.edu)
Settings Detected Tampering
Detection Resolution Unit: 0 🔻 days 0 💌 hrs 1 💌 mins.
Forensic Algorithm: Monochromatic 🔽
Number of Resolution Units Between Notarizations 1
Time between notarizations: 0 days 0 hrs 1 mins.
Number of Notarizations Between Validations 1
Time between validations: 0 days 0 hrs 1 mins.
Cost Per Unit: \$ 0.01
Predicted Cost:
Tampering Detection:
Per Day \$: 0.02 Per Year \$: 7.30
Forensic Analysis (Worst Case):
One Corruption \$: 0.00 🛛 🔽 Corruptions \$: 0.00
Start On: 7 ▼/ 22 ▼/ 2012 ▼ at 23 ▼: 52 ▼
Save These Settings

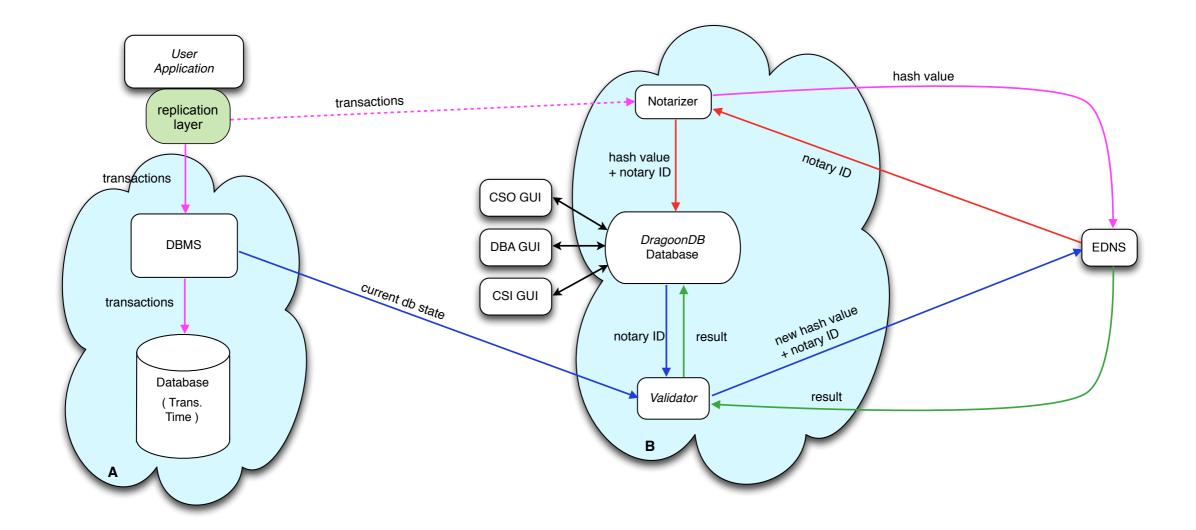












The new threat model may give rise to other temporal concepts.

Also holds for concurrency, replication, and distribution.

# Summary

- Information Accountability
- Reference Architecture & Execution Phases
- Forensic Analysis
- Refinements
- Enterprise Considerations

# Summary (2)

- Need to be able to capture history.
- Need to be able to revisit history.
- Need a trusted witness or at least consensus opinion to provide continuous assurance over time.

# The Challenge

### The Challenge

As we have seen time arises naturally in many aspects of database information accountability (and in many guises).

#### The Challenge

As we have seen time arises naturally in many aspects of database information accountability (and in many guises).

What is the deeper structure of the fundamental connection between temporal databases and information security?

# Thank You!

# Questions?